

## Lecture 6

# House Allocation

# Housing Market

[Shapley and Scarf, *JME* 1974]

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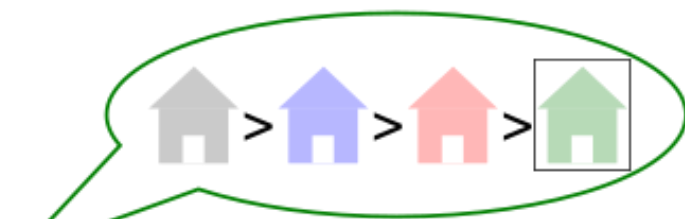
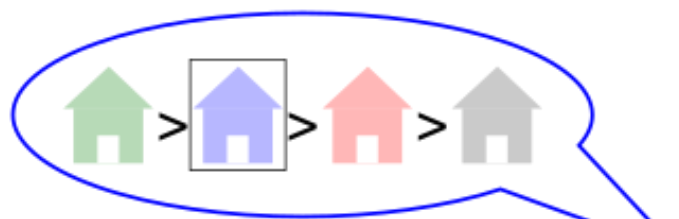
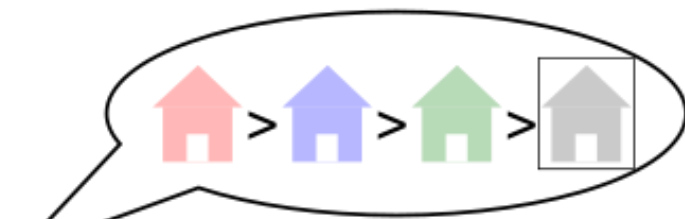
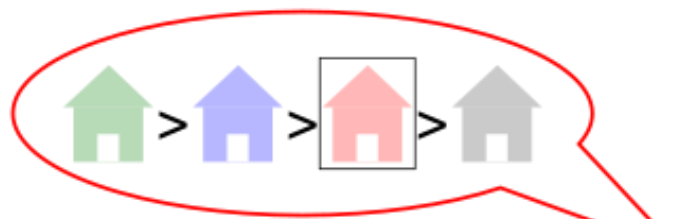
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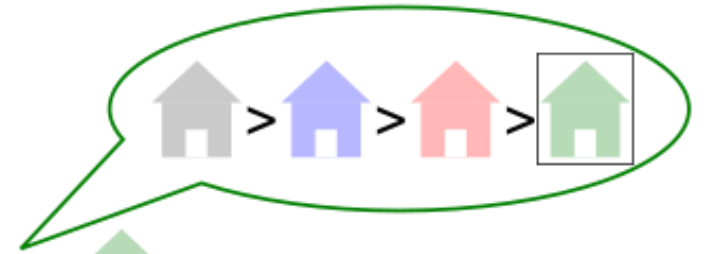
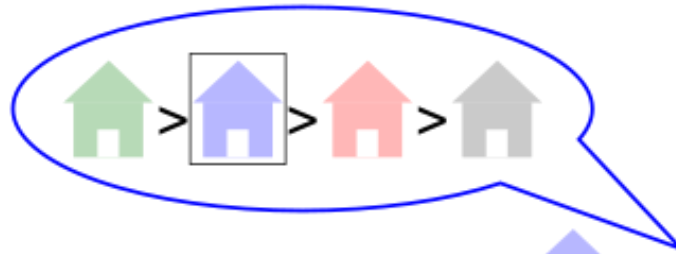
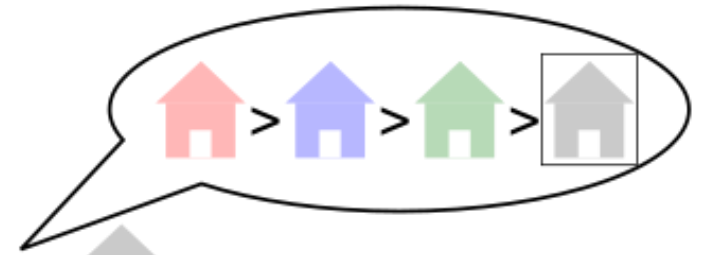
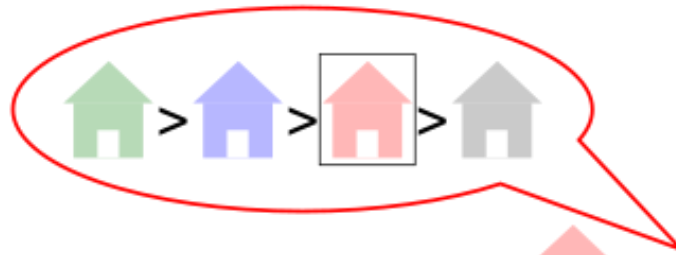
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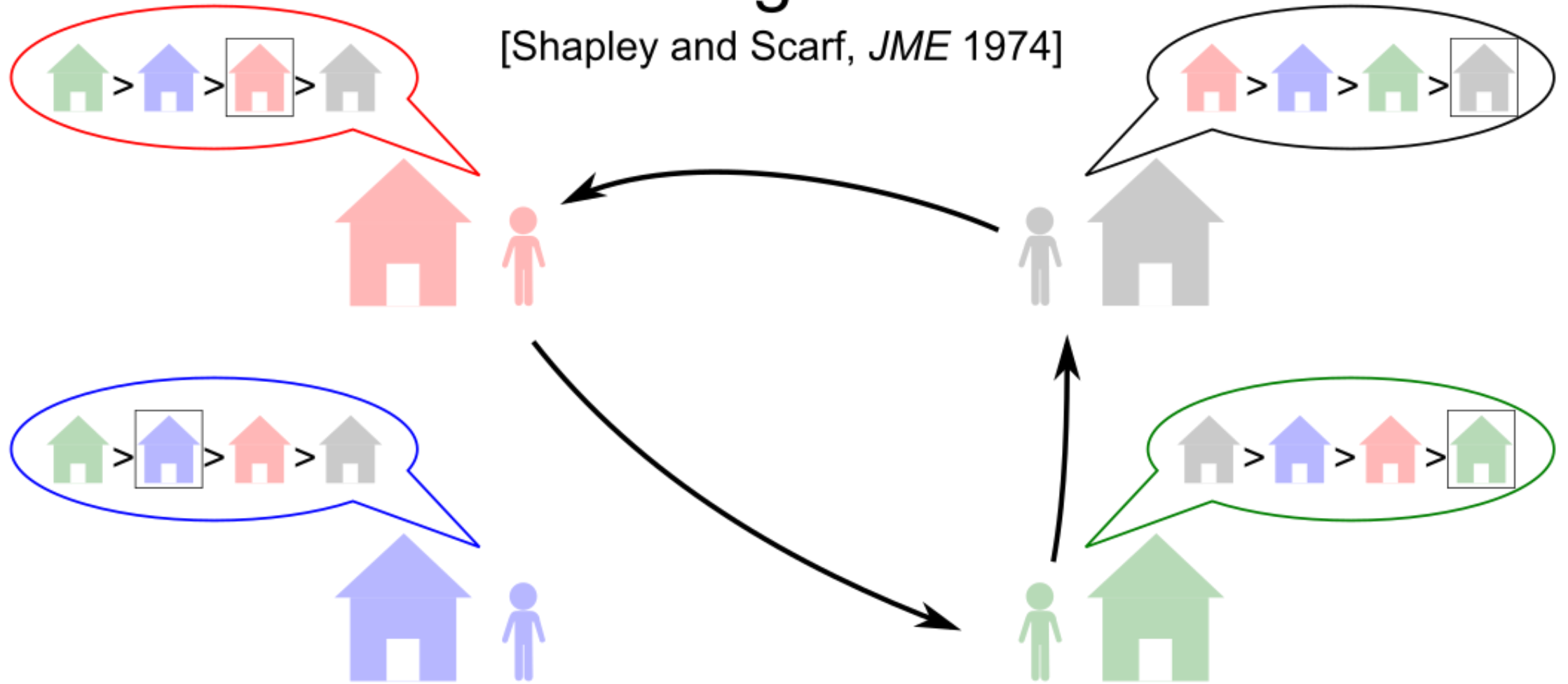
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Can't use money. Only way to make agents happier is via *exchanges*.

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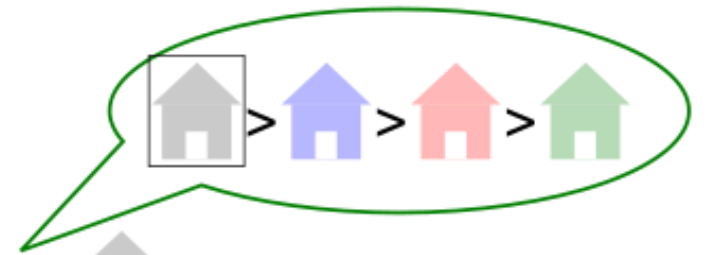
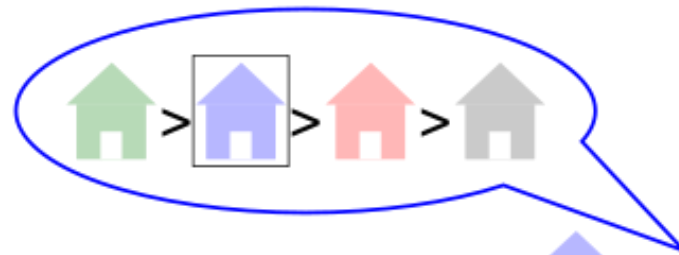
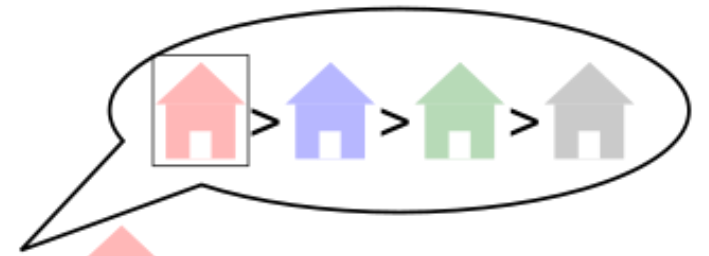
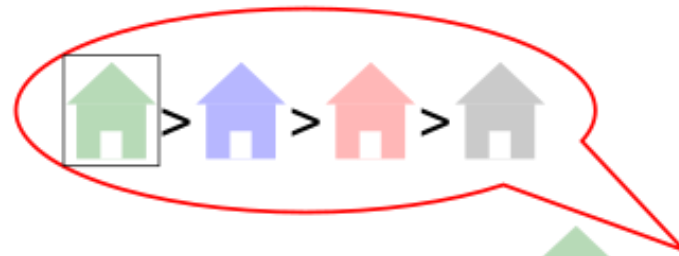
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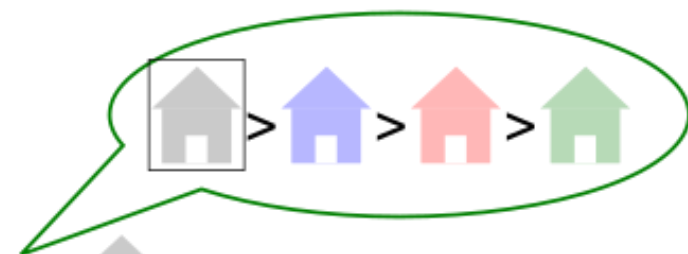
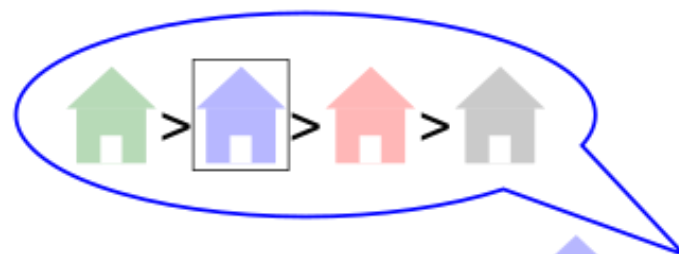
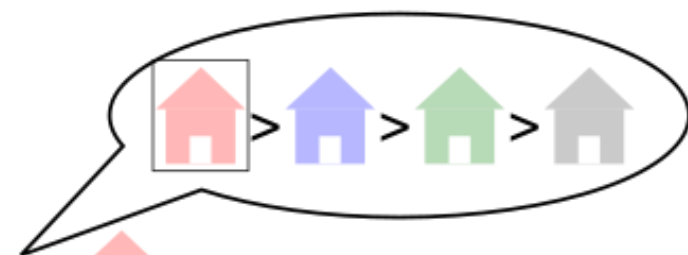
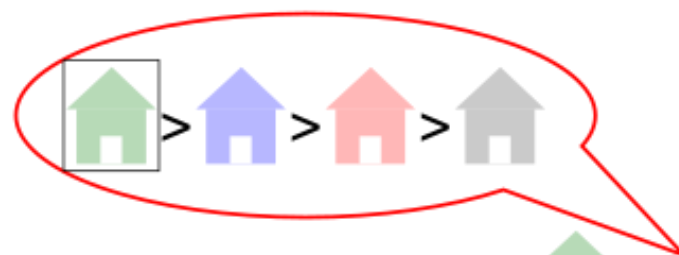


