

STCS Vigyan Vidushi 2025

Cake Cutting

Rohit Vaish

The Model

- The resource: Cake $[0,1]$
- Set of agents $\{1,2,\dots,n\}$
- *Piece of cake*: Finite union of subintervals of $[0,1]$

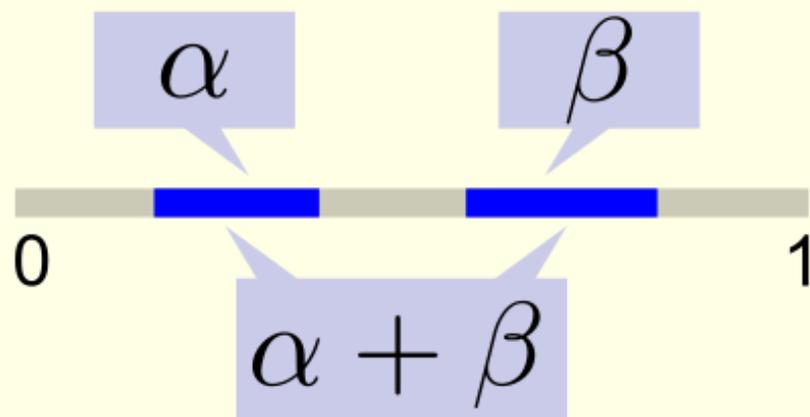


Preferences of Agents

- **Valuation function** v_i : Assigns a non-negative value to any piece of cake

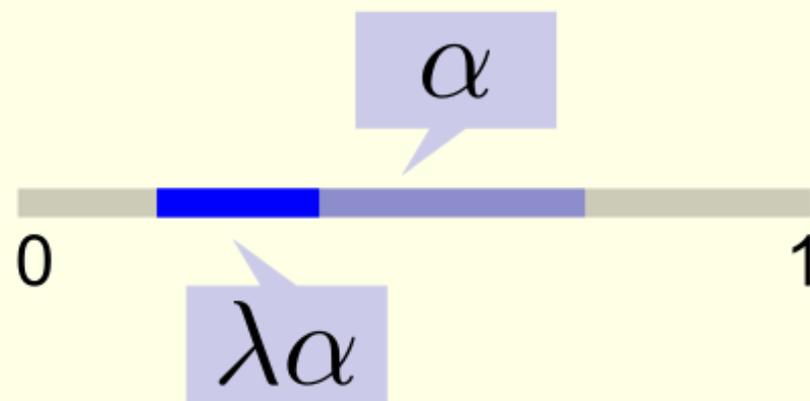
Additivity

for disjoint $X, Y \subseteq [0, 1]$,
 $v_i(X \cup Y) = v_i(X) + v_i(Y)$



Divisibility

for any $X \subseteq [0, 1]$ and any $\lambda \in [0, 1]$,
there exists $Y \subseteq X$ s.t. $v_i(Y) = \lambda v_i(X)$



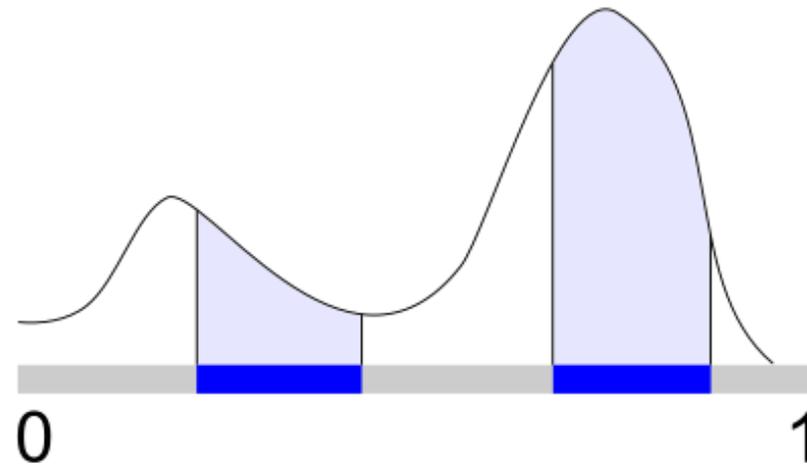
Normalization: for each agent i , $v_i([0, 1]) = 1$.

Preferences of Agents

- **Valuation function** v_i : Assigns a non-negative value to any piece of cake

$$v_i(X) = \int_{x \in X} f_i(x) dx$$

value density function



Fairness notions

- *Allocation/Division*: A partition (A_1, A_2, \dots, A_n) of the cake $[0, 1]$ where each A_i is a piece of cake.



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Proportionality

[Steinhaus, 1948]

for each agent i ,

$$v_i(A_i) \geq \frac{1}{n}$$

Envy-freeness

[Gamow and Stern, 1958; Foley, 1967]

for every pair of agents i, j ,

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For two agents ($n=2$), is one property stronger than the other?

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What about three or more agents?

