

# COL 202: DISCRETE MATHEMATICAL STRUCTURES

## MAJOR EXAM SOLUTIONS

# PROBLEM 1 (a)

(a) [5 points] Prove or disprove: Every graph  $G = (V, E)$  has a bipartite subgraph with at least  $|E|/2$  edges.

Proof by probabilistic argument

Assign each vertex to the "left" set w.p.  $1/2$  and "right" w.p.  $1/2$  independently of other vertices.

Fix any edge  $e = \{u, v\}$ .

Define  $X_e = \begin{cases} 1 & \text{if edge } e \text{ is crossing} \\ 0 & \text{otherwise} \end{cases}$

# PROBLEM 1 (a)

(a) [5 points] Prove or disprove: Every graph  $G = (V, E)$  has a bipartite subgraph with at least  $|E|/2$  edges.

$$\Pr(X_e = 1) = \Pr(u \text{ on left and } v \text{ on right or vice versa})$$

$$\text{disjoint events} = \Pr(u \text{ left, } v \text{ right}) + \Pr(u \text{ right, } v \text{ left})$$

$$\text{independence} = \Pr(u \text{ left}) \cdot \Pr(v \text{ right}) + \Pr(u \text{ right}) \cdot \Pr(v \text{ left})$$

$$= \frac{1}{2} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{2}$$

$$\text{Define } X = \sum_{e \in E} X_e$$

Then  $\mathbb{E}[X]$  is expected number of crossing edges.

# PROBLEM 1 (a)

(a) [5 points] Prove or disprove: Every graph  $G = (V, E)$  has a bipartite subgraph with at least  $|E|/2$  edges.

By linearity of expectation:

$$\begin{aligned} \mathbb{E}[X] &= \sum_e \mathbb{E}[X_e] \\ &= \frac{|E|}{2} \end{aligned}$$

$X$  is a random variable whose expectation is  $\frac{|E|}{2}$ .

$$\Rightarrow \mathbb{P}_x \left( X \geq \frac{|E|}{2} \right) > 0 \quad \leftarrow \text{probabilistic method}$$

$\Rightarrow \exists$  a vertex partition with at least  $\frac{|E|}{2}$  crossing edges.  $\square$

## PROBLEM 1(b)

(b) [10 points] Prove or disprove: Every graph  $G = (V, E)$  where  $|V|$  is even and  $|E| > 0$  has a bipartite subgraph with strictly more than  $|E|/2$  edges.

Proof by probabilistic argument.

Let  $|V| = 2n$ .

We will divide  $V$  into two sets, say  $A$  and  $B$ , of size  $n$  each

No. of equipartitions =  $2^n C_n$ .

Fix an edge  $e = \{u, v\}$ .

# PROBLEM 1(b)

(b) [10 points] Prove or disprove: Every graph  $G = (V, E)$  where  $|V|$  is even and  $|E| > 0$  has a bipartite subgraph with strictly more than  $|E|/2$  edges.

Let us count the number of partitions in which  $e$  is crossing

(I) If  $u \in A$  and  $v \in B$

picking  $n-1$  vertices other than  $u$  in the set  $A$

The number of such partitions is  $\binom{2n-2}{n-1}$ .

(II) If  $u \in B$  and  $v \in A$

The number of such partitions is  $\binom{2n-2}{n-1}$ .

## PROBLEM 1(b)

(b) [10 points] Prove or disprove: Every graph  $G = (V, E)$  where  $|V|$  is even and  $|E| > 0$  has a bipartite subgraph with strictly more than  $|E|/2$  edges.

Define  $X_e$  as in part (a).

Suppose each equipartition is chosen uniformly at random.

$$\Pr(X_e = 1) = \frac{2 \cdot 2^{n-2} C_{n-1}}{2^n C_n} = \frac{n}{2n-1} > \frac{1}{2}.$$

Desired bipartite subgraph exists by the same argument as in part (a).