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# Housing Market

[Shapley and Scarf, *JME* 1974]



Is there a way to exchange houses to make the agents "maximally happy"?

# Top-Trading Cycle Algorithm (TTCA)

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

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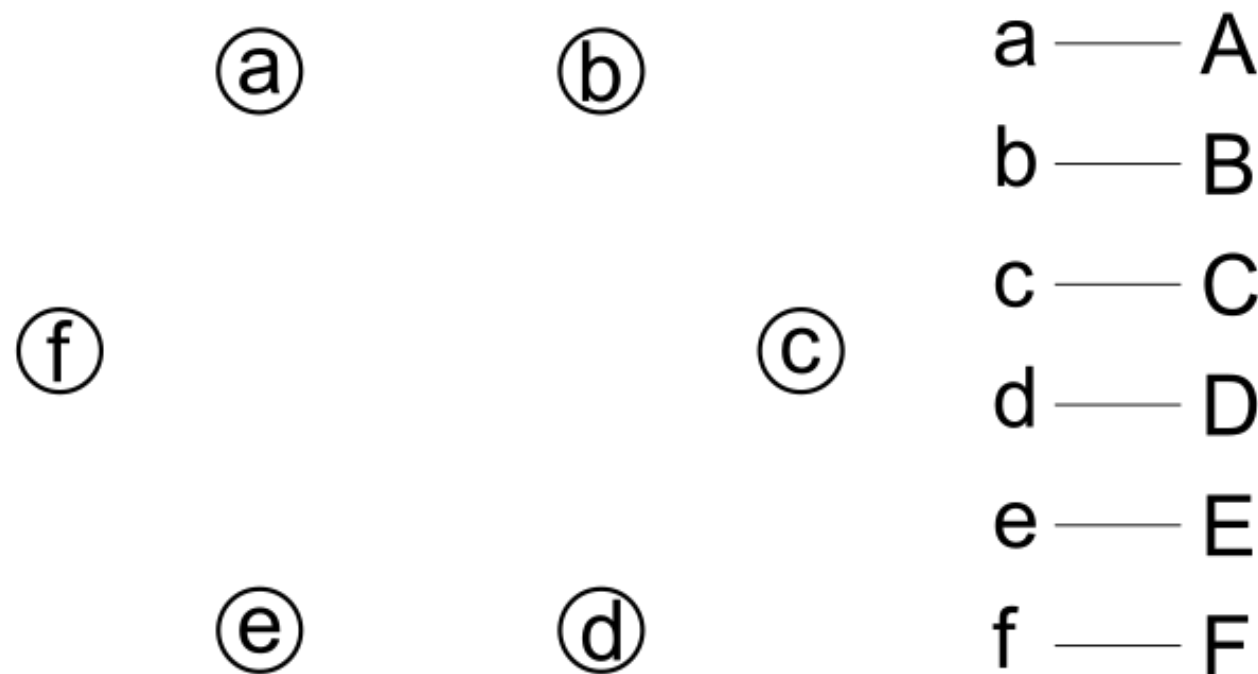
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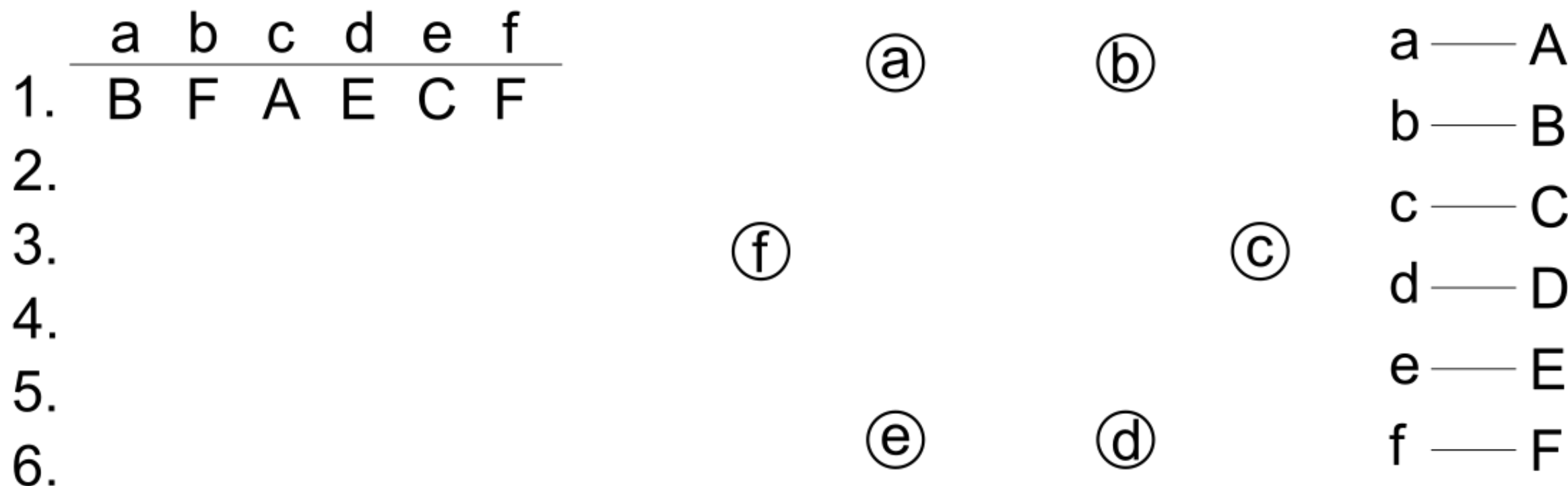