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EF implies Prop for *any* number of agents

Prop implies EF for *two* agents (but no more)

Robertson-Webb Query Model

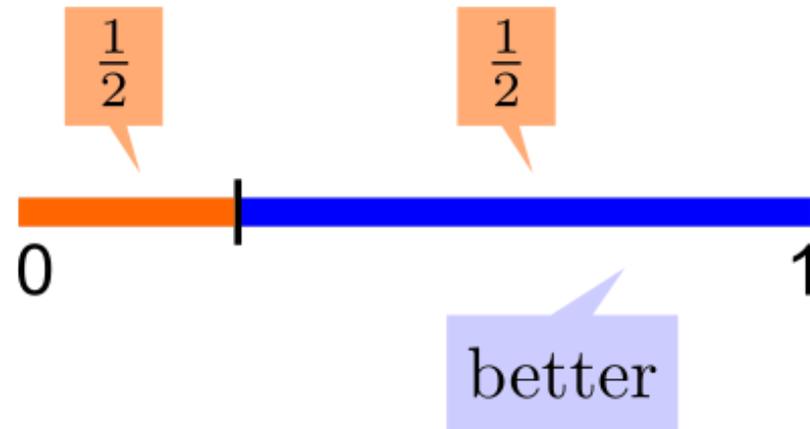
Types of queries that can be used to access the valuation functions

$\text{eval}_i(x, y)$: returns $v_i([x, y])$

$\text{cut}_i(x, \alpha)$: returns y such that $v_i([x, y]) = \alpha$

Cut and Choose

1. Agent 1 cuts the cake into two equally-valued pieces (as per v_1).
2. Agent 2 chooses its preferred piece (as per v_2), and agent 1 gets the remaining piece.



For two agents, an envy-free/proportional cake division can be computed using two queries.