# PROBLEM 1(a) [5 pts]

- \* Mention "We will prove the statement." \_\_\_\_\_ 0.5 pts
- \* Mention proof technique. \_\_\_\_\_ 0.5 pts
- \* Mention the experiment (random partitioning) \_\_\_\_\_ 0.5 pts
- \* Correctly define indicator random variables  
and their sum \_\_\_\_\_ 1 pt
- \* Correctly compute expected values \_\_\_\_\_ 1 pt
- \* Apply probabilistic method to finish the proof \_\_\_\_\_ 1.5 pts



# PROBLEM 1(b) [10 pts]

- \* Mention "We will prove the statement." \_\_\_\_\_ 1 pt
- \* Mention proof technique. \_\_\_\_\_ 1 pt
- \* Mention the experiment (random partitioning) \_\_\_\_\_ 1 pt
- \* Correctly define indicator random variables  
and their sum \_\_\_\_\_ 1 pt
- \* Correctly compute expected values \_\_\_\_\_ 4 pts  
(should be strictly more than  $|E|/2$ )
- \* Apply probabilistic method to finish the proof \_\_\_\_\_ 2 pts

# PROBLEM 2 (a)

## Problem 2 [6+4+5=15 points]

For any  $n \in \mathbb{N}$ , let  $[n]$  denote the set  $\{1, 2, \dots, n\}$ . We will assume that  $n \geq 3$ .

A permutation  $\sigma$  of  $[n]$  is said to be *concave* if, for every  $i \in \{2, 3, \dots, n-1\}$ ,  $\sigma(i) \geq \frac{\sigma(i-1) + \sigma(i+1)}{2}$ . For example, when  $n = 4$ , the permutation  $(1, 2, 3, 4)$  is concave but the permutation  $(4, 1, 3, 2)$  is not.

A permutation  $\sigma$  of  $[n]$  is said to be *bitonic* if there exists some  $i \in [n]$  such that

- for all  $j \in [n-1]$  such that  $j < i$ ,  $\sigma(j) < \sigma(j+1)$ , and
- for all  $k \in [n-1]$  such that  $k \geq i$ ,  $\sigma(k) > \sigma(k+1)$ .

For example, when  $n = 4$ , the permutation  $(1, 2, 3, 4)$  is bitonic but the permutation  $(4, 1, 3, 2)$  is not.

## PROBLEM 2 (a)

(a) [6 points] Prove or disprove: Every concave permutation is bitonic.

Proof by contradiction.

Let  $\sigma$  be any concave permutation of  $[n]$ .

Let  $i^* \in [n]$  be such that  $\sigma(i^*) = n$ .

Suppose, for contradiction, that  $\sigma$  is not bitonic. Then,

(i) either  $\exists j < i^*$  such that  $\sigma(j) > \sigma(j+1)$

(ii) or  $\exists k \geq i^*$  such that  $\sigma(k) < \sigma(k+1)$ .

## PROBLEM 2 (a)

(a) [6 points] Prove or disprove: Every concave permutation is bitonic.

(i)  $\exists j < i^*$  such that  $\sigma(j) > \sigma(j+1)$

let  $j^*$  be the closest index to  $i^*$  that satisfies case (i).

Observe that  $j^* \neq i^* - 1$ ; thus  $j^* < i^* - 1$ .

Then,  $\sigma(j^*) > \sigma(j^* + 1)$  and  $\sigma(j^* + 1) < \sigma(j^* + 2)$ .

$\Rightarrow$  concavity violated at  $j^* + 1$ . well-defined

Contradiction!

## PROBLEM 2 (a)

(a) [6 points] Prove or disprove: Every concave permutation is bitonic.

(ii)  $\exists k \geq i^*$  such that  $\sigma(k) < \sigma(k+1)$ .

Let  $k^*$  be the index closest to  $i^*$  that satisfies case (ii).

Then,  $k^* \neq i^*$ , and thus  $k^* > i^*$ .

We have  $\sigma(k^*) < \sigma(k^*+1)$  and  $\sigma(k^*) < \sigma(k^*-1)$

$\Rightarrow$  concavity violated at  $k^*$ .