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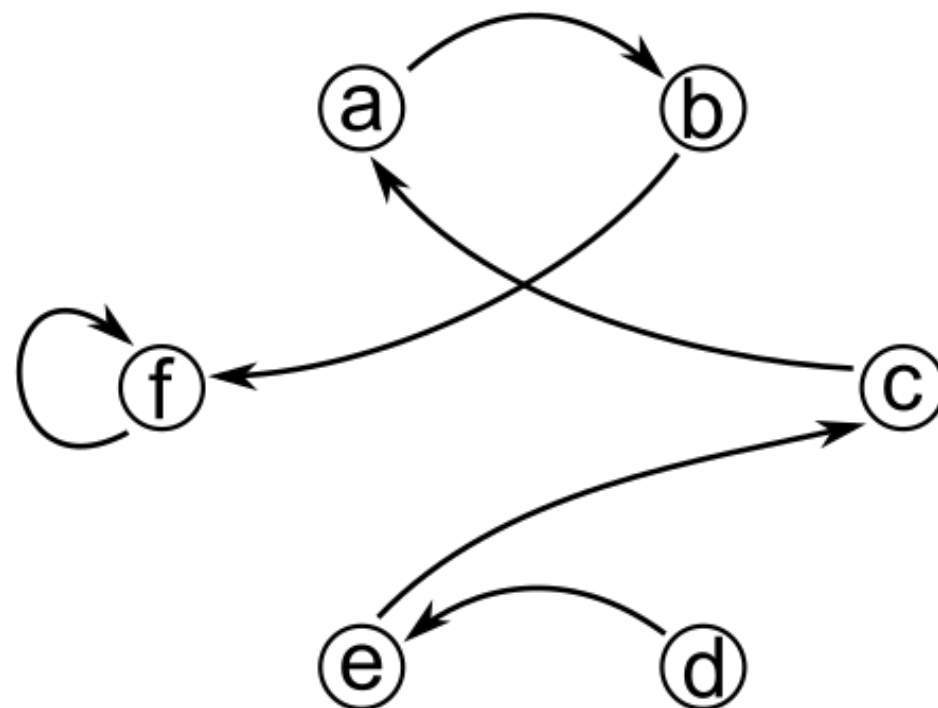
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3.						
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c	—	C
d	—	D
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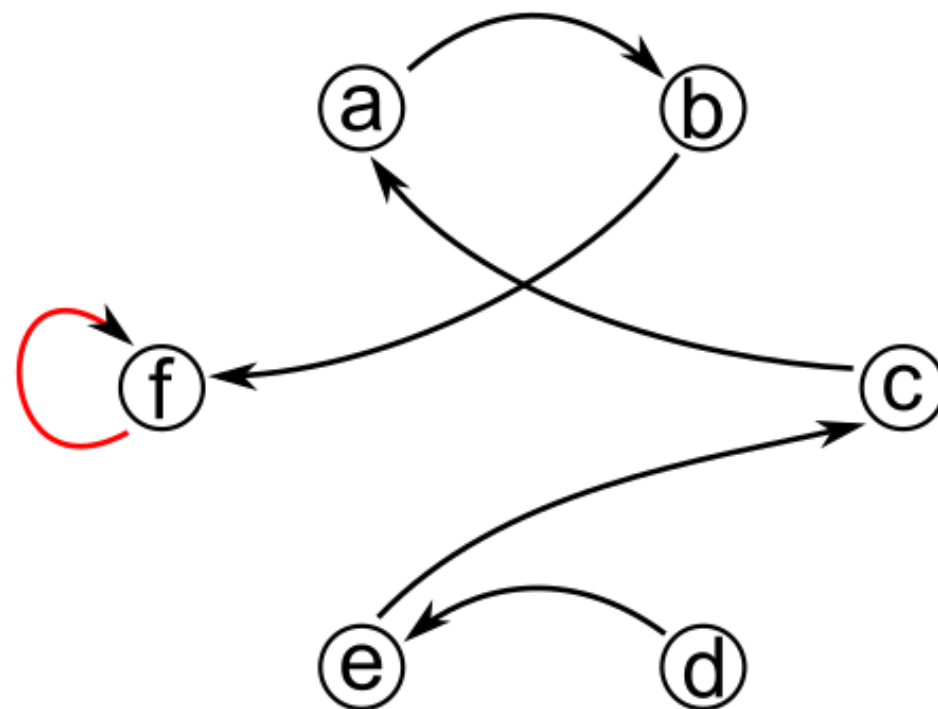
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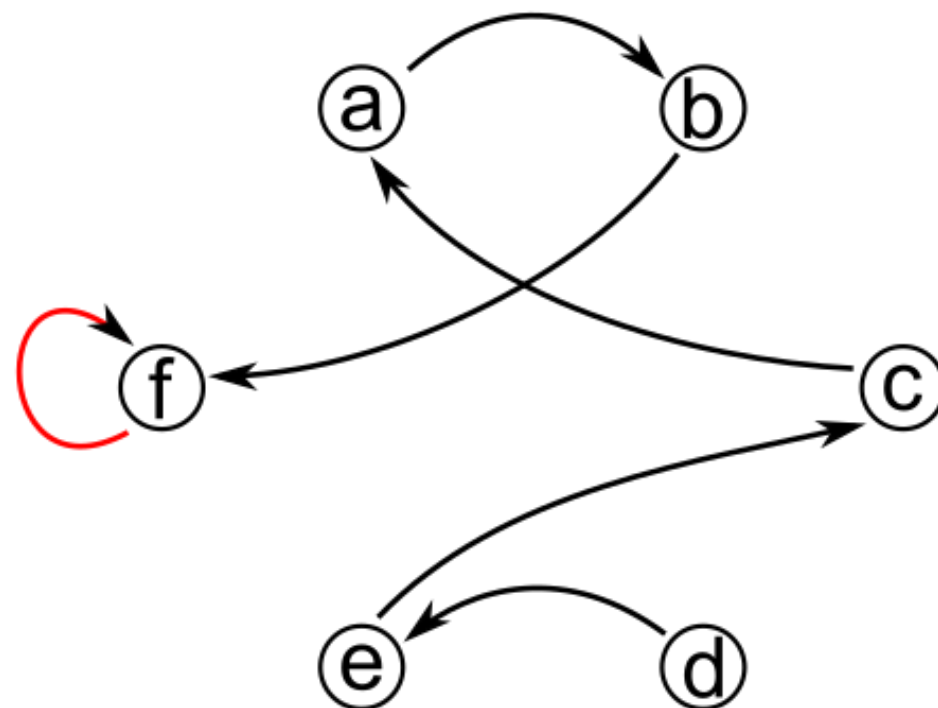
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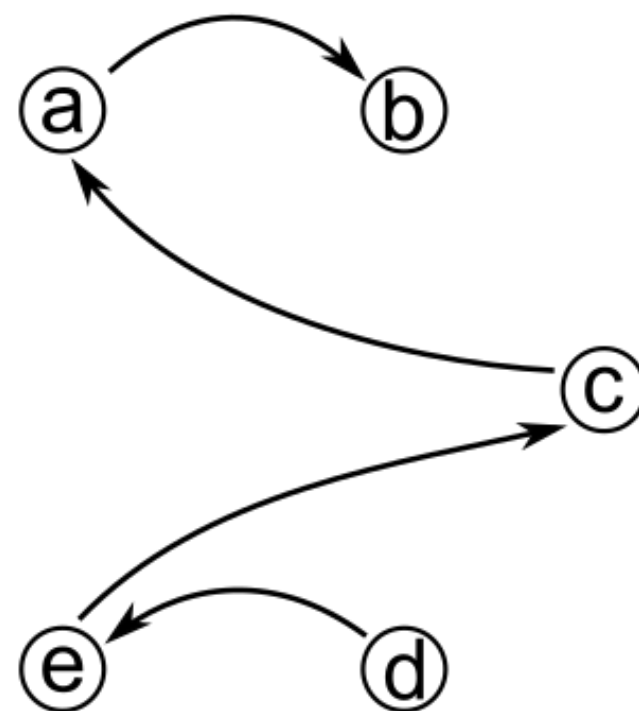
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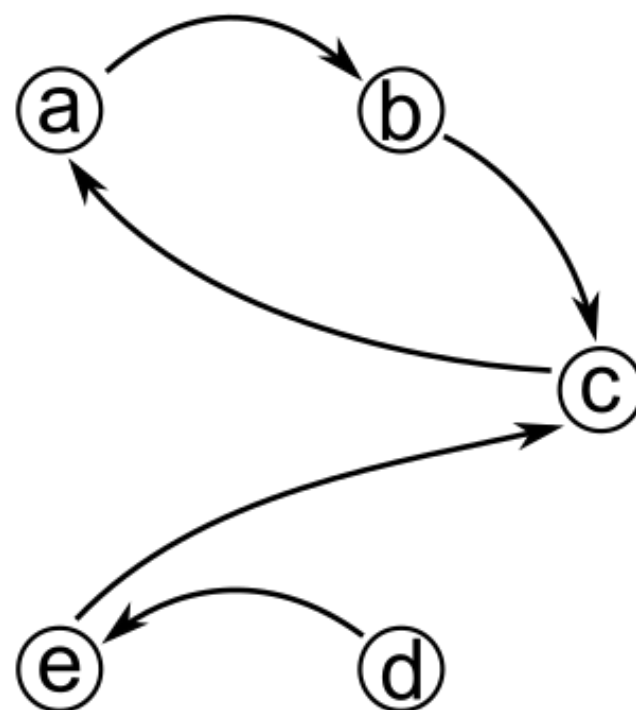
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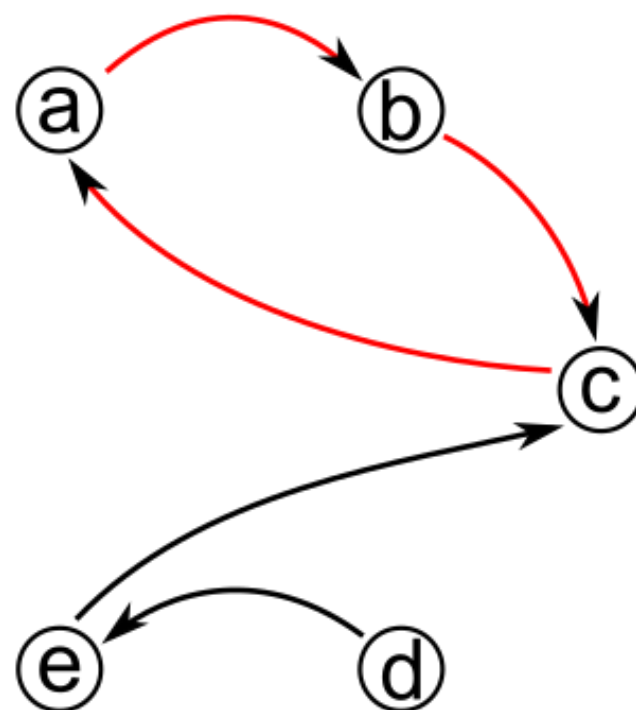
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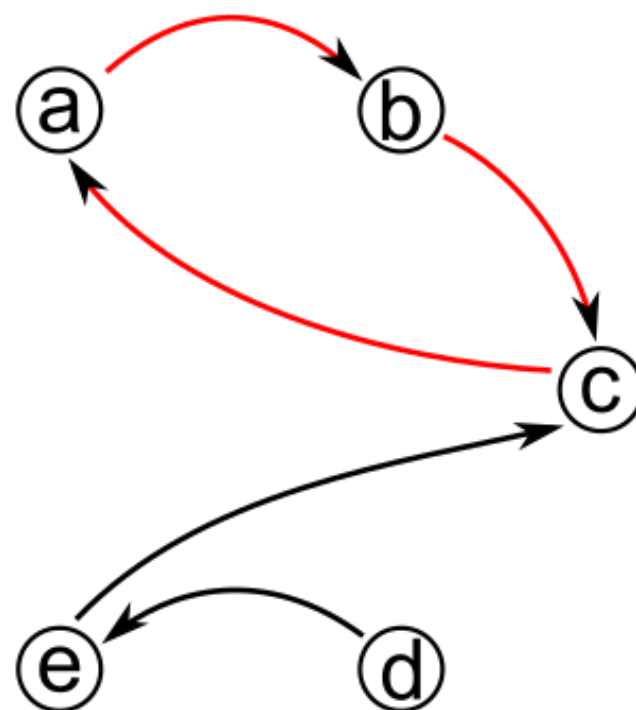
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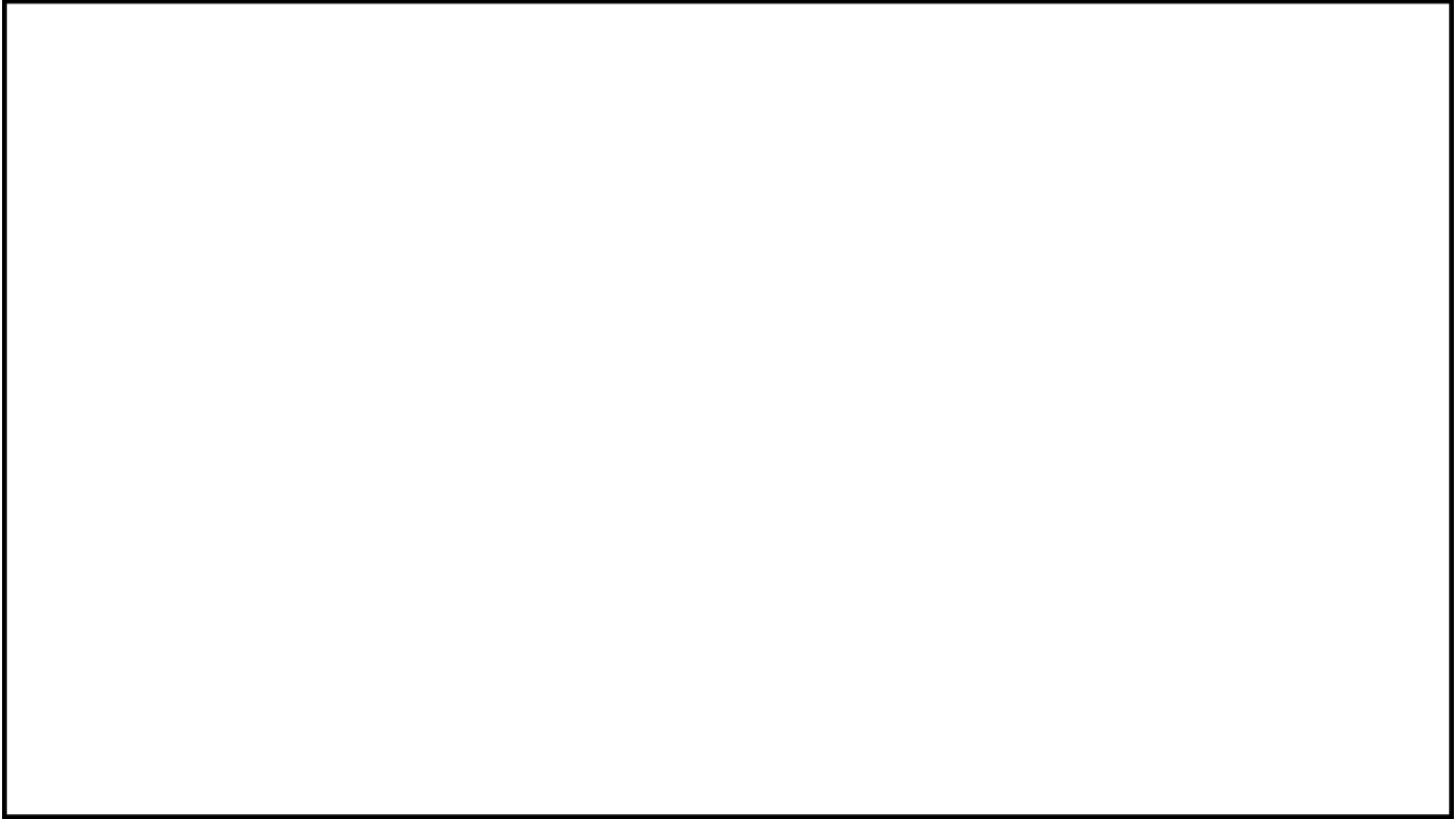
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Imagine a graph where the *vertices* are the (remaining) agents and the *edge*  $i \rightarrow j$  means that agent  $i$ 's favorite house is owned by agent  $j$ .

