

COL 202: DISCRETE MATHEMATICAL STRUCTURES

LECTURE 24

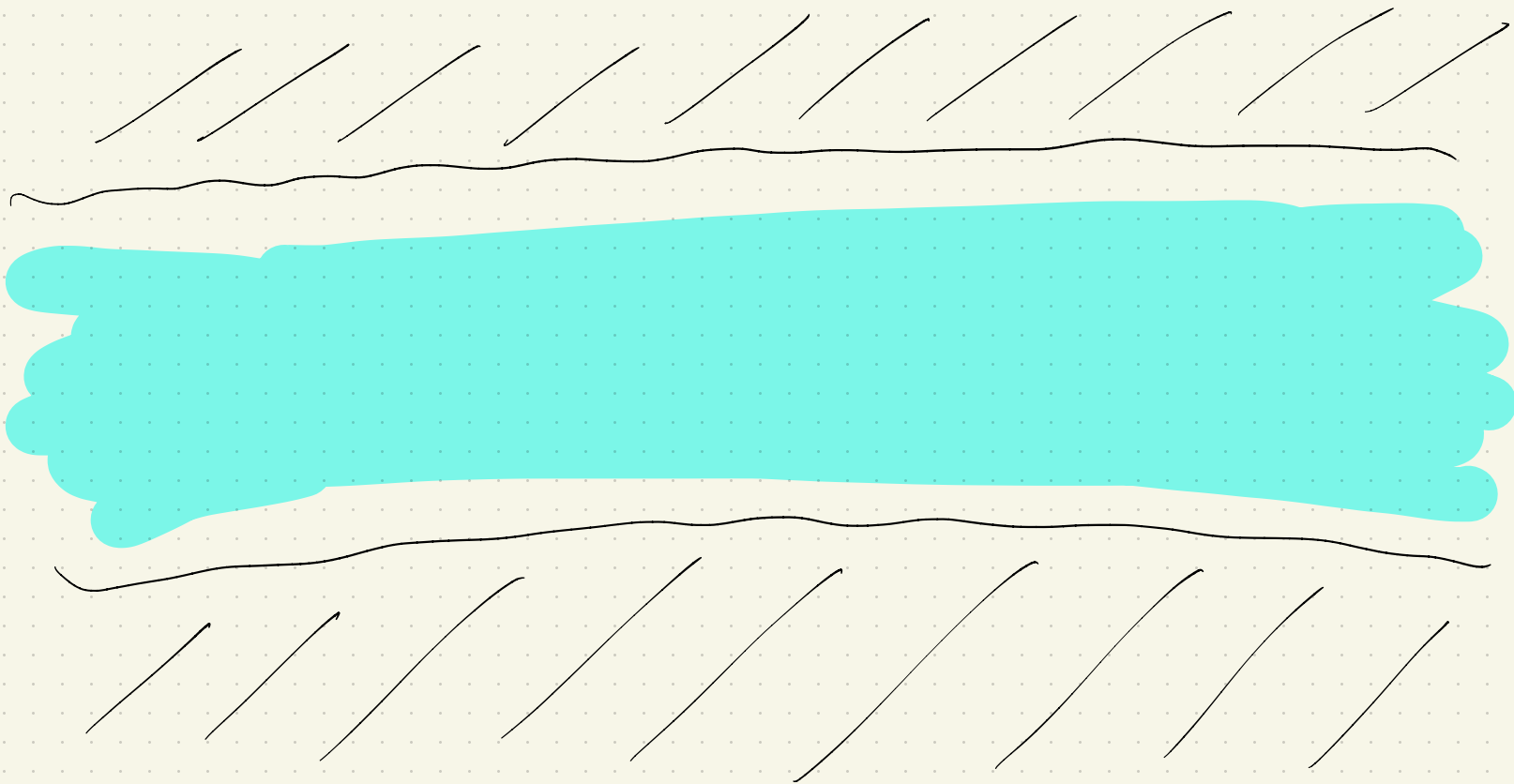
EULER TOURS

MAR 06, 2024

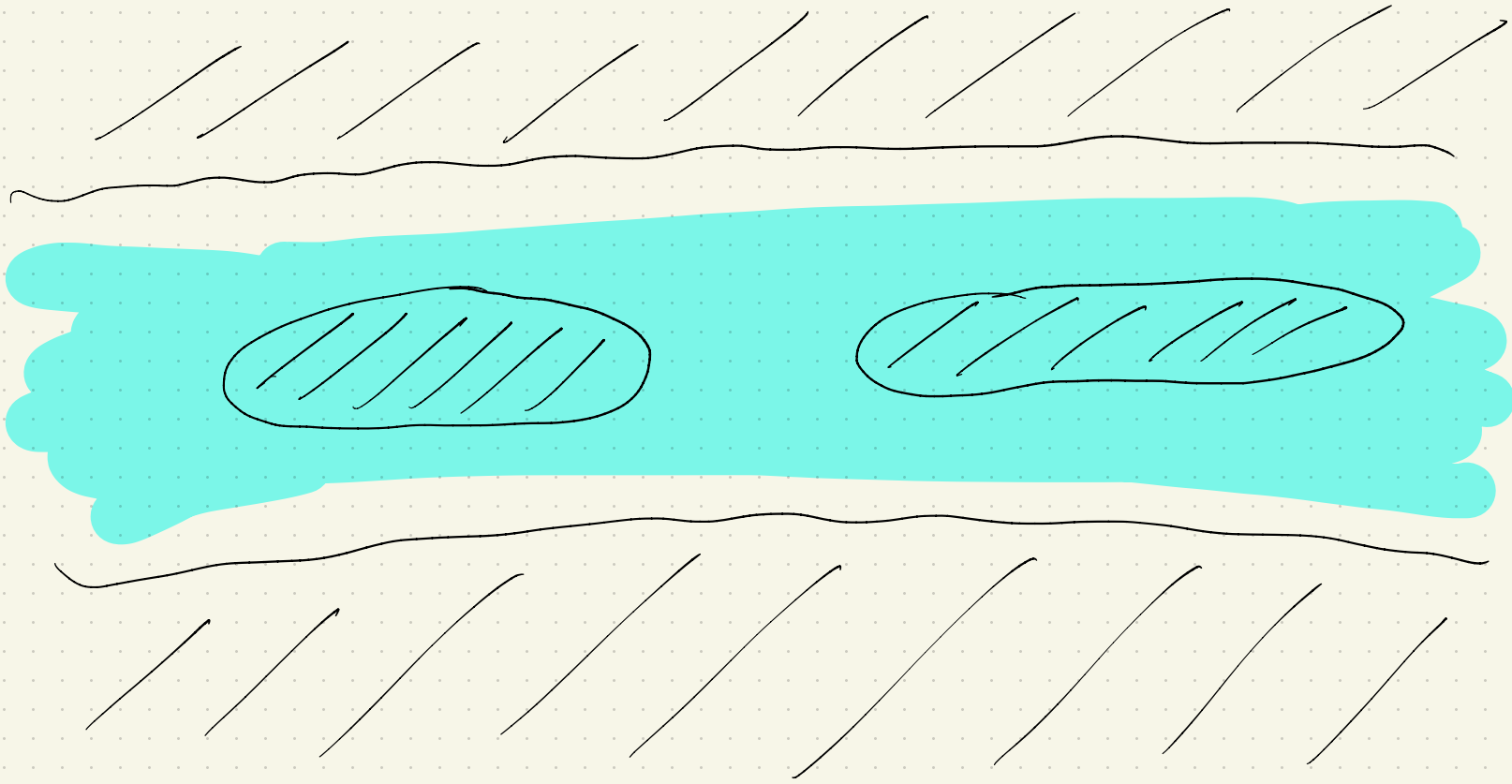