Dubins-Spanier Procedure

A proportional cake division protocol for any number of agents

Dubins-Spanier Procedure

1. A referee gradually moves a knife from left to right.

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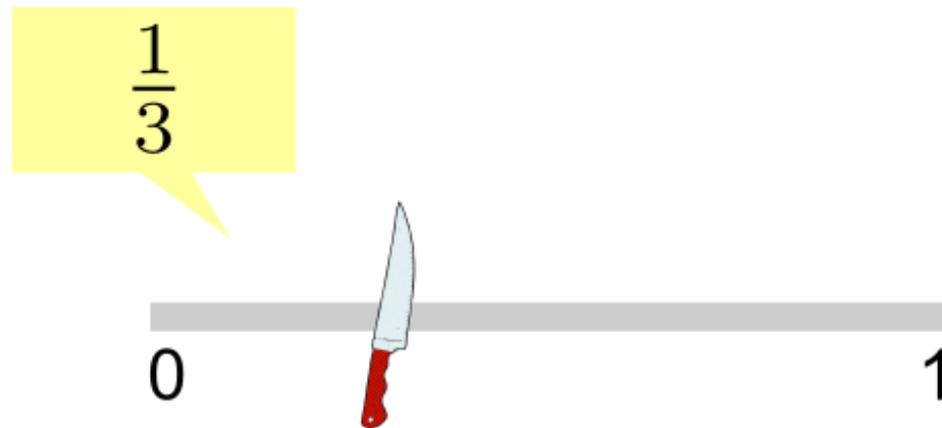
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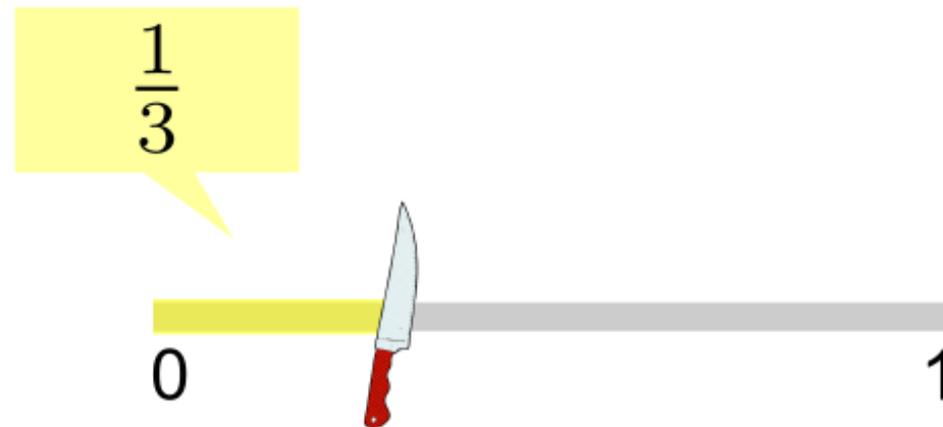
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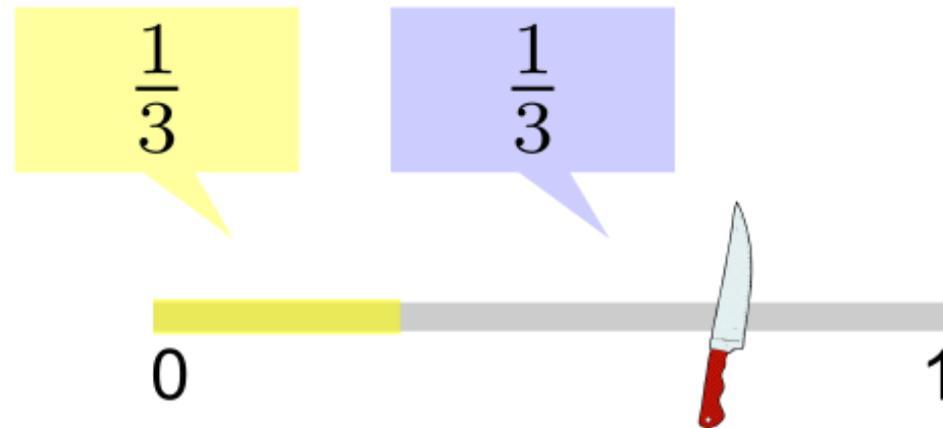
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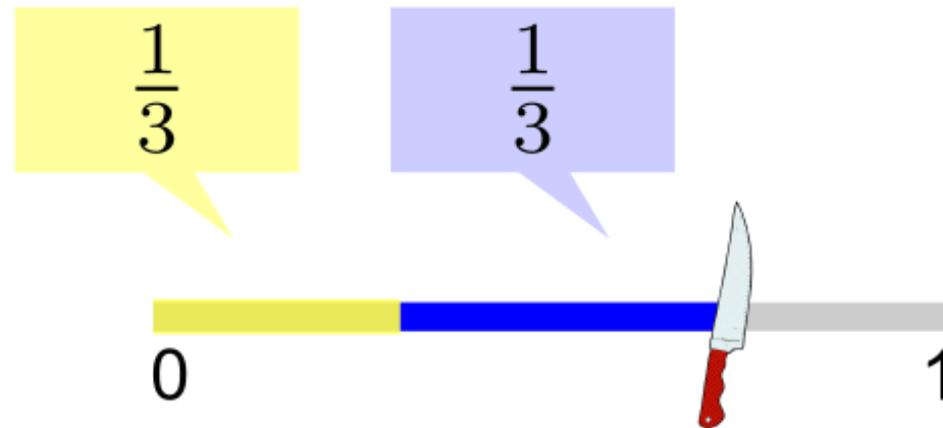
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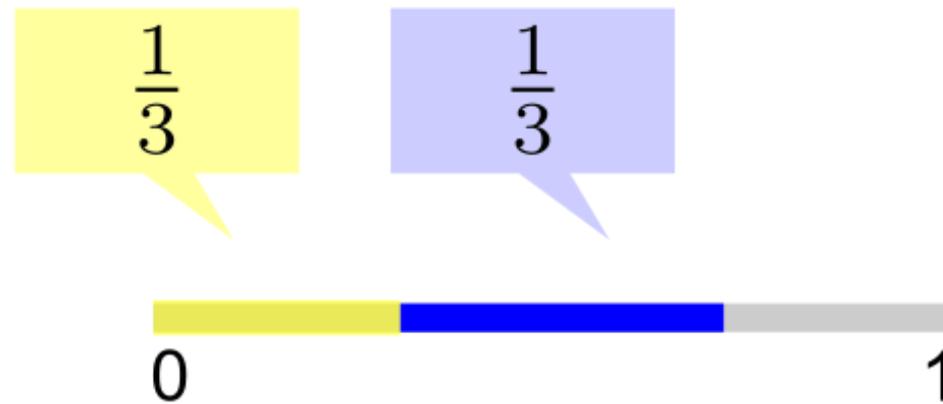
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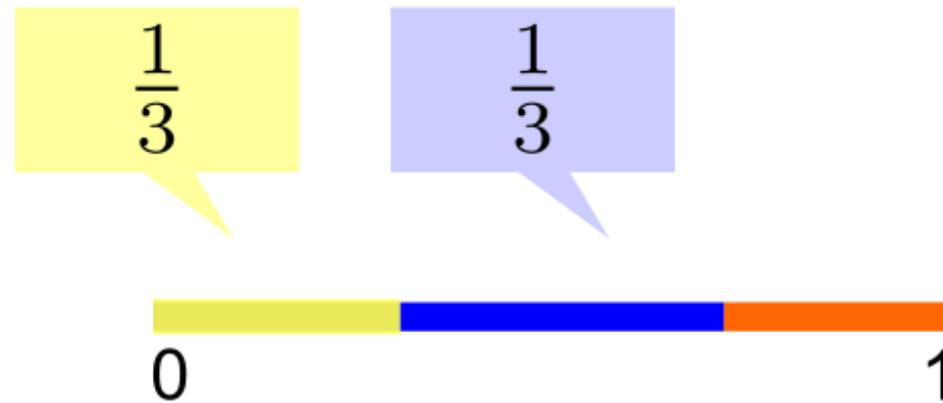
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Why is the resulting allocation proportional?

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Why is the resulting allocation proportional?

Every agent except for the last one gets *exactly* $1/n$.
The last agent gets *at least* $1/n$.

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Can this procedure be implemented in the Robertson-Webb model?

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Yes!

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Query complexity in the Robertson-Webb model?

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Query complexity in the Robertson-Webb model?

$\mathcal{O}(n^2)$ queries (Exercise)

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For n agents, a proportional cake division can be computed using $O(n^2)$ queries.