At each step, each vertex has an outgoing edge.

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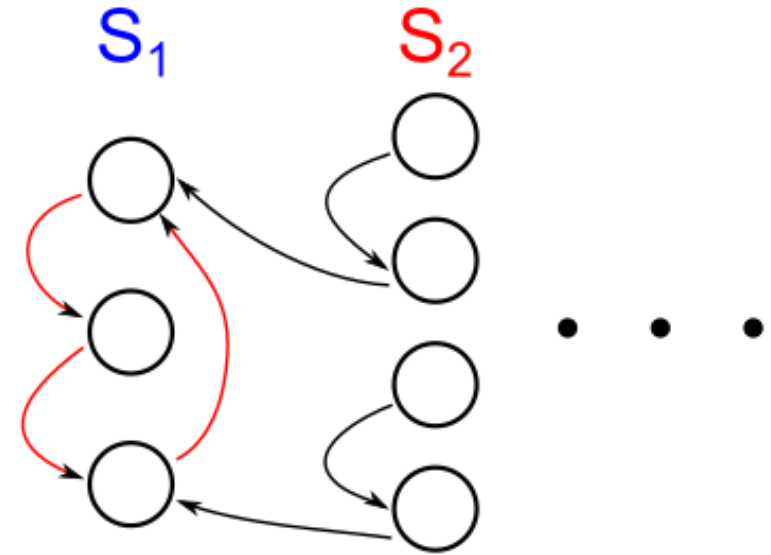
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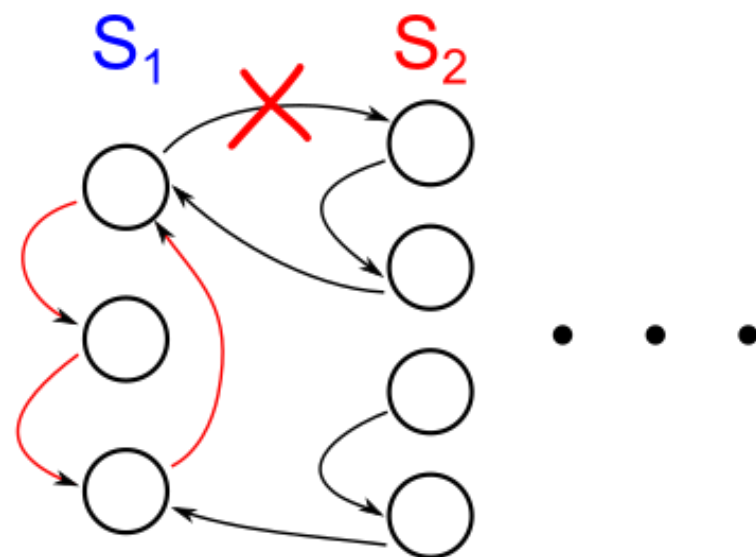
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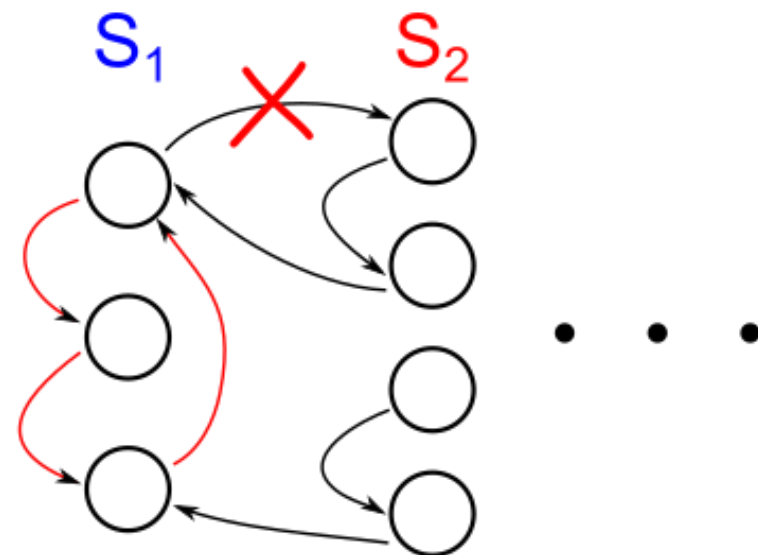


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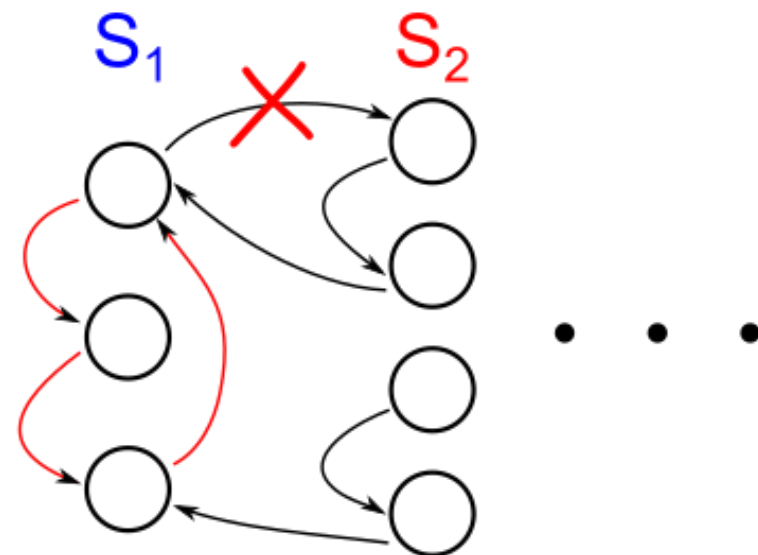
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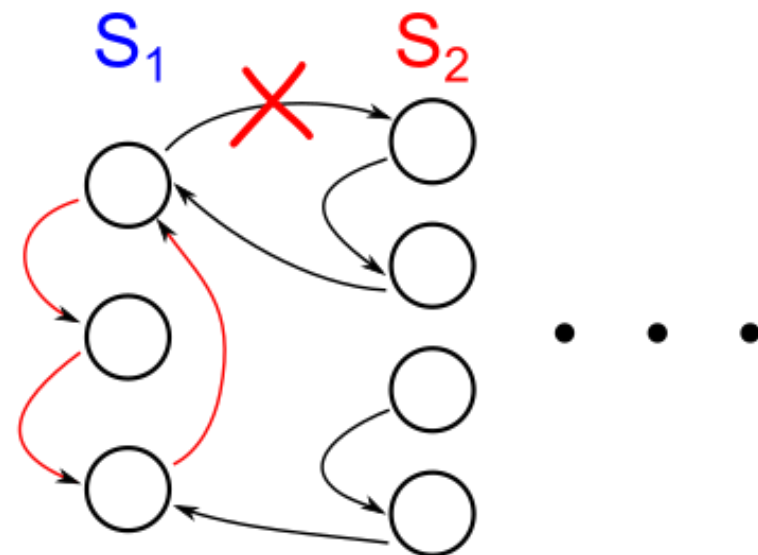


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Is TTCA group-strategyproof?

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Can a coalition of agents misreport their preferences such that no one in the coalition is worse off and someone is strictly better off w.r.t. their true preferences?

# Is TTCA group-strategyproof?

Monday, September 23, 2024

## [A 40 year old proof about top trading cycles is corrected \(by two Berkeley grad students\)](#)

Science (and math) can be self-correcting, sometimes slowly. Here's an article that corrects the first proof that the top trading cycles algorithm is group strategy proof. That's a true result, with multiple subsequent proofs. But apparently the first proof presented wasn't the best one. That's good to know.

One reason this may have taken a long time to spot is that the result is correct, and that there are subsequent proofs that connect the result to properties of other mechanisms.


[Will Sandholtz](#) and [Andrew Tai](#), the authors, did this work as Ph.D. students at UC Berkeley. (good for them!)

[Group incentive compatibility in a market with indivisible goods: A comment](#) by Will Sandholtz and Andrew Tai

### "Highlights

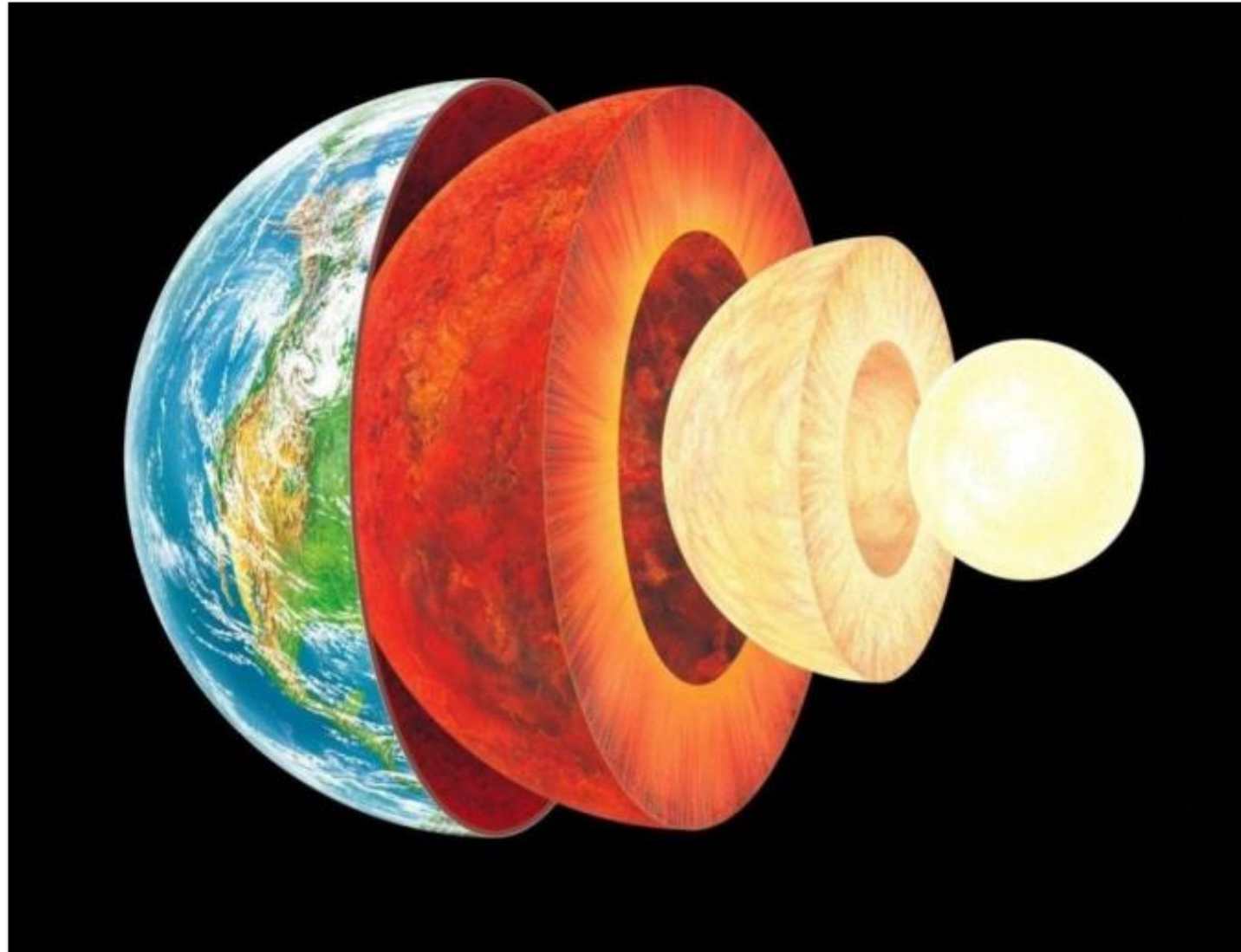
- Bird (1984), first to show top trading cycles is group strategy-proof, has errors.
- We present corrected results and proofs.
- We present a novel proof of strong group strategy-proofness without non-bossiness.