$\uparrow$   
well-defined

Contradiction.

Therefore,  $\sigma$  must be bitonic.

$\square$

## PROBLEM 2 (b)

(b) [4 points] Identify all concave permutations of the set [5]. No explanation is required.

1 2 3 4 5

5 4 3 2 1

1 3 4 5 2

2 5 4 3 1

1 3 5 4 2

2 4 5 3 1

1 5 4 3 2

2 3 4 5 1

## PROBLEM 2 (c)

(c) [5 points] How many bitonic permutations of  $[n]$  are there? Explain your reasoning.

There are  $2^{n-1}$  bitonic permutations

Observe:

- ① 1 must always be at one of extremes of any bitonic permutation
- ② After eliminating 1, the remaining permutation of  $\{2, 3, \dots, n\}$  is also bitonic.

Recurrence:  $f(n) = 2 f(n-1) \Rightarrow f(n) = 2^{n-1}$ .

Verify by induction using above observations.

## PROBLEM 2(a) [6 pts]

- \* Mention "we will prove the statement." \_\_\_\_\_ 1 pt
- \* Mention proof technique. \_\_\_\_\_ 1 pt
- \* Correctly derive contradiction for  
the left of the peak \_\_\_\_\_ 2 pts
- \* Correctly derive contradiction for  
the right of the peak \_\_\_\_\_ 2 pts



## PROBLEM 2(b) [4 pts]

0.5 pt for each correct answer

-0.5 pt for each incorrect answer

Minimum marks : 0 / 4 .

( even if the solution consists of more incorrect answers than correct ones )

## PROBLEM 2(c) [5 pts]

- \* Mention the correct answer \_\_\_\_\_ 1 pt
- \* Making the relevant observations \_\_\_\_\_ 1 pt
- \* Correct recurrence \_\_\_\_\_ 2 pts
- \* Verify via induction \_\_\_\_\_ 1 pt

# PROBLEM 3(a)

(a) [2 points] Prove that for any non-negative random variable  $X$ ,

$$\Pr(X \geq 1) \leq \mathbf{E}[X].$$

This result is a special case of what's called Markov's inequality:  $\Pr(X \geq k) \leq \frac{\mathbf{E}[X]}{k}$

For any  $k \geq 0$ ,

if  $\Pr(X \geq k) = p$ , then  $\mathbf{E}[X] \geq k \cdot p$ .

The desired inequality follows when  $k=1$ .



# PROBLEM 3 (b)

(b) [13 points] Given any  $n \in \mathbb{N}$ , consider a random graph  $G = (V, E)$  on  $n$  vertices in which for any pair of vertices  $u, v \in V$ , the edge  $\{u, v\}$  exists with probability  $1/2$  independently of any other pair of vertices.

An *independent set* of a graph is a subset of vertices in which no two vertices are adjacent.

Show that the probability that the largest independent set of the random graph  $G$  is larger than  $\lceil 3 \log_2 n + 1 \rceil$  is  $o(n^{-\log_2 n})$ , where  $o(\cdot)$  stands for little-o notation.

$$\text{Fix } k = \lceil 3 \log_2 n + 1 \rceil.$$

Fix any subset of vertices  $S \subseteq V$  such that  $|S| = k$

$$\begin{aligned} \Pr(S \text{ is independent}) &= \Pr(\text{no edge between any of the } \binom{k}{2} \\ &\quad \text{pairs of vertices in } S) \\ &= \left(\frac{1}{2}\right)^{\binom{k}{2}} \quad \text{--- } \textcircled{1} \end{aligned}$$

# PROBLEM 3 (b)

(b) [13 points] Given any  $n \in \mathbb{N}$ , consider a random graph  $G = (V, E)$  on  $n$  vertices in which for any pair of vertices  $u, v \in V$ , the edge  $\{u, v\}$  exists with probability  $1/2$  independently of any other pair of vertices.

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Let  $S_1, S_2, \dots, S_{n C_k}$  be all  $k$ -sized subsets of vertices.