The Story of Proportionality

The Story of Proportionality

query complexity



The Story of Proportionality

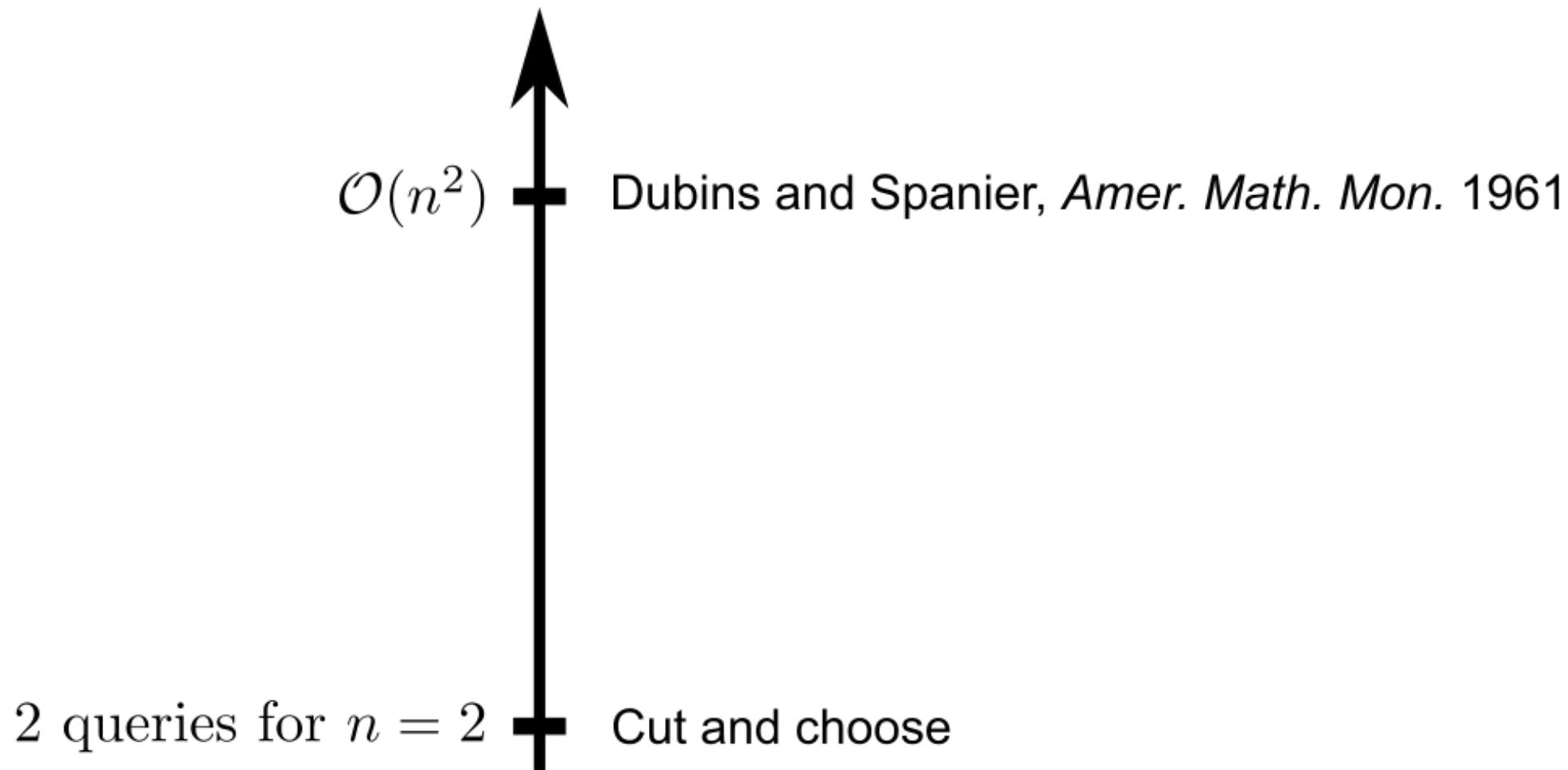
query complexity