"Abstract: We note that the proofs of Bird (1984), the first to show group strategy-proofness of top trading cycles (TTC), require correction. We provide a counterexample to a critical claim and present corrected proofs in the spirit of the originals. We also present a novel proof of strong group strategy-proofness using the corrected results."

Posted by Al Roth at [5:58 AM](#) 

Labels: [mathematics](#), [papers](#), [peer review](#), [science](#), [TTC](#)

# The Core



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An allocation is in the **core** if no coalition blocks it.

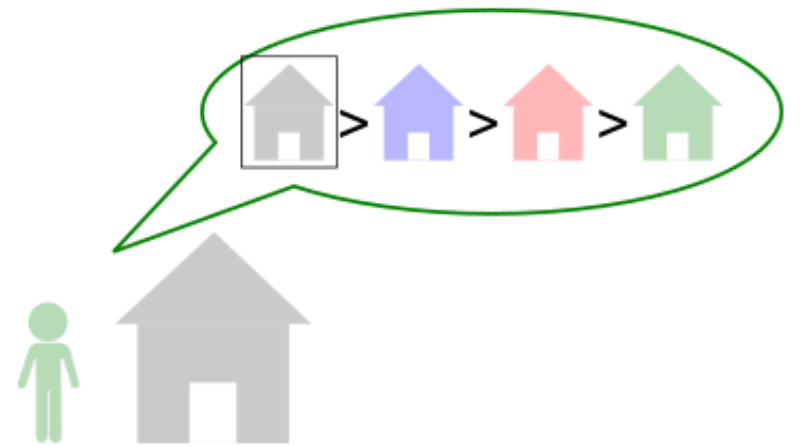
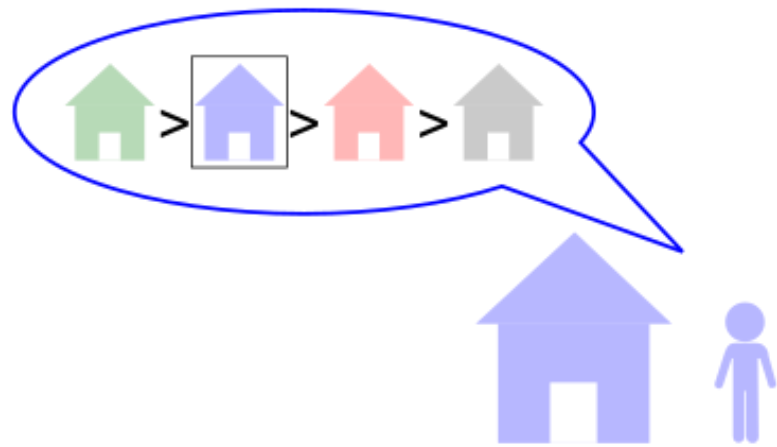
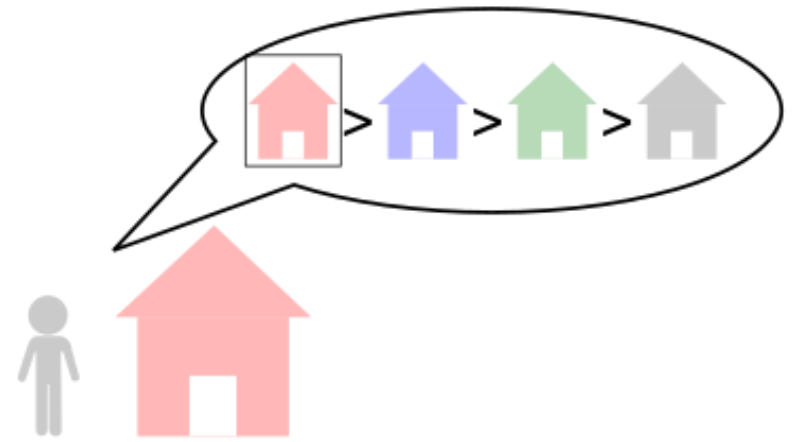
# The Core

An allocation is in the **core** if no coalition blocks it.

A coalition of agents **blocks** an allocation **A** if:

they can redistribute their **endowed houses** among themselves such that, compared to **A**, none of them is worse off and at least one of them is strictly better off (i.e., redistributing endowments is a Pareto improvement over **A**).

# Example of a Core Allocation





**TTCA outcome is in the core.**

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

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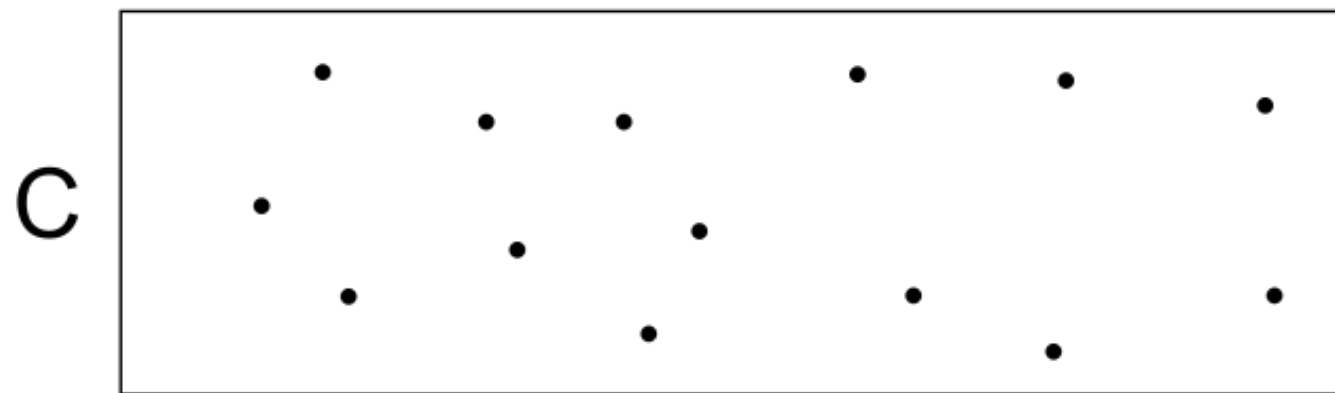
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Suppose not. Then, a coalition  $C$  of agents must block the TTCA allocation  $T$ . Let  $R$  be a redistribution of endowments among agents in  $C$  that they find Pareto better than  $T$ .

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[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

Suppose not. Then, a coalition  $C$  of agents must block the TTCA allocation  $T$ . Let  $R$  be a redistribution of endowments among agents in  $C$  that they find Pareto better than  $T$ .

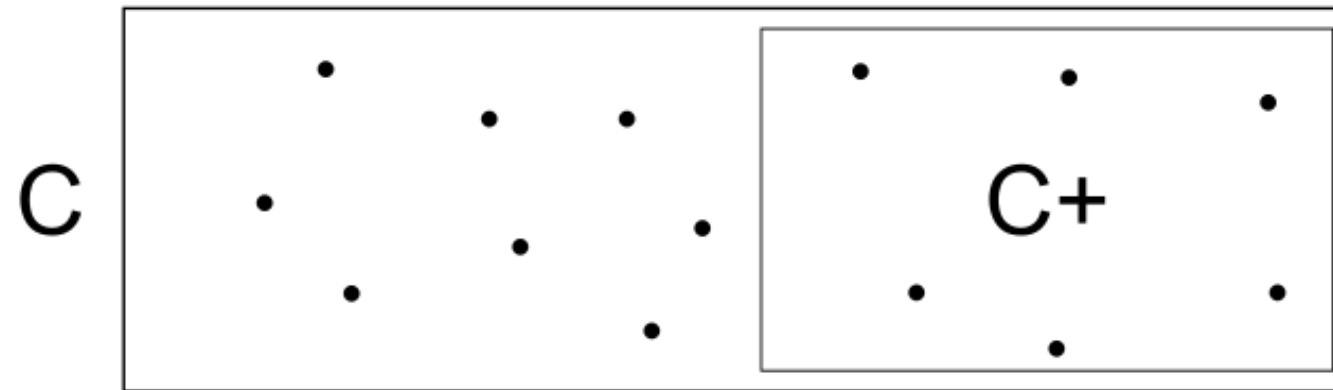


$T$  = TTCA outcome,  $R$  = redistribution of endowments within  $C$

# TTCA outcome is in the core.

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

Let  $C^+ \subseteq C$  be the agents in  $C$  who *strictly prefer*  $R$  over  $T$ .