Let  $X_i = \begin{cases} 1 & \text{if } S_i \text{ is independent} \\ 0 & \end{cases}$

Let  $X = \sum_{i=1}^{n C_k} X_i$

Then  $\mathbb{E}[X] = \sum_i \mathbb{E}[X_i] = \sum_i \Pr(X_i=1) \stackrel{\text{by linearity of expectation}}{=} \sum_i \Pr(X_i=1) \stackrel{\text{being } \textcircled{1}}{=} n C_k \cdot \left(\frac{1}{2}\right)^{k C_2}$

# PROBLEM 3(b)

(b) [13 points] Given any  $n \in \mathbb{N}$ , consider a random graph  $G = (V, E)$  on  $n$  vertices in which for any pair of vertices  $u, v \in V$ , the edge  $\{u, v\}$  exists with probability  $1/2$  independently of any other pair of vertices.

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Show that the probability that the largest independent set of the random graph  $G$  is larger than  $\lceil 3 \log_2 n + 1 \rceil$  is  $o(n^{-\log_2 n})$ , where  $o(\cdot)$  stands for little-o notation.

$$\mathbb{E}[X] = \binom{n}{k} \cdot \left(\frac{1}{2}\right)^{kC_2}$$

$$\leq n^k \cdot \left(\left(\frac{1}{2}\right)^{(k-1)/2}\right)^k$$

$$\leq \left[ n \cdot \left(\frac{1}{2}\right)^{\frac{3}{2} \log_2 n} \right]^k$$

$$= \left[ n \cdot n^{-3/2} \right]^k$$

$$\text{since } \binom{n}{k} \leq n^k$$

$$\text{since } k \geq 3 \log_2 n$$

$$\leq n^{-k/2}$$

————— (2)

# PROBLEM 3(b)

(b) [13 points] Given any  $n \in \mathbb{N}$ , consider a random graph  $G = (V, E)$  on  $n$  vertices in which for any pair of vertices  $u, v \in V$ , the edge  $\{u, v\}$  exists with probability  $1/2$  independently of any other pair of vertices.

An *independent set* of a graph is a subset of vertices in which no two vertices are adjacent.

Show that the probability that the largest independent set of the random graph  $G$  is larger than  $\lceil 3 \log_2 n + 1 \rceil$  is  $o(n^{-\log_2 n})$ , where  $o(\cdot)$  stands for little-o notation.

From part (a), we have  $\Pr(X \geq 1) \leq \mathbb{E}[X]$ .

$$\Rightarrow \Pr(X \geq 1) \leq n^{-k/2} \quad (\text{from } \textcircled{2})$$

$$= o(n^{-\log_2 n}) \quad \text{---} \quad \textcircled{3}$$

# PROBLEM 3(b)

(b) [13 points] Given any  $n \in \mathbb{N}$ , consider a *random graph*  $G = (V, E)$  on  $n$  vertices in which for any pair of vertices  $u, v \in V$ , the edge  $\{u, v\}$  exists with probability  $1/2$  independently of any other pair of vertices.

An *independent set* of a graph is a subset of vertices in which no two vertices are adjacent.

Show that the probability that the largest independent set of the random graph  $G$  is larger than  $\lceil 3 \log_2 n + 1 \rceil$  is  $o(n^{-\log_2 n})$ , where  $o(\cdot)$  stands for little-o notation.

$$\begin{aligned}
 & \Pr(\text{size of largest independent set} \geq k) \\
 &= \Pr(\text{there exists an independent set of size} \geq k) \quad \leftarrow \text{same events} \\
 &\leq \Pr(\text{" " " " " } = k) \\
 &= \Pr(X \geq 1) \quad \leftarrow (\text{being } A \subseteq B \Rightarrow \Pr(A) \leq \Pr(B)) \\
 &= o(n^{-\log_2 n}) \quad \text{from } \textcircled{3} \quad \text{as desired.}
 \end{aligned}$$