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ROHIT VAISH

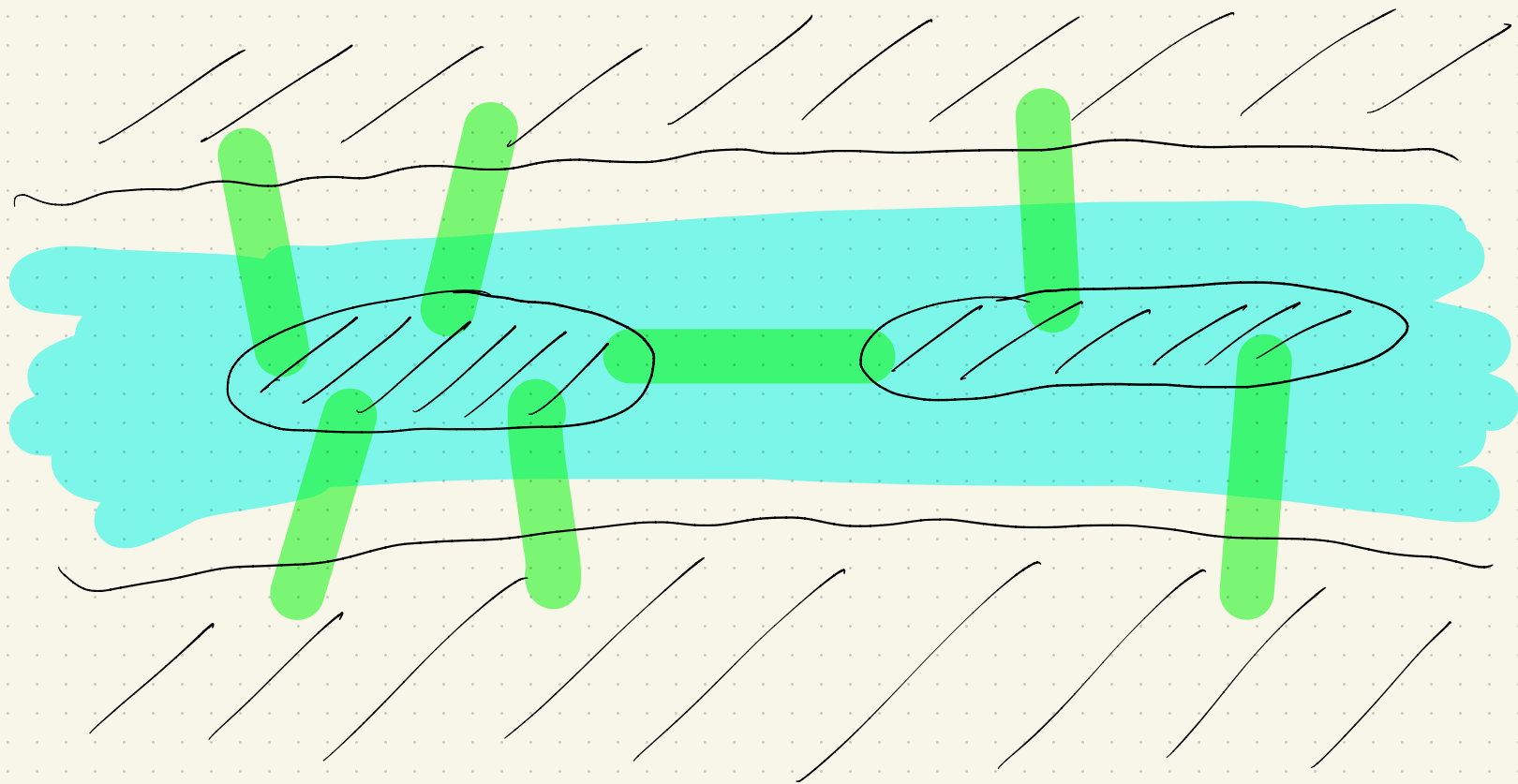
SEVEN BRIDGES OF KÖNIGSBERG



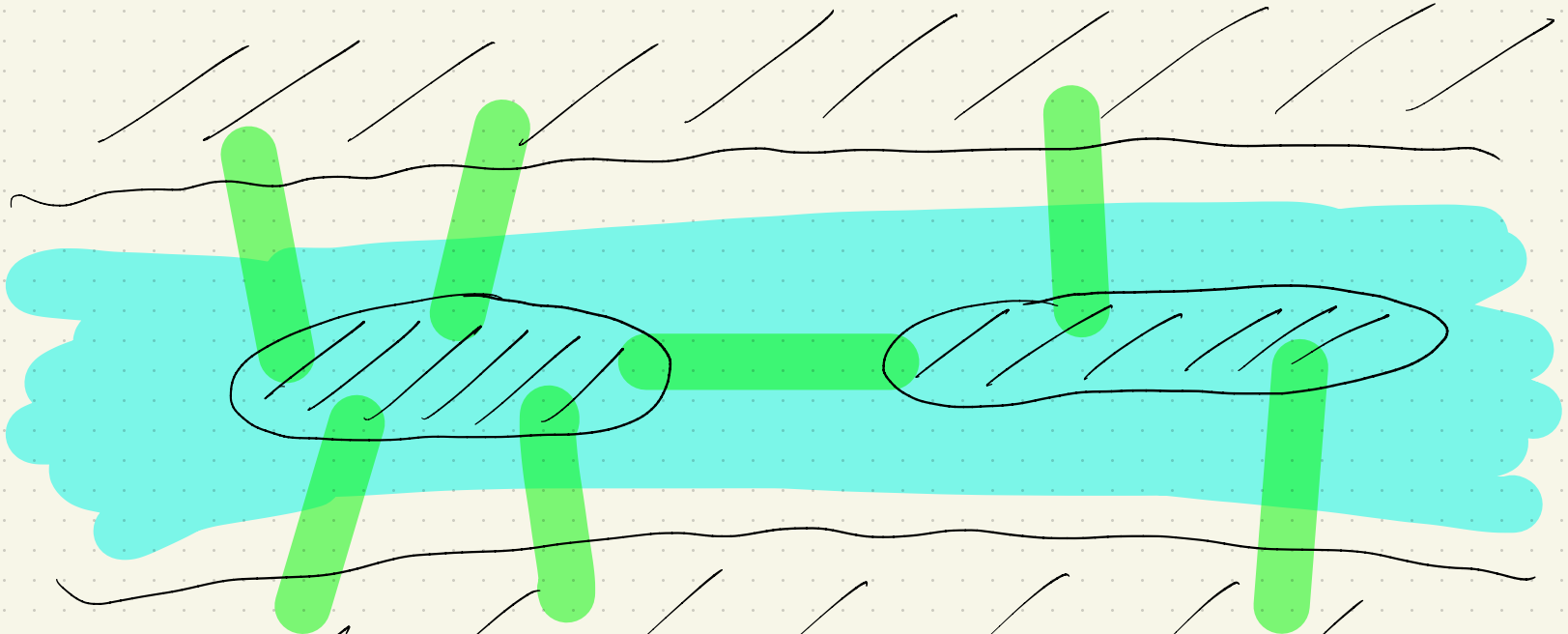