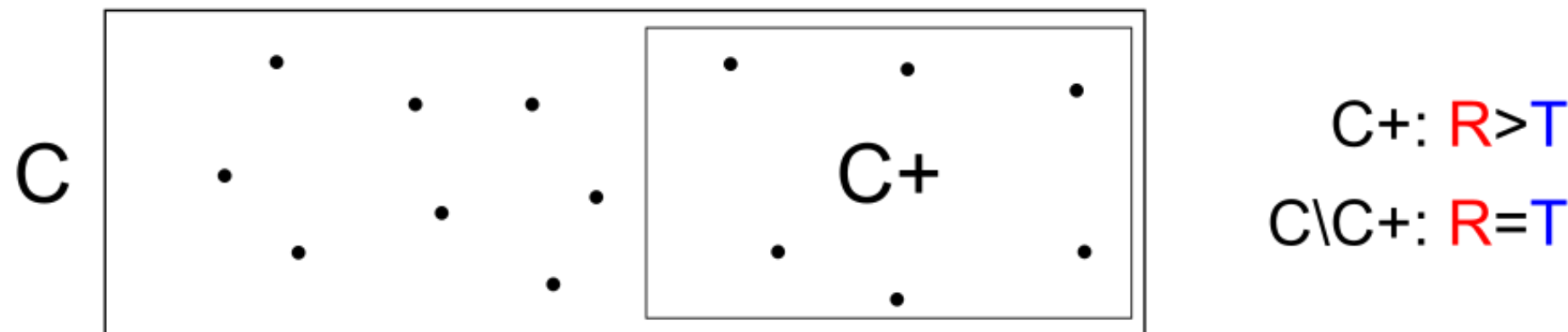


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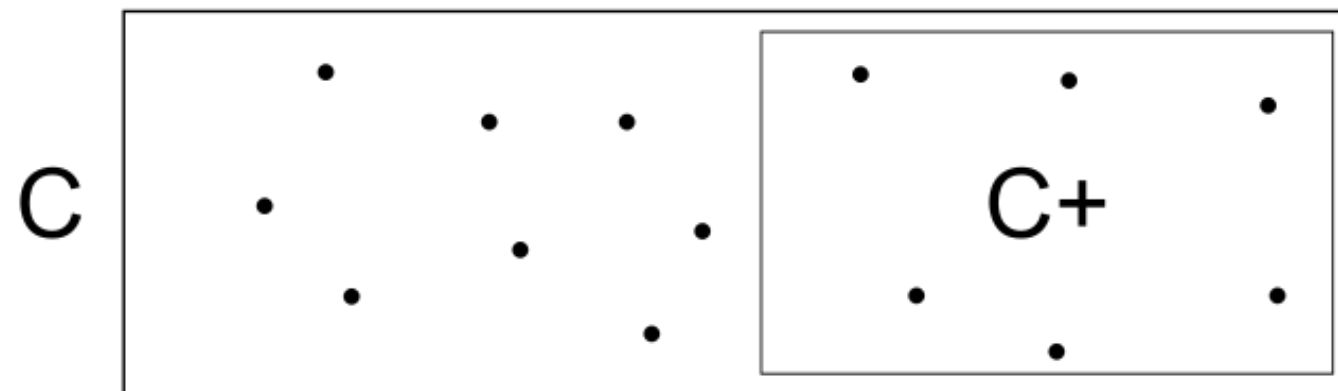


$T$  = TTCA outcome,  $R$  = redistribution of endowments within  $C$

# TTCA outcome is in the core.

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

Let  $x \in C^+$  be the agent in  $C^+$  who is the *earliest* to be eliminated under TTCA (say, in round  $r$ ).



$C^+$ :  $R > T$

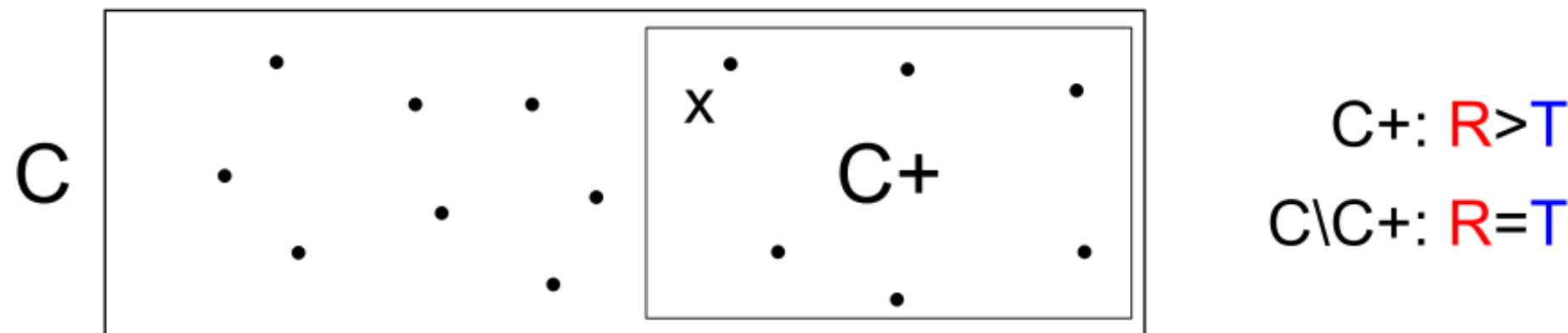
$C \setminus C^+$ :  $R = T$

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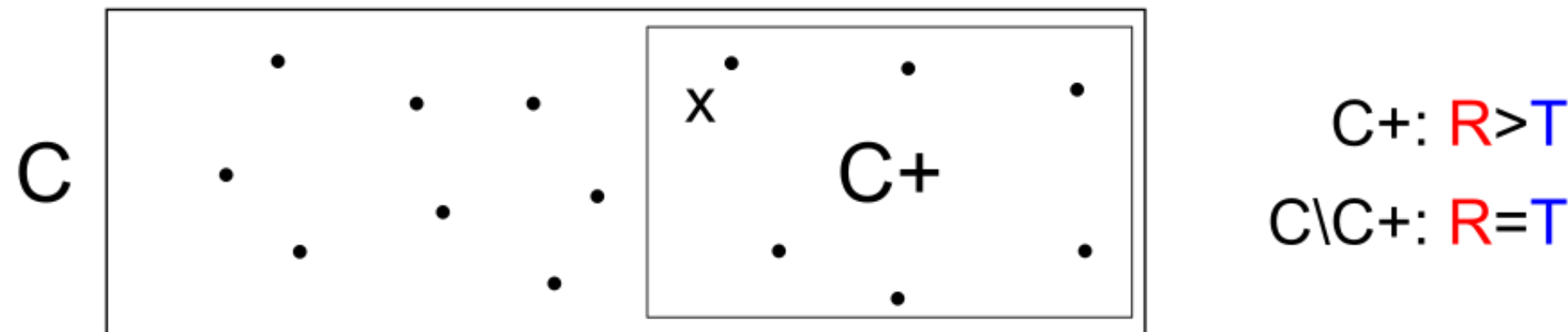


$T$  = TTCA outcome,  $R$  = redistribution of endowments within  $C$

# TTCA outcome is in the core.

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

At the time  $x$  is eliminated (round  $r$ ), all agents in  $C^+$  are still available under TTCA along with their endowed houses.



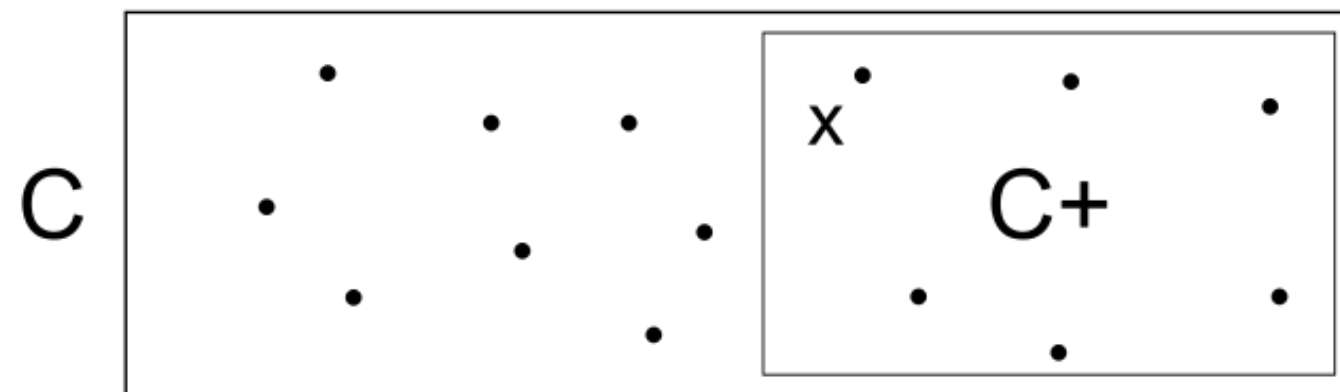
$T$  = TTCA outcome,  $R$  = redistribution of endowments within  $C$



# TTCA outcome is in the core.

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

The house  $x$  gets under  $T$  is at least as good (according to  $x$ 's preference) as the endowed house of any agent in  $C^+$ .



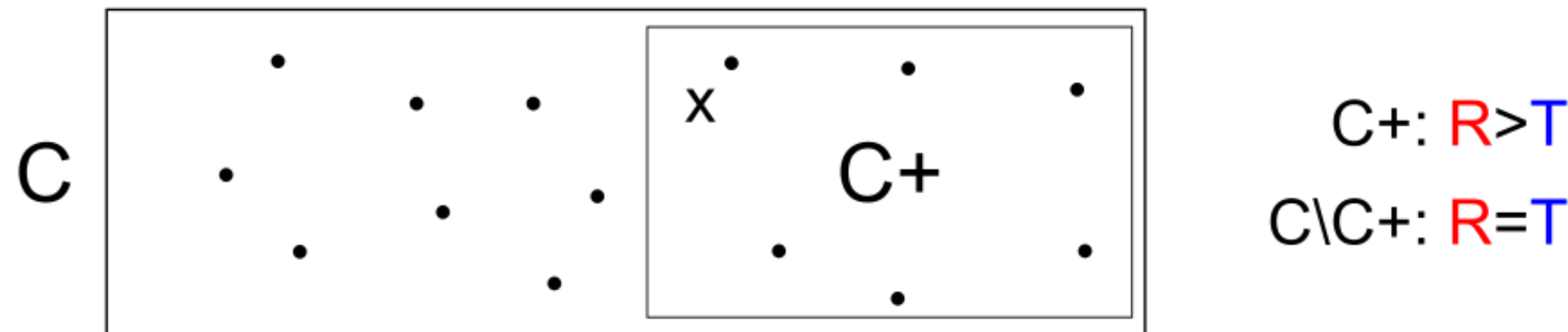
$C^+$ :  $R > T$   
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The house  $x$  gets under  $R$  is strictly better (according to  $x$ 's preference) than the endowed house of any agent in  $C^+$ .

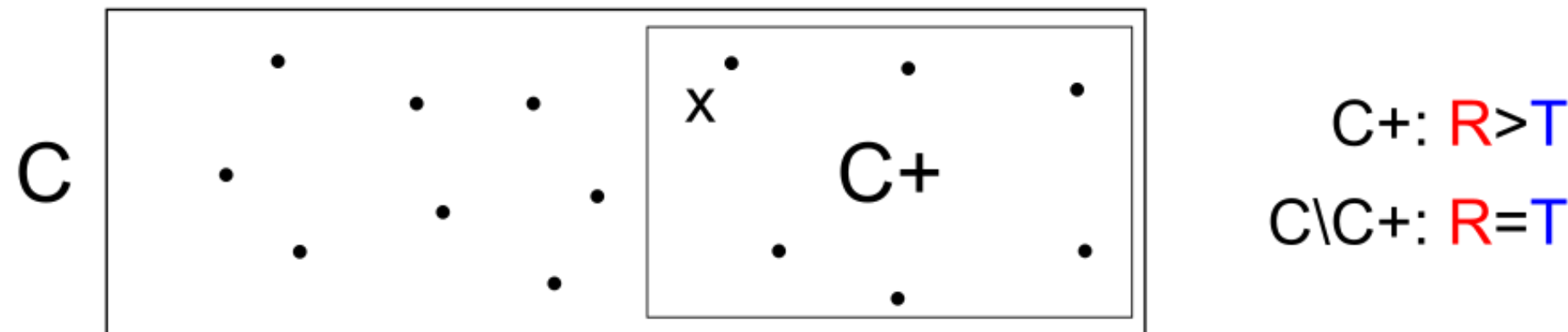


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# TTCA outcome is in the core.

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The house  $x$  gets under  $R$  must be the endowed house of some agent  $y \in C \setminus C^+$ .