2 queries for $n = 2$ + Cut and choose

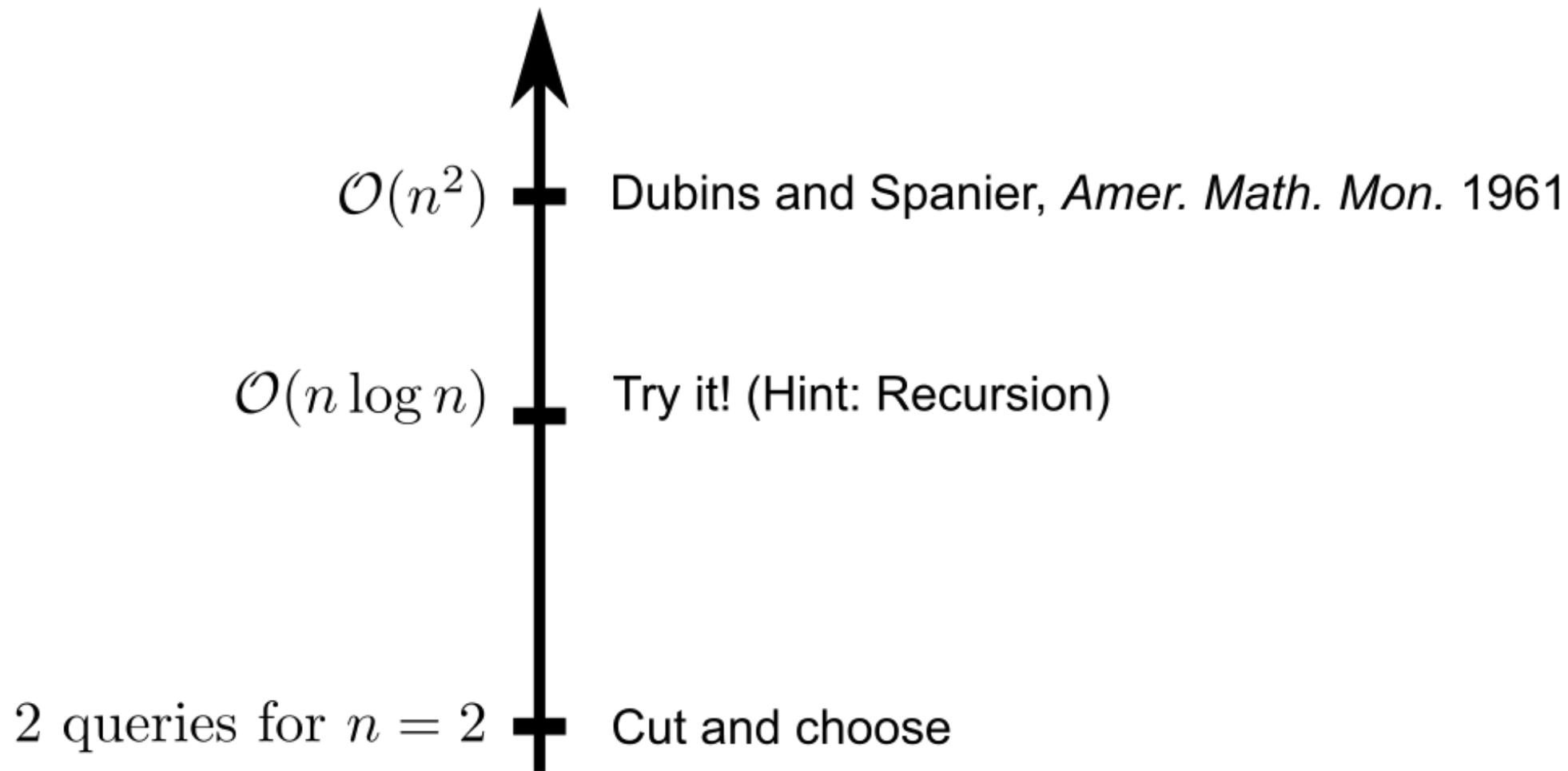
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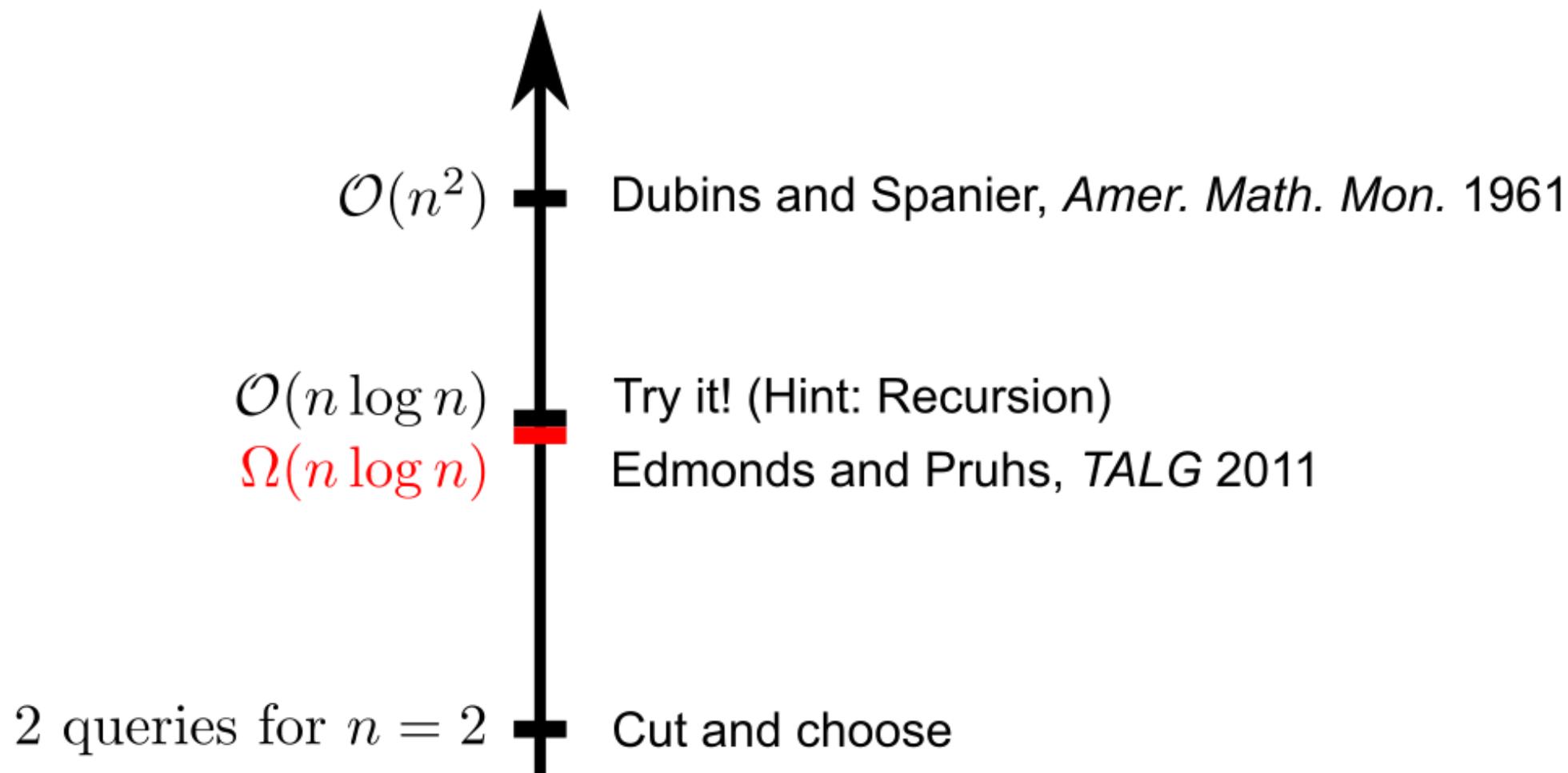
The Story of Proportionality