## PROBLEM 3 (a) [2 pts]

- \* Proving the inequality for all  $k \geq 0$  \_\_\_\_\_ 1.5 pt
- \* Substituting  $k=1$  \_\_\_\_\_ 0.5 pt

## PROBLEM 3 (b) [13 pts]

\* Computing expected value of indicator variables — 3 pts

\* Deriving  $O\left(\frac{1}{n \lg_2 n}\right)$  bound on  $\Pr(X \geq 1)$  — 8 pts

\* Finishing the proof by observing that the bound on  $\Pr(X \geq 1)$  gives a bound on the desired probability — 2 pts

## PROBLEM 4 (a)

(a) [5 points] Let  $a, b, c, d$ , and  $m$  be positive integers. Prove or disprove: If  $a \equiv b \pmod{m}$ ,  $c \equiv d \pmod{m}$ , and  $\gcd(c, m) = 1$ , then  $a \cdot c^{-1} \equiv b \cdot d^{-1} \pmod{m}$ , where  $c^{-1}$  and  $d^{-1}$  are the multiplicative inverses  $\pmod{m}$  of  $c$  and  $d$ , respectively.

Proof by using standard properties of congruence.

Observe:

$$\textcircled{1} \quad \gcd(c, m) = 1 \text{ and } c \equiv d \pmod{m} \Rightarrow \gcd(d, m) = 1.$$

$\textcircled{2}$  By  $\textcircled{1}$ ,  $c^{-1}$  and  $d^{-1}$  are well-defined.

$$\text{Then, } c \cdot (ac^{-1} - bd^{-1}) \pmod{m}$$

$$\equiv acc^{-1} - bcd^{-1} \pmod{m}$$

$$\equiv a \cdot 1 - b \cdot 1 \pmod{m}$$

$$\equiv a - b \pmod{m}$$

$$\left[ \text{Note: } c \equiv d \pmod{m} \text{ and } dd^{-1} \equiv 1 \pmod{m} \right] \\ \Rightarrow cd^{-1} \equiv 1 \pmod{m}$$

## PROBLEM 4 (a)

(a) [5 points] Let  $a, b, c, d$ , and  $m$  be positive integers. Prove or disprove: If  $a \equiv b \pmod{m}$ ,  $c \equiv d \pmod{m}$ , and  $\gcd(c, m) = 1$ , then  $a \cdot c^{-1} \equiv b \cdot d^{-1} \pmod{m}$ , where  $c^{-1}$  and  $d^{-1}$  are the multiplicative inverses  $\pmod{m}$  of  $c$  and  $d$ , respectively.

$$\text{Thus, } c \cdot (ac^{-1} - bd^{-1}) \pmod{m} \equiv a - b \pmod{m} \equiv 0 \pmod{m}$$

Since  $c$  and  $m$  are relatively prime

we have  $ac^{-1} - bd^{-1} \equiv 0 \pmod{m}$  as desired.



## PROBLEM 4 (b)

(b) [5 points] Let  $a, b, c, d$ , and  $m$  be positive integers such that  $b$  and  $m$  are relatively prime. Prove or disprove: If  $b^a \equiv 1 \pmod{m}$ ,  $b^c \equiv 1 \pmod{m}$ , and  $d = \gcd(a, c)$ , then  $b^d \equiv 1 \pmod{m}$ . How does your answer change if you are not given that  $b$  and  $m$  are relatively prime?

Proof by using gcd - spe equivalence and part (a).

$$d = \gcd(a, c) \Rightarrow \exists \text{ integers } \alpha, \beta \text{ such that } d = \alpha a + \beta c.$$

Without loss of generality,  $\alpha \geq 0$  (can achieve by adding enough copies of 'a')

Thus, we must have that  $\beta \leq 0$ .

$$b^a \equiv 1 \pmod{m} \Rightarrow b^{\alpha a} \equiv 1 \pmod{m} \longrightarrow \textcircled{1}$$

$$b^c \equiv 1 \pmod{m} \Rightarrow b^{-\beta c} \equiv 1 \pmod{m} \longrightarrow \textcircled{2}$$

## PROBLEM 4 (b)

(b) [5 points] Let  $a, b, c, d$ , and  $m$  be positive integers such that  $b$  and  $m$  are relatively prime. Prove or disprove: If  $b^a \equiv 1 \pmod{m}$ ,  $b^c \equiv 1 \pmod{m}$ , and  $d = \gcd(a, c)$ , then  $b^d \equiv 1 \pmod{m}$ . How does your answer change if you are not given that  $b$  and  $m$  are relatively prime?