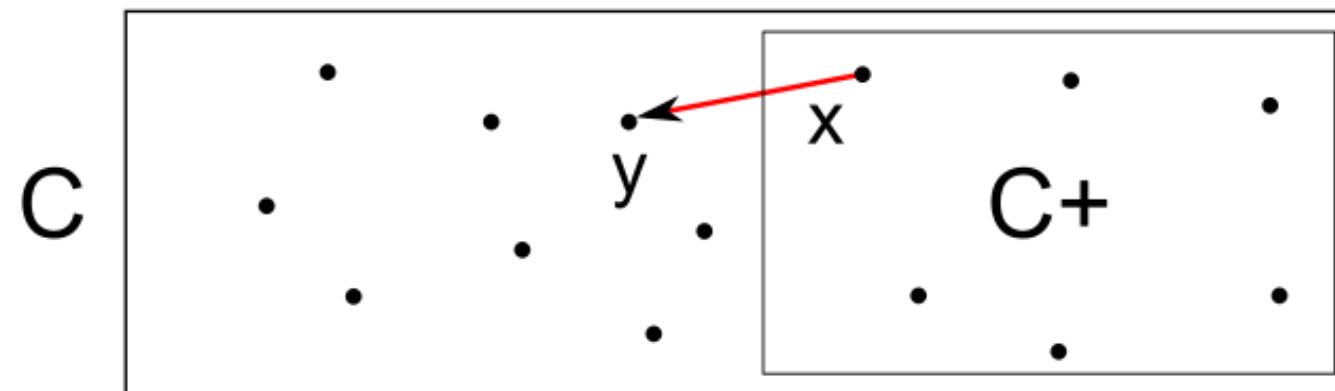


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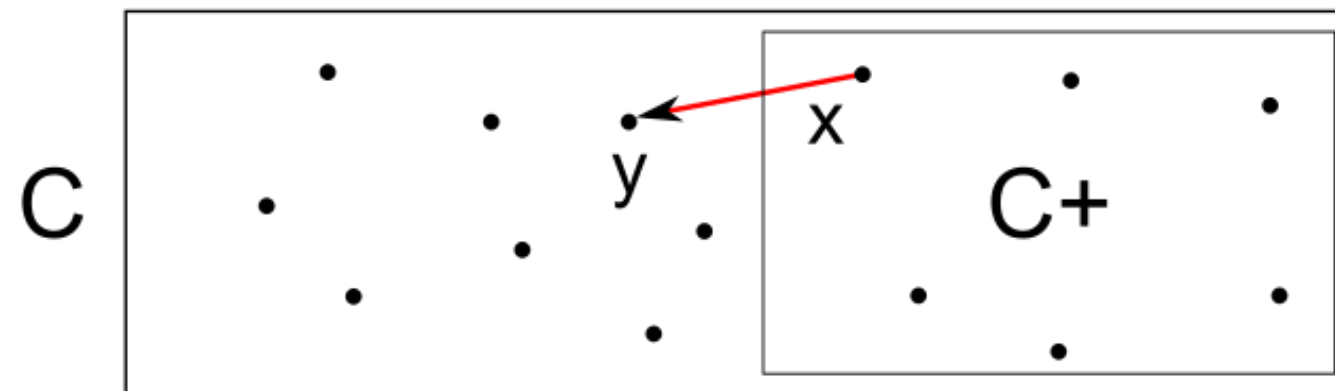
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$T$  = TTCA outcome,  $R$  = redistribution of endowments within  $C$

# TTCA outcome is in the core.

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

Then, under TTCA, agent  $y$  must have been eliminated in round  $r-1$  or earlier (i.e., strictly before  $x$ ).



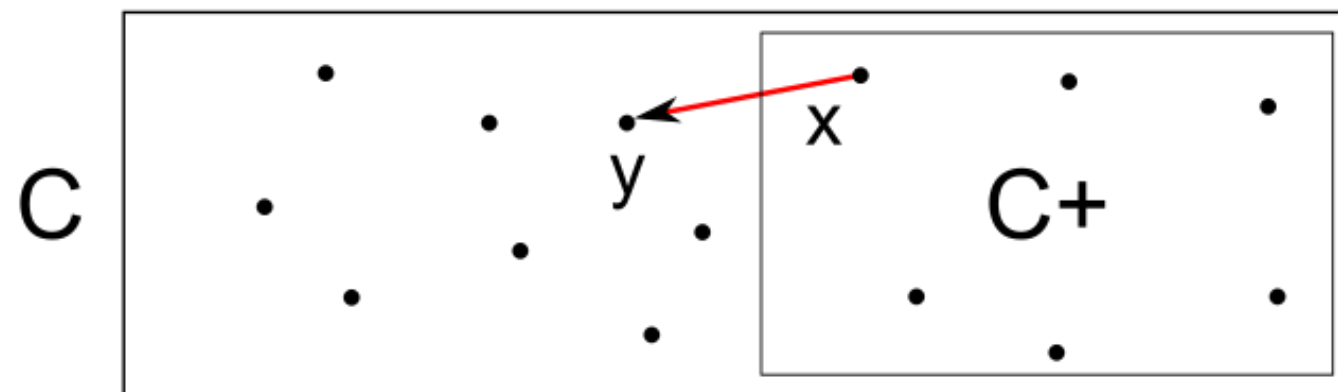
$$C^+: R > T$$
$$C \setminus C^+: R = T$$

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# TTCA outcome is in the core.

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

Agent  $y$  cannot get its own endowed house under the TTCA outcome  $T$ .



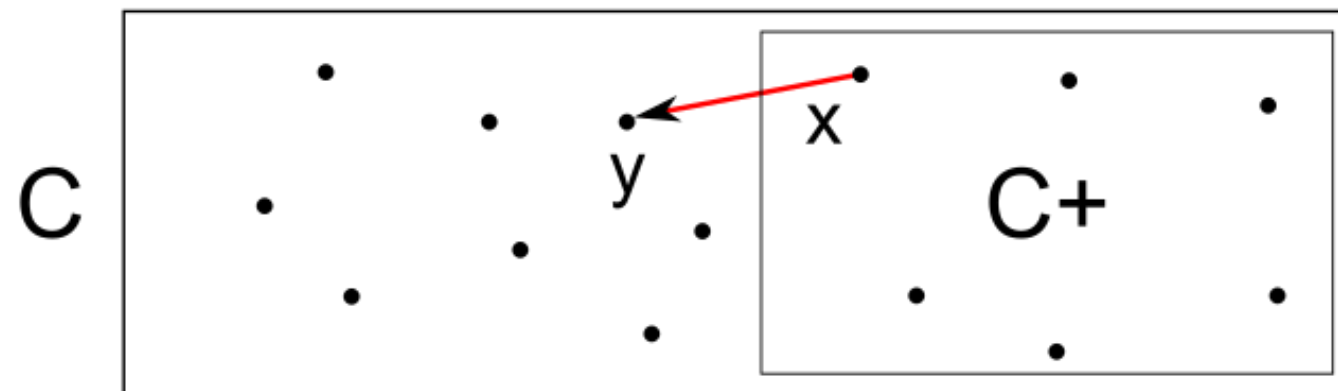
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# TTCA outcome is in the core.

[Shapley and Scarf, *JME* 1974 (attributed to David Gale)]

So, under **T** (as well as **R**),  $y$  gets the endowed house of agent  $z \in C$ .



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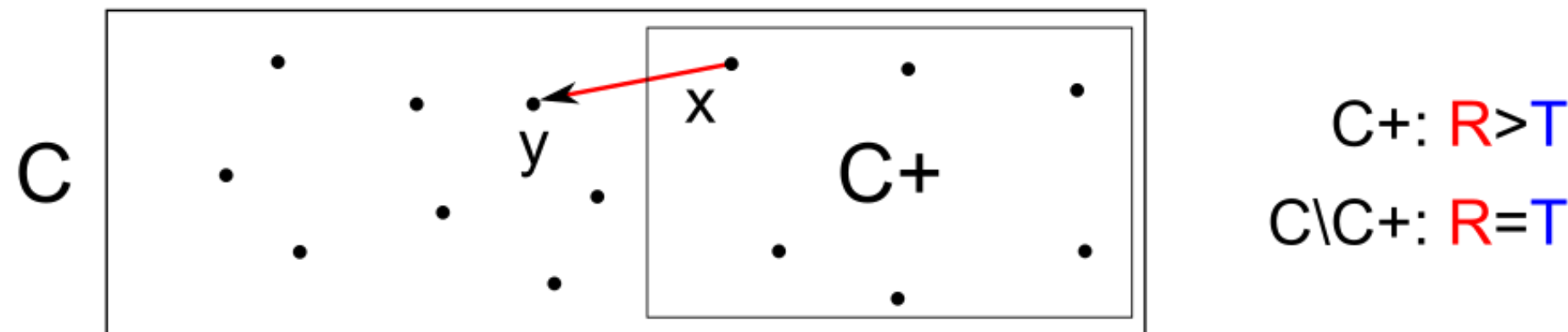
**T** = TTCA outcome, **R** = redistribution of endowments within  $C$

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Note that  $z \in C \setminus C^+$  by semantics of TTCA.



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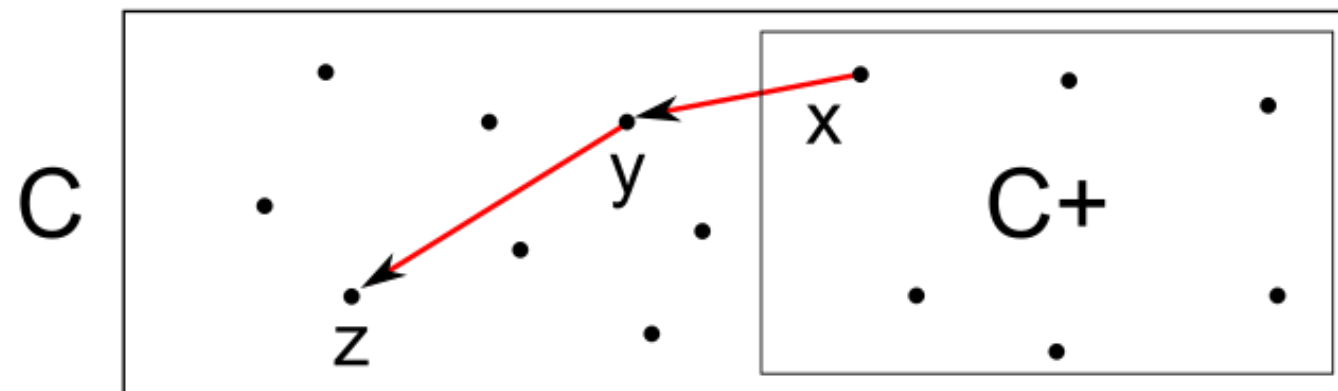


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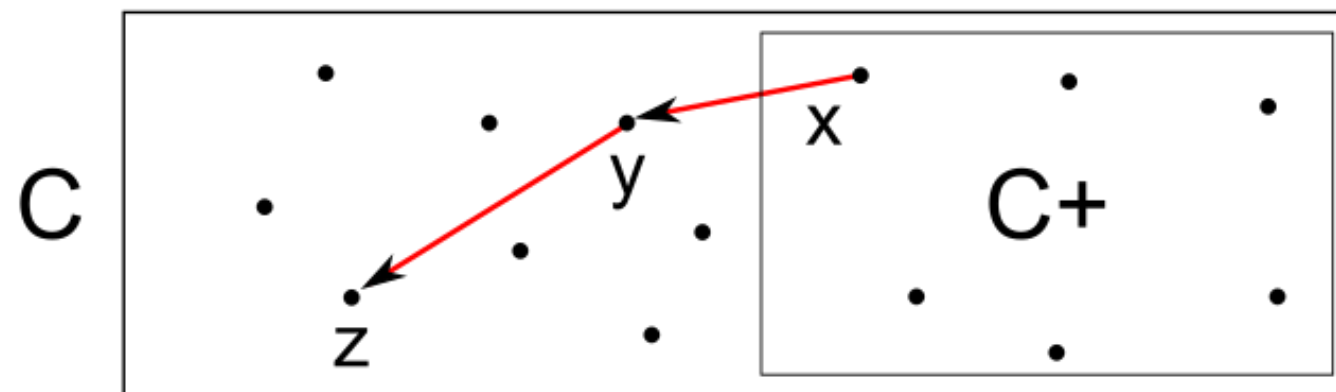
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# TTCA outcome is in the core.