query complexity



The Story of Proportionality

query complexity



The Story of Envy-freeness



Selfridge-Conway Procedure

An envy-free cake division protocol for three agents

Envy-free for three

Envy-free for three

A



B

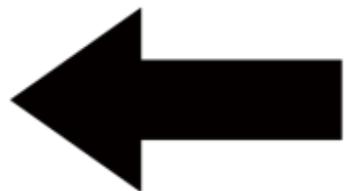


C



Envy-free for three

A



B



C



Envy-free for three

A



B



C



Envy-free for three

A

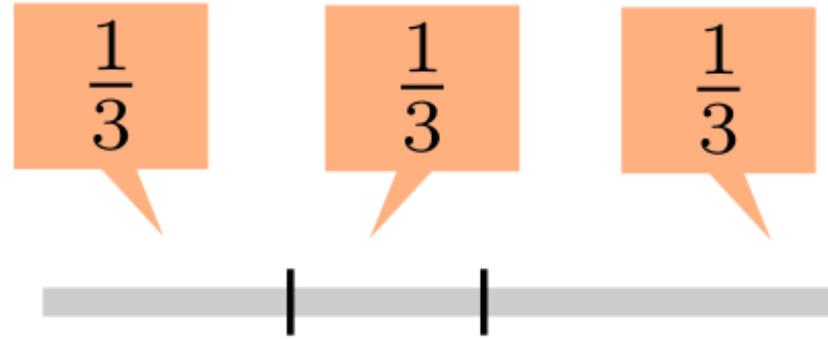


equal in
my view

B



C



Envy-free for three

A



B



C



Envy-free for three

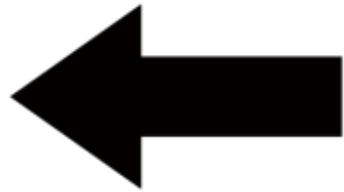
A



B



C



Envy-free for three

A



B



two-way
tie

C



Envy-free for three

A



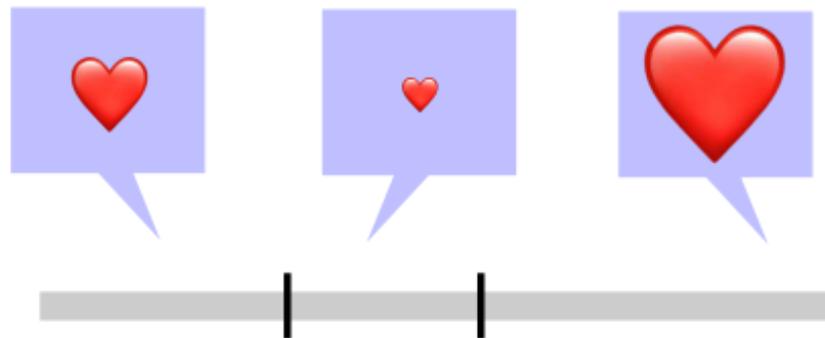
B



C



two-way
tie



Envy-free for three

A

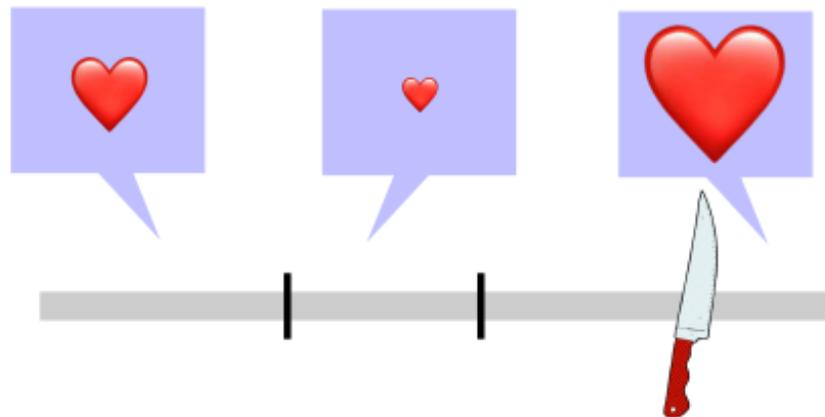


B



two-way
tie

C



Envy-free for three

A

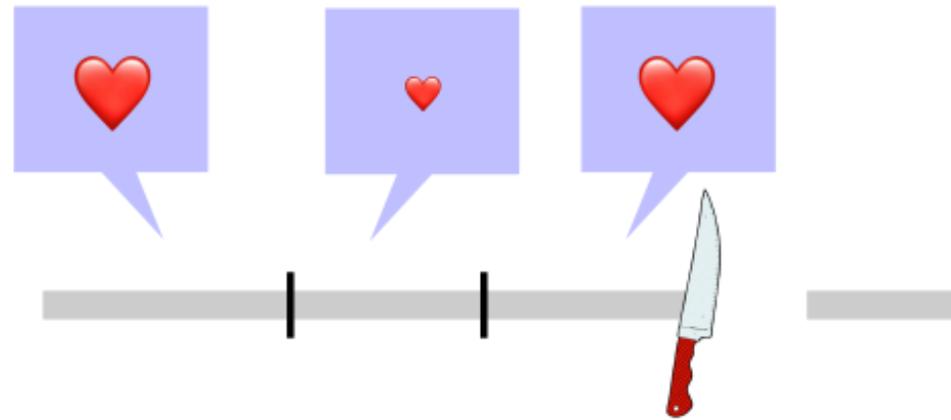


B



two-way
tie

C



Envy-free for three

A

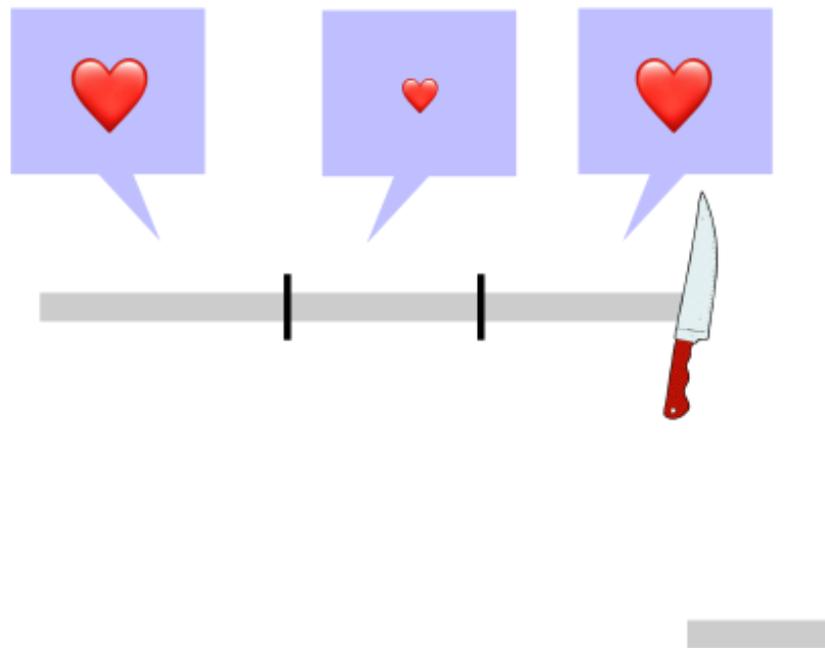


B



two-way
tie

C

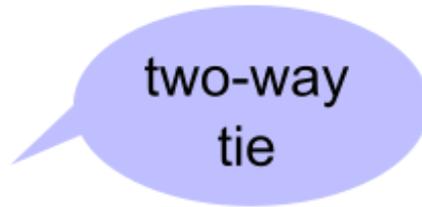


Envy-free for three

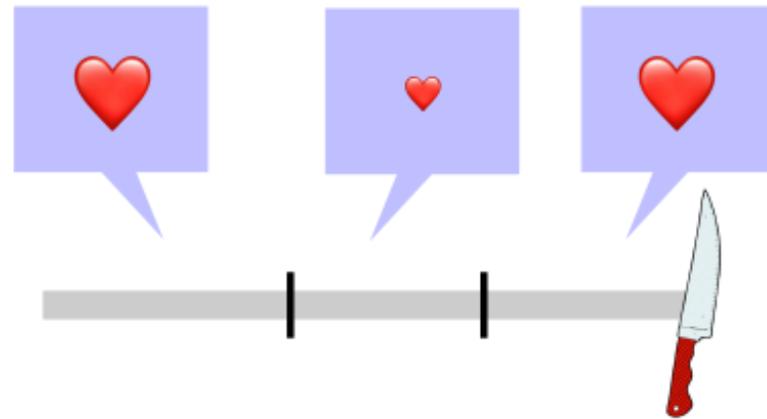
A



B



C



Trimmings

Envy-free for three

A

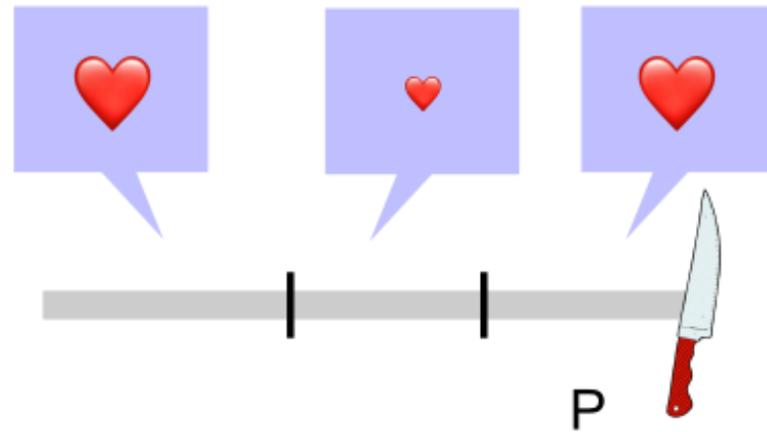


B



two-way
tie

C



Trimmings

Envy-free for three

A



B



C



Trimmings

Envy-free for three

A



B



C



1st



Trimmings

Envy-free for three

A



B



2nd

C



1st



Trimmings

Envy-free for three

A



3rd

B



2nd

C



1st



Trimmings

Envy-free for three

A



3rd

B



2nd

I pick P if C doesn't

C



1st



Trimmings