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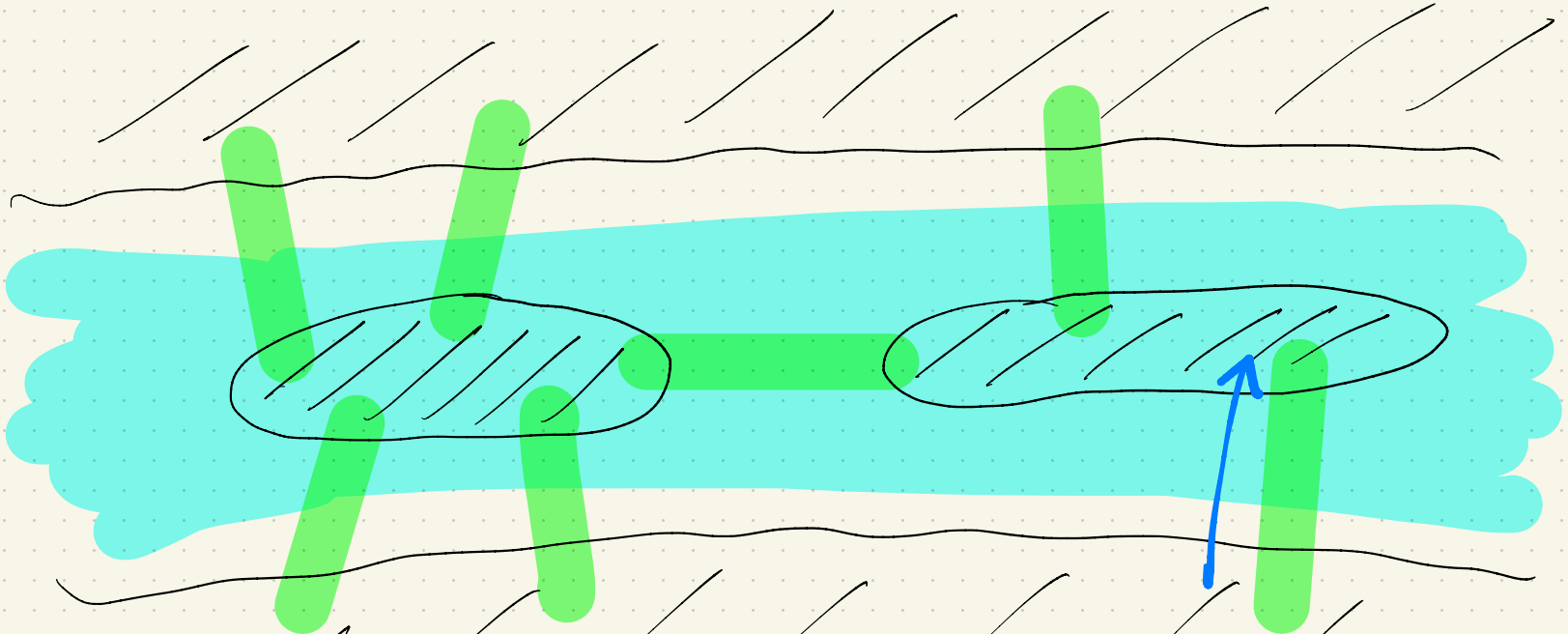
SEVEN BRIDGES OF KÖNIGSBERG



Return to starting point after
crossing each bridge exactly once.



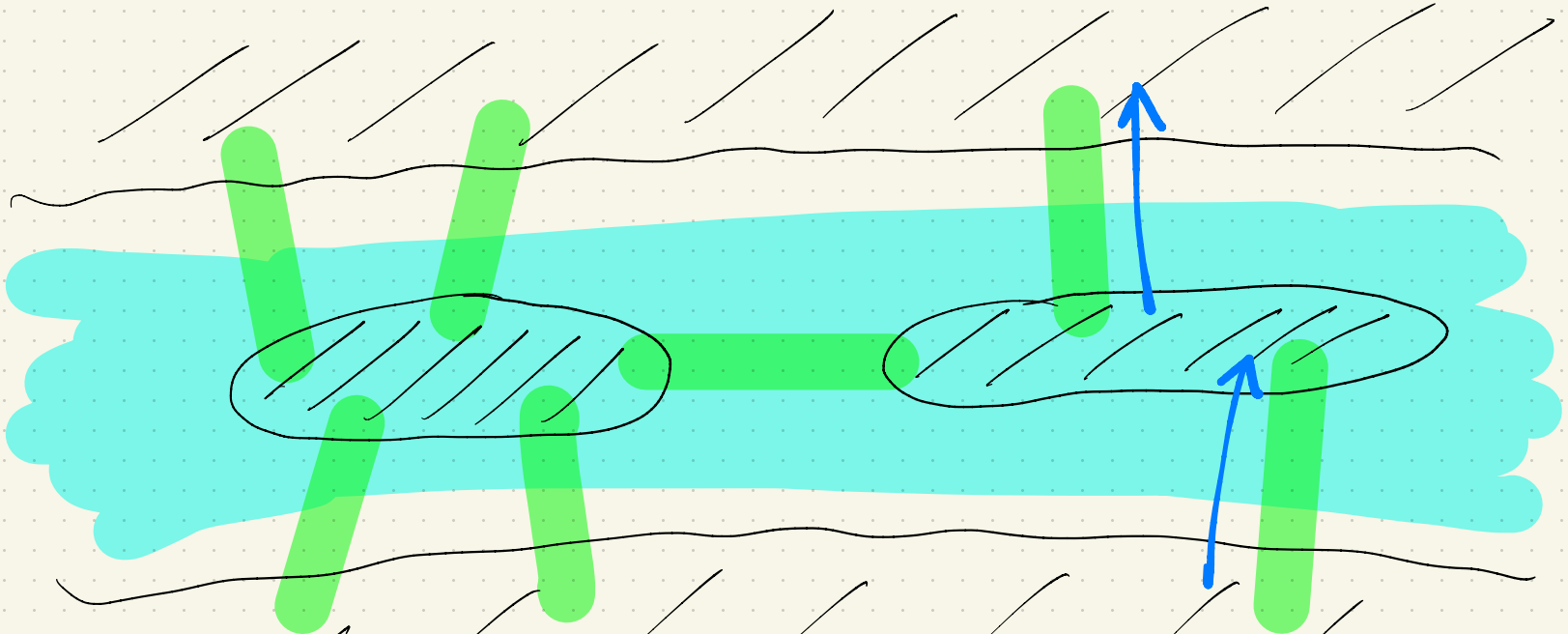
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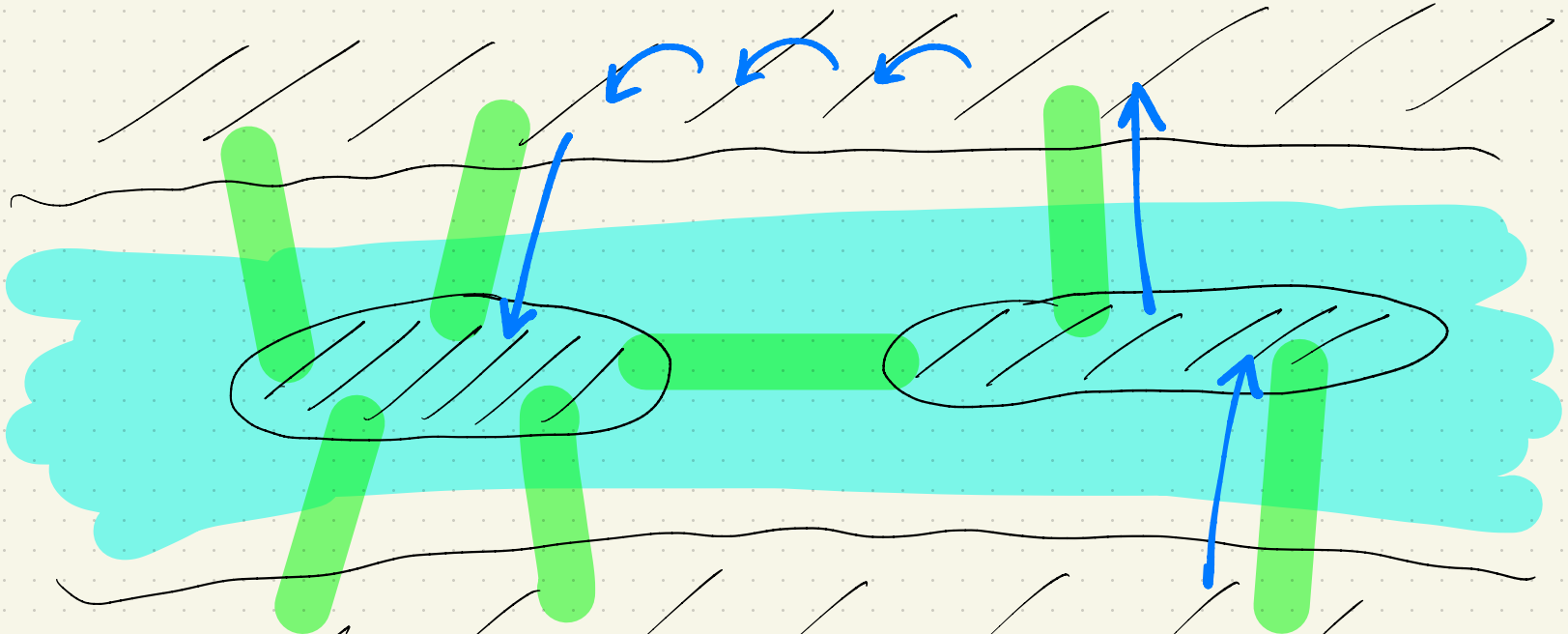
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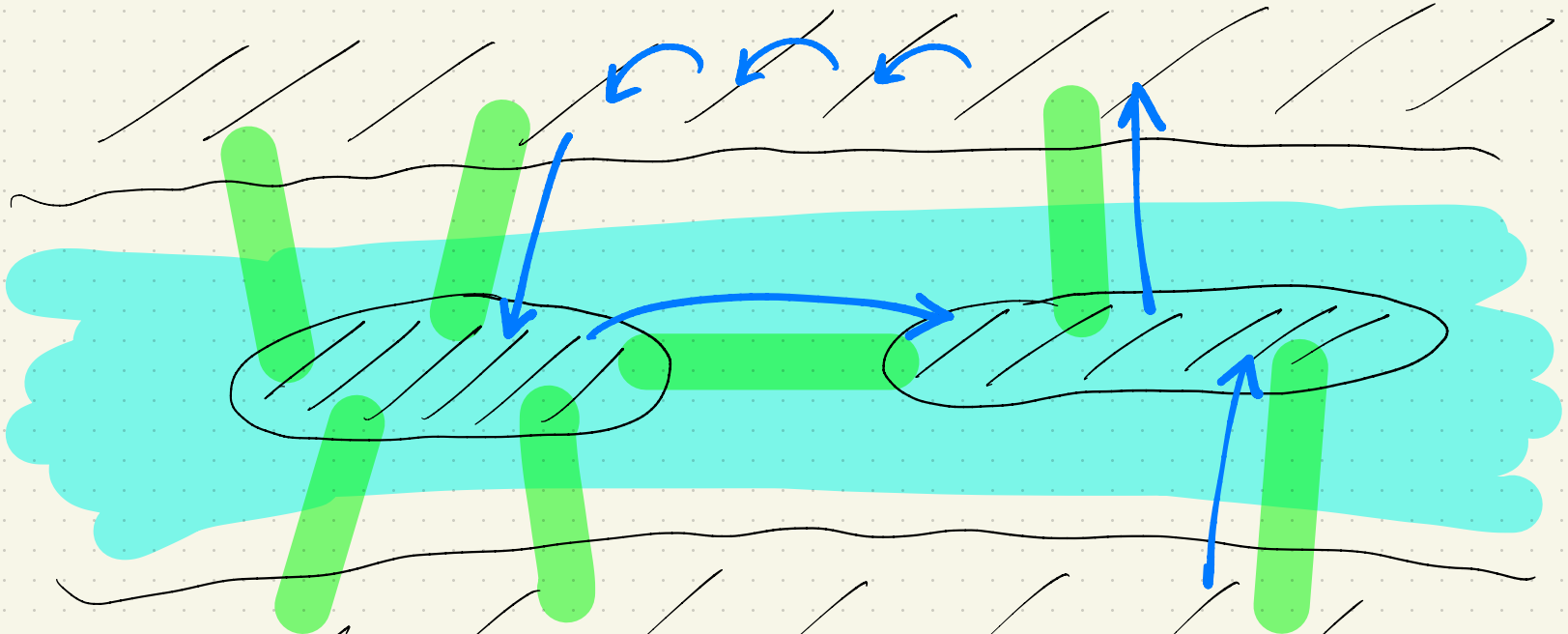
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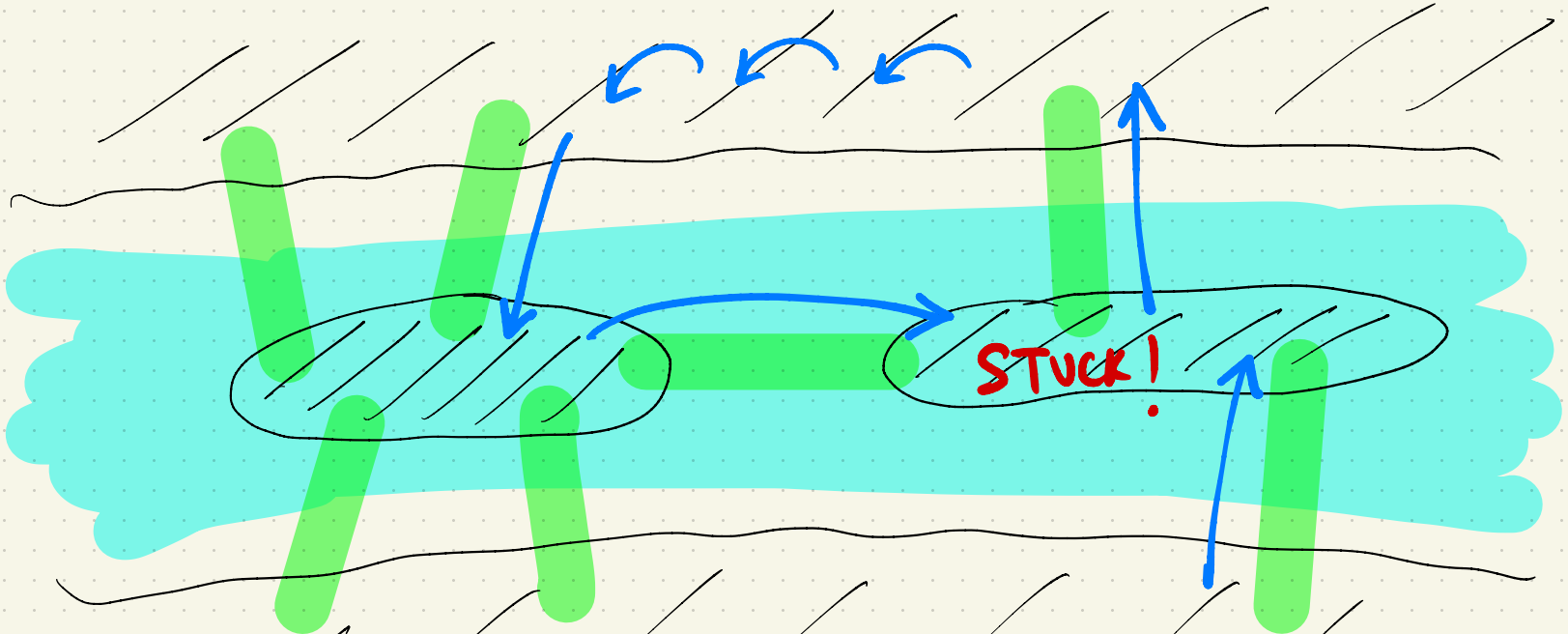
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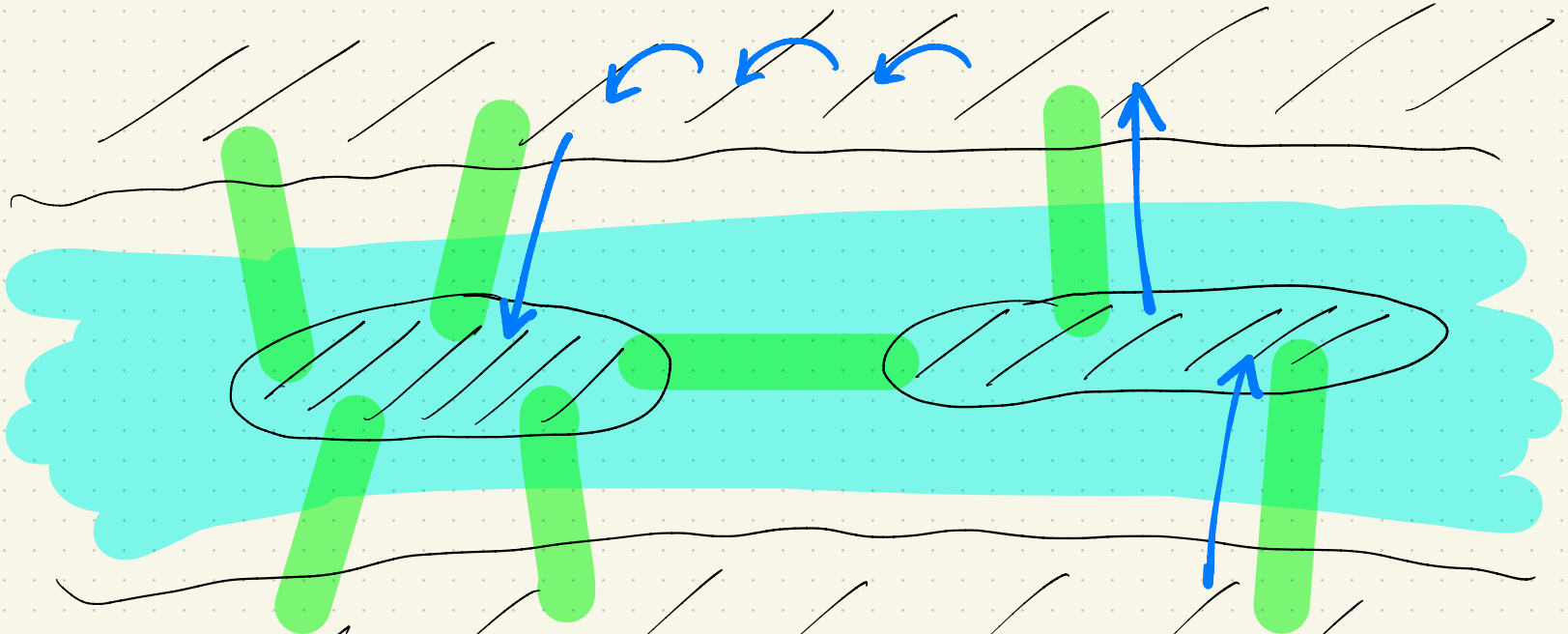
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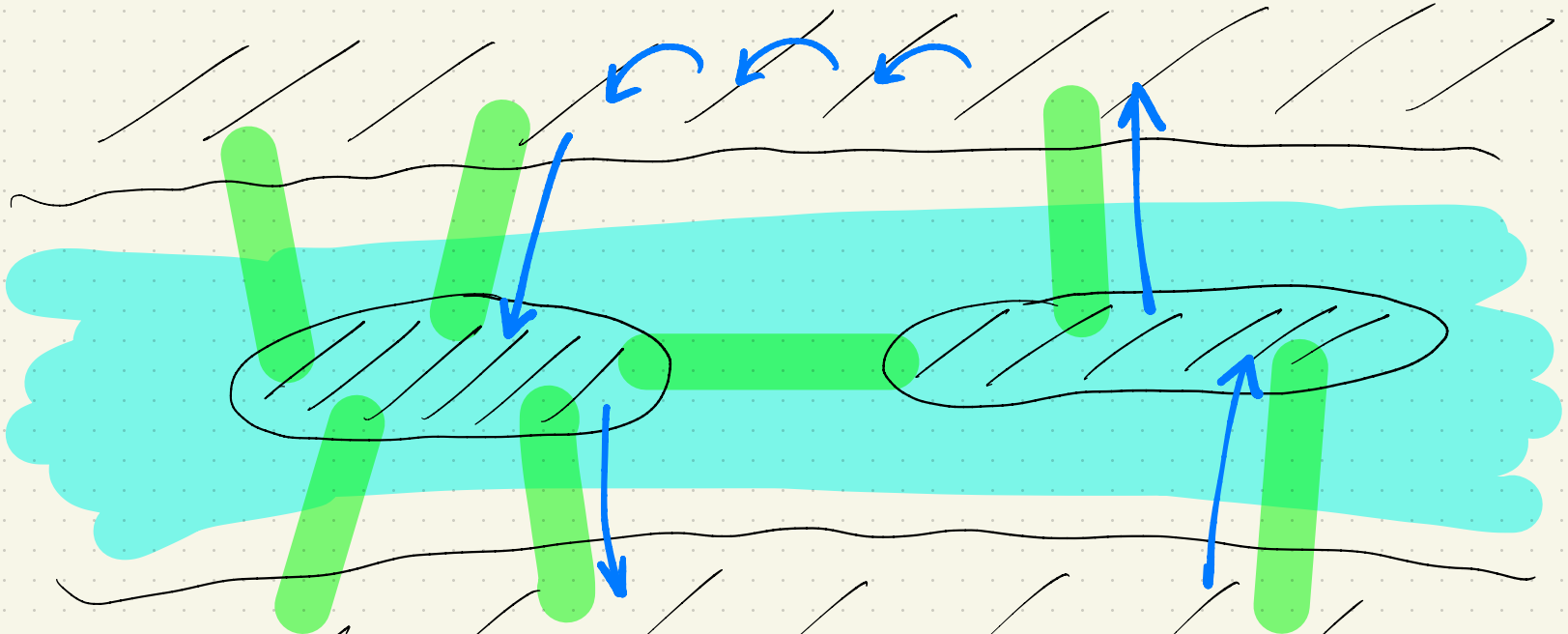
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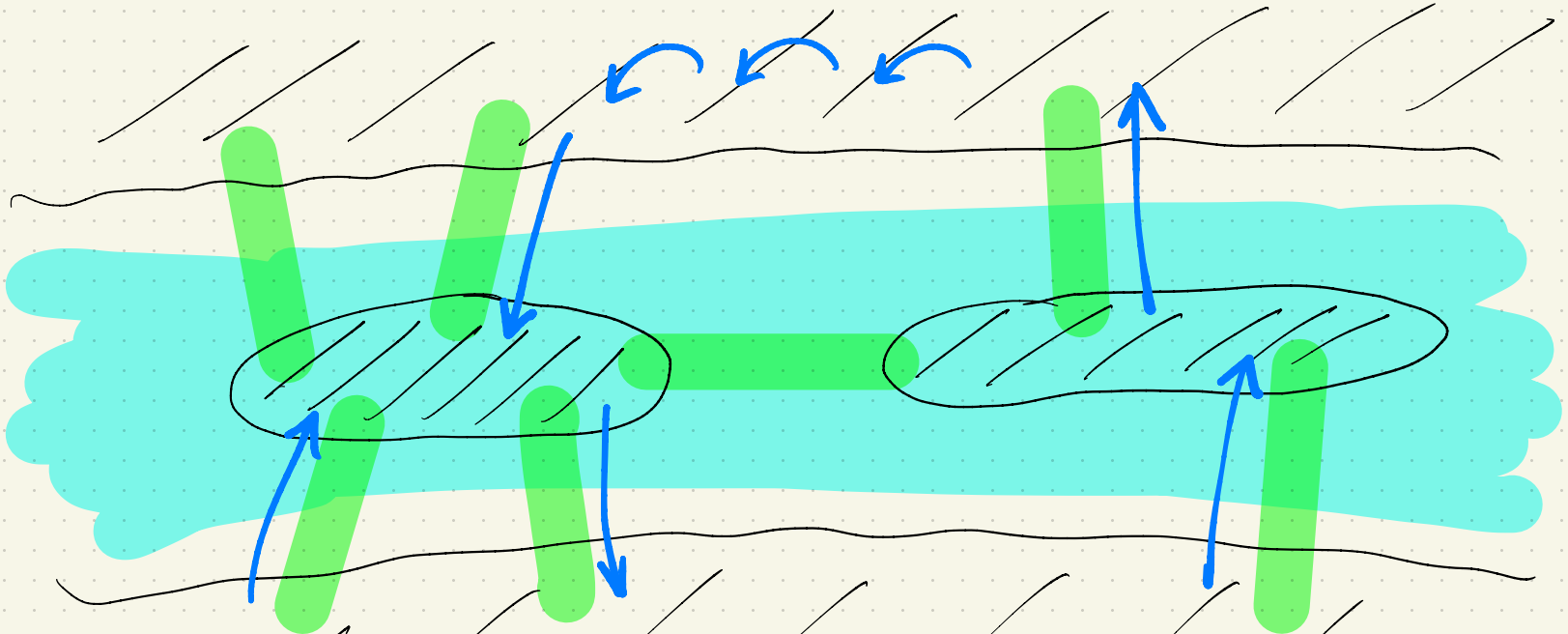
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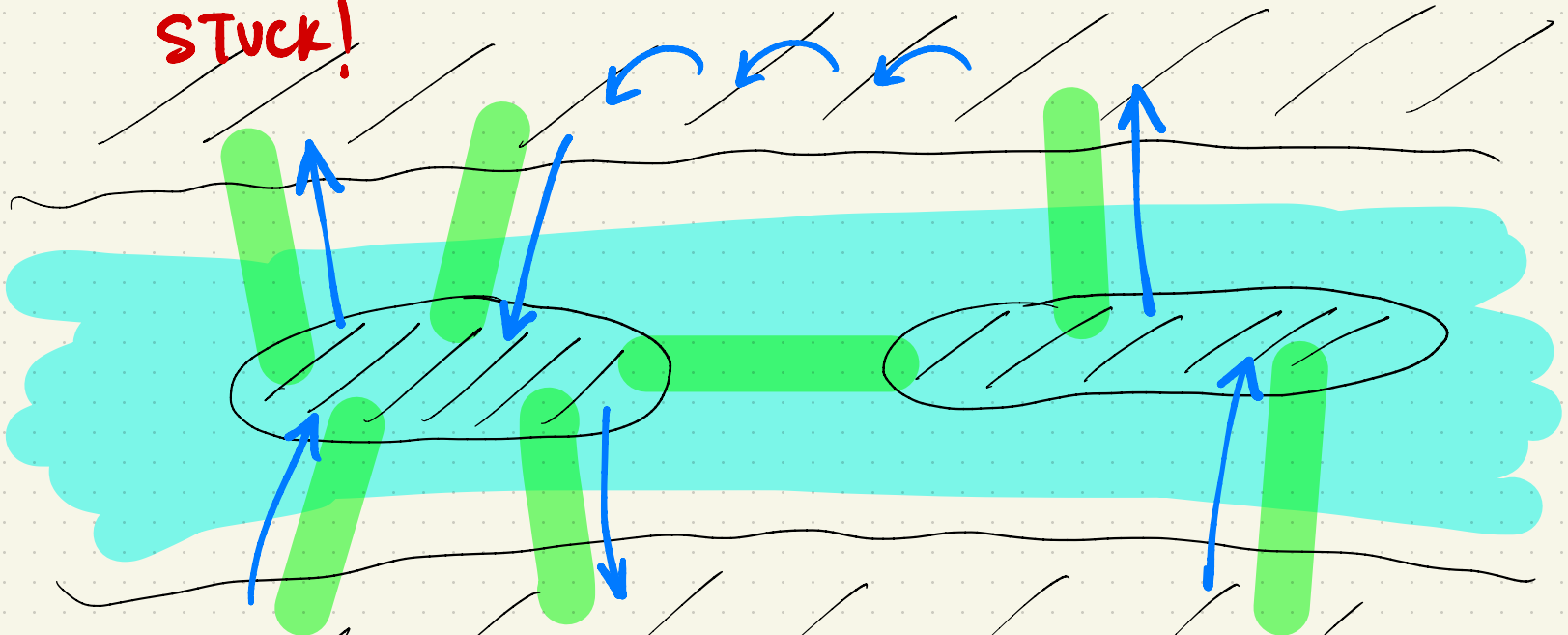


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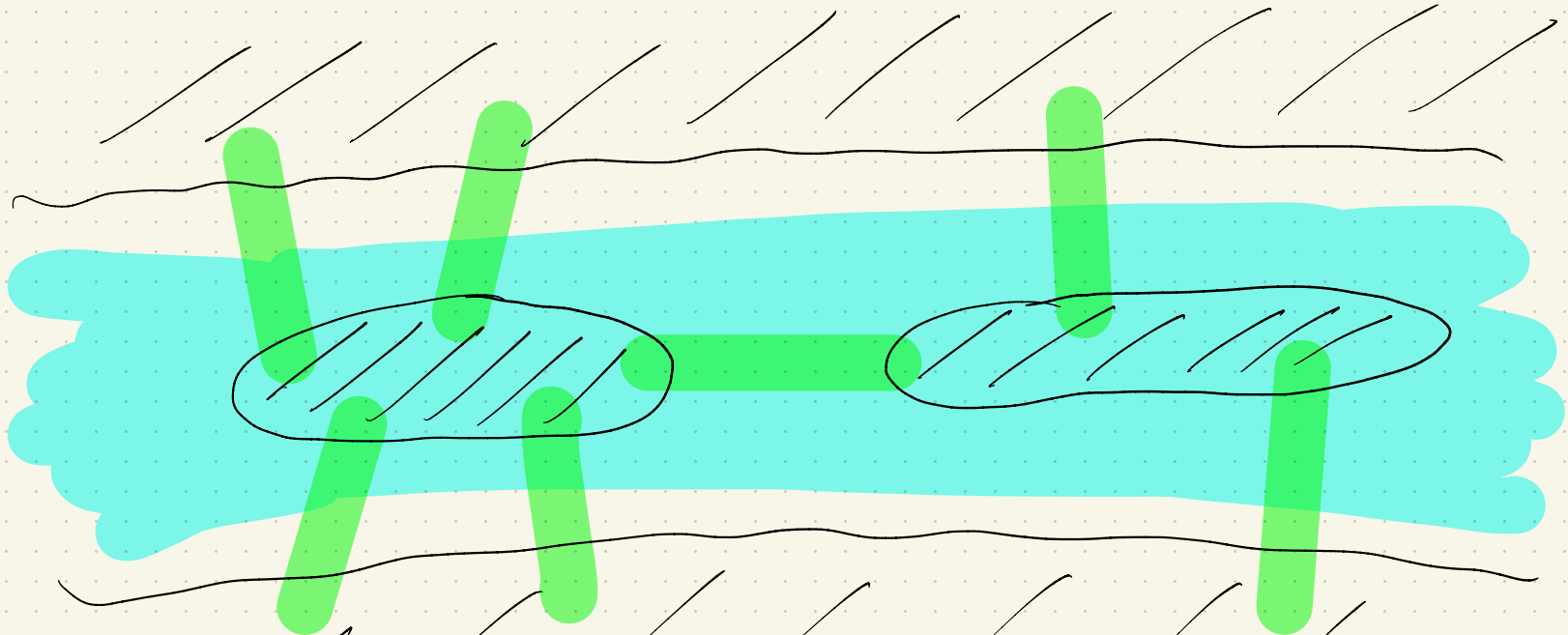
STUCK!



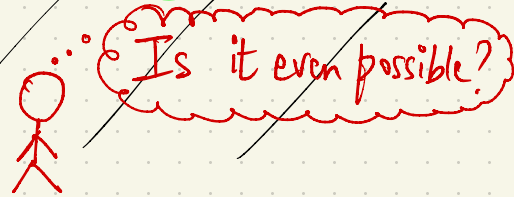
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