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Then, under TTCA, agent  $z$  must be eliminated in the *same round* as  $y$ .



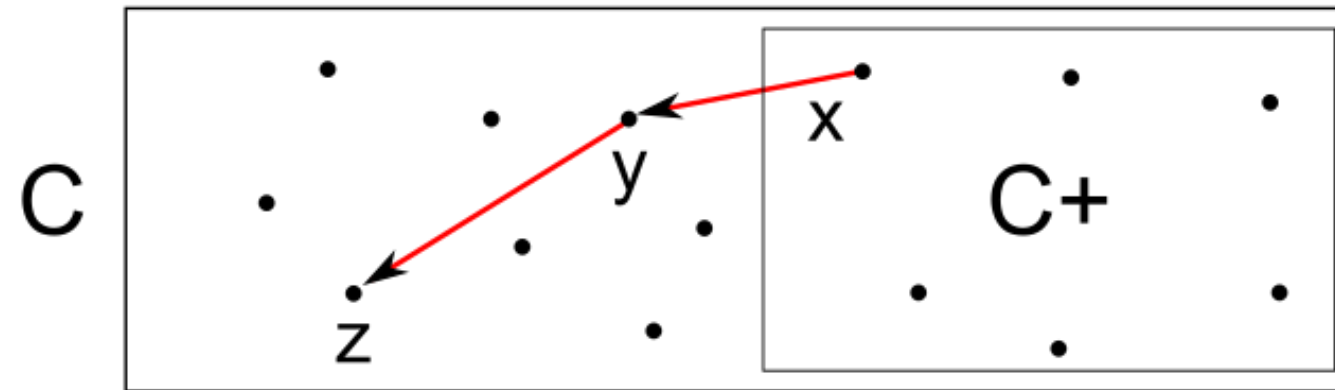
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Then, under TTCA, agent  $z$  must be eliminated in the *same round* as  $y$ .  
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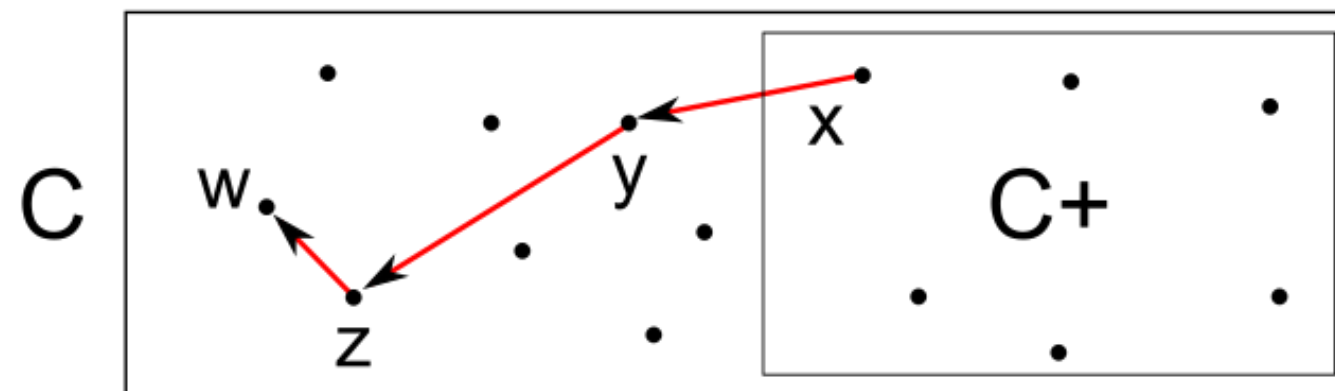
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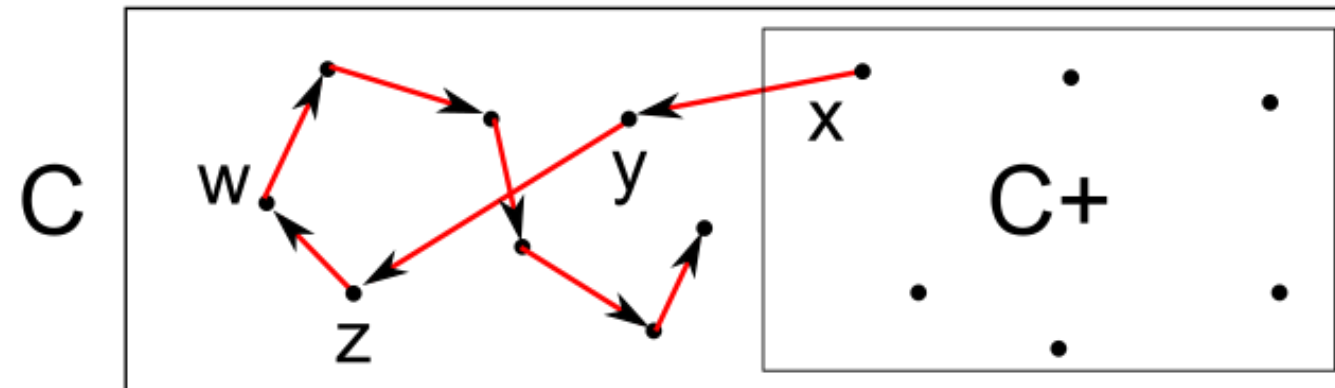
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Continuing in this manner, we get an "unending" chain within  $C \setminus C^+$ .



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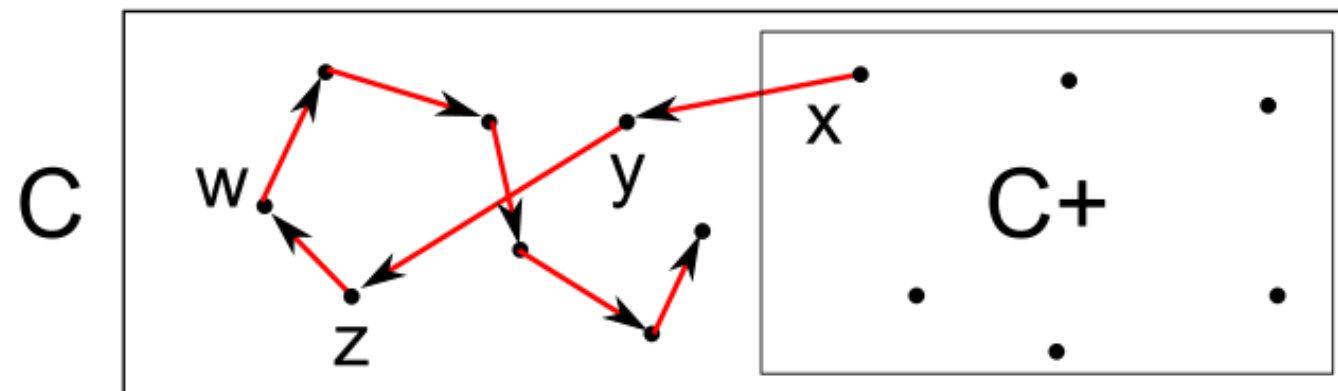
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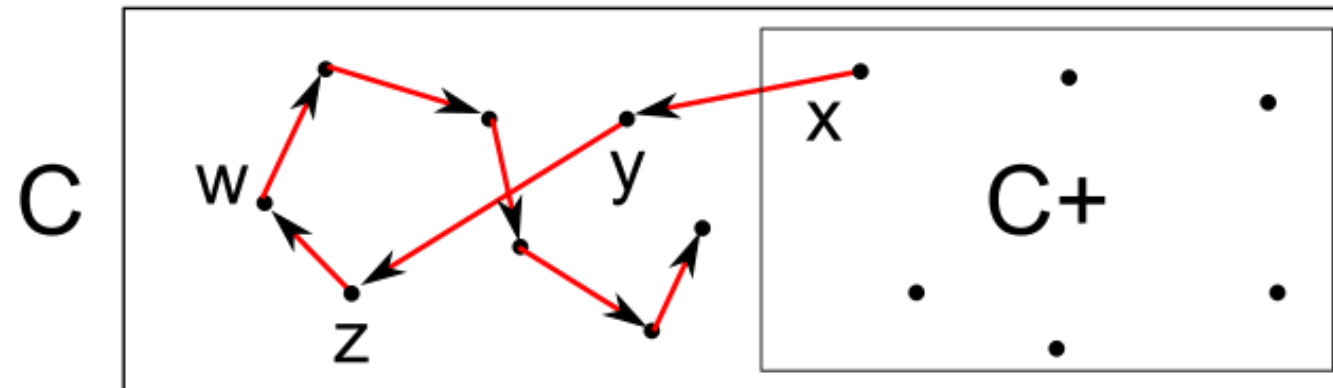
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# Awesome Properties of TTCA

- Polynomial running time
- (Group) strategyproof
- Core
- Pareto optimal (**extra credit!**)

no other allocation where all agents are weakly better off and someone is strictly better off

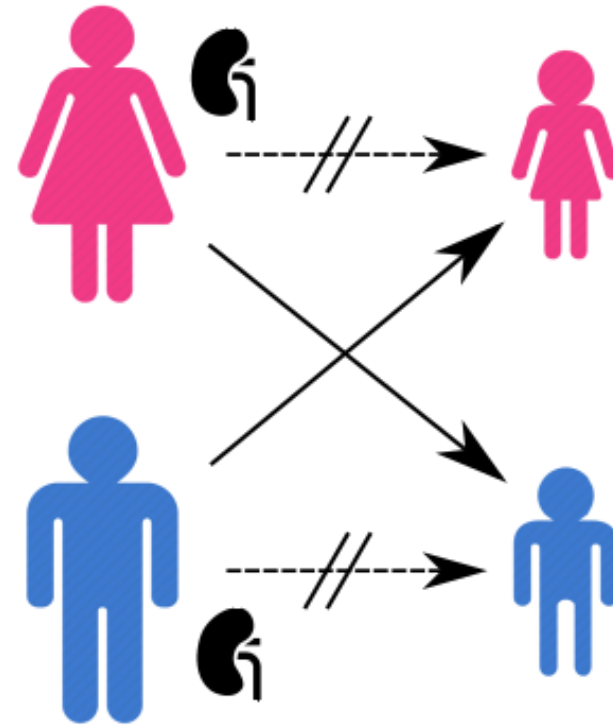


# Reminders

Project groups: Due Jan 26

Assignment 1: Due Jan 28

# Next Time: Application of TTCA



# Quiz

# Quiz

Prove that the core outcome is unique.

# References

- Housing markets and TTCA

Lloyd Shapley and Herbert Scarf

*“On Cores and Indivisibility”*

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- Truthfulness under TTCA

Alvin E. Roth

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