Envy-free for three

A



3rd

B



2nd

I pick P if C doesn't

C



1st



Trimmings

Envy-free for three

A



3rd

B



2nd

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C



1st



Trimmings

Envy-free for three

A



3rd

B



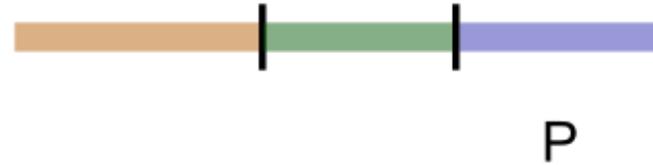
2nd

I pick P if C doesn't

C



1st



Trimmings

Envy-free for three

A



EF because
of equal cuts

B



EF because
of two-way tie

C



EF because
I picked first



Trimmings

Envy-free for three

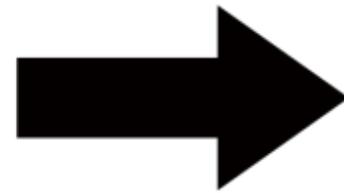
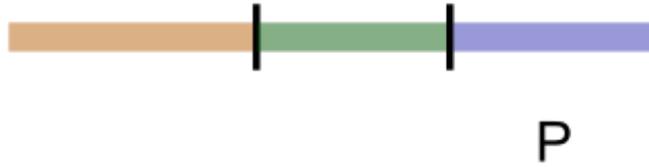
A



B



C



Trimmings

Envy-free for three

A



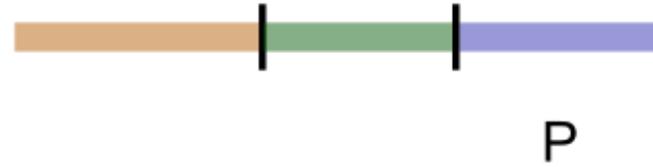
B



C



I equidivide



Trimmings

Envy-free for three

A



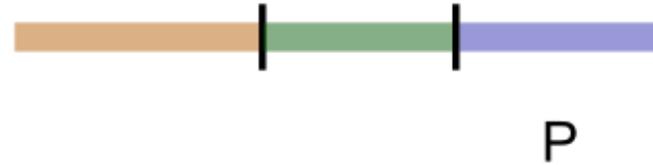
B



C



I equidivide



Envy-free for three

A



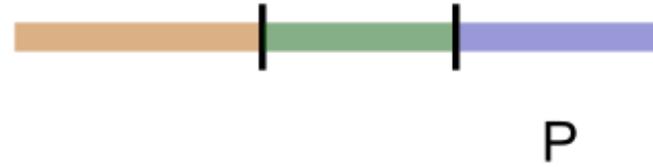
B



C



I equidivide



Trimmings

Envy-free for three

A



B

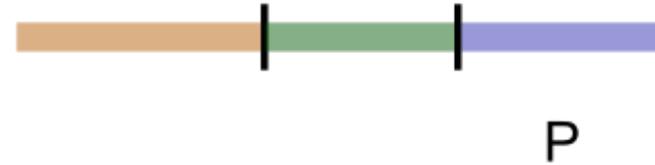


I pick first
yay!

C



I equidivide



Trimmings

Envy-free for three

A



I can pick after B does

B



I pick first yay!

C



I equidivide



 Trimmings

Envy-free for three

A



I can pick after B does

B



I pick first yay!

C



I equidivide and pick last



 Trimmings

Envy-free for three

A



I can pick after B does

B



I pick first yay!

C



I equidivide and pick last



 Trimmings

Envy-free for three

A



I can pick after B does

B



I pick first yay!

C



I equidivide and pick last



 Trimmings

Envy-free for three

A



I can pick after B does

B



I pick first yay!

C



I equidivide and pick last



 Trimmings

Envy-free for three

A



Irrevocable
advantage

B



EF because
I picked first

C



EF because
I equidivided



 Trimmings

Exercise

How many queries does the three-person EF protocol require?

The Story of Envy-freeness

The Story of Envy-freeness

query complexity



2 queries for $n = 2$  Cut and choose

The Story of Envy-freeness

query complexity



$\mathcal{O}(1)$ queries for $n = 3$ — Selfridge-Conway

2 queries for $n = 2$ — Cut and choose

The Story of Envy-freeness

query complexity

A finite but *unbounded* protocol

Brams and Taylor, *Amer. Math. Mon.* 1995

$\mathcal{O}(1)$ queries for $n = 3$

Selfridge-Conway

2 queries for $n = 2$

Cut and choose



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$\Omega(n^2)$

Procaccia, *IJCAI* 2009

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Cut and choose



The Story of Envy-freeness

query complexity

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$\mathcal{O}(1)$ queries for $n = 3$

Selfridge-Conway

2 queries for $n = 2$

Cut and choose

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Cut and choose

References

- Introduction to cake-cutting algorithms.

Ariel Procaccia

“*Cake Cutting Algorithms*”

Chapter 13 in Handbook of Computational Social Choice

- Lecture by Ariel Procaccia on “Cake cutting” in the *Optimized Democracy* course.

<https://sites.google.com/view/optdemocracy/schedule>