Observe that  $\gcd(b^{-\beta c}, m) = 1$ . This is because

$$b^{-\beta c} \equiv 1 \pmod{m} \quad \text{and} \quad \gcd(1, m) = 1. \quad \left( \begin{array}{l} \text{Do not need to assume} \\ \text{that } b, m \text{ are rel. prime} \end{array} \right)$$

By applying part (a), we can divide ① by ② to get

$$b^{\alpha a + \beta c} \equiv 1 \pmod{m}$$

$$\text{or } b^d \equiv 1 \pmod{m} \quad \text{as desired.} \quad \square$$

# PROBLEM 4 (c)

(c) [5 points] Let  $b$ ,  $p$ , and  $n$  be positive integers. Prove or disprove: If  $p$  is a prime such that  $p \mid (b^n - 1)$ , then:

- either  $p \mid (b^d - 1)$  for some proper divisor  $d$  of  $n$  (a proper divisor of  $n$  is any positive divisor of  $n$  excluding  $n$  itself),
- or  $p \equiv 1 \pmod{n}$ .

Proof by using Euler's theorem and part (b).

$p$  is prime  $\Rightarrow b^{p-1} \equiv 1 \pmod{p}$  by Euler's thm since  $\phi(p) = p-1$ .

Given  $b^n \equiv 1 \pmod{p}$ .

Let  $d = \gcd(n, p-1)$ .

By part (b),  $b^d \equiv 1 \pmod{p}$

# PROBLEM 4 (c)

(c) [5 points] Let  $b$ ,  $p$ , and  $n$  be positive integers. Prove or disprove: If  $p$  is a prime such that  $p \mid (b^n - 1)$ , then:

- either  $p \mid (b^d - 1)$  for some proper divisor  $d$  of  $n$  (a proper divisor of  $n$  is any positive divisor of  $n$  excluding  $n$  itself),
- or  $p \equiv 1 \pmod{n}$ .

If  $d = n$ , then  $\gcd(n, p-1) = n \implies n \mid p-1$   
 $\implies p \equiv 1 \pmod{n}$ .

If  $d < n$ ,  $p \mid b^d - 1$  for some divisor  $d < n$  of  $n$   
 $\searrow$  proper divisor.





## PROBLEM 4 (a) [5 pts]

- \* Mentioning "We will prove the statement" ——— 0.5 pt
- \* Observing that  $c^{-1}$  and  $d^{-1}$  are well-defined ——— 1 pt
- \* Observing that  $c \cdot (ac^{-1} - bd^{-1}) \equiv 0 \pmod{m}$  ——— 2.5 pts
- \* Using relative primality of  $c$  and  $m$  ——— 1 pt  
to finish the proof

## PROBLEM 4 (b) [5 pts]

- \* Mentioning "We will prove the statement" ——— 0.5 pt
- \* Invoking gcd-spc equivalence and observing that  $\alpha \geq 0$  and  $\beta \leq 0$  ——— 2 pts
- \* Observing that part (a) can be used to divide the congruences in ① and ② ——— 1.5 pts
- \* Stating that relative primality of  $b$  and  $m$  is not needed. ——— 1 pts

## PROBLEM 4 (c) [5 pts]

- \* Mentioning "We will prove the statement" ——— 0.5 pt
- \* Using Euler's theorem ————— 1 pt
- \* Using part (b) ————— 1.5 pts
- \* Case analysis for  $d=n$  and  $d < n$  ————— 2 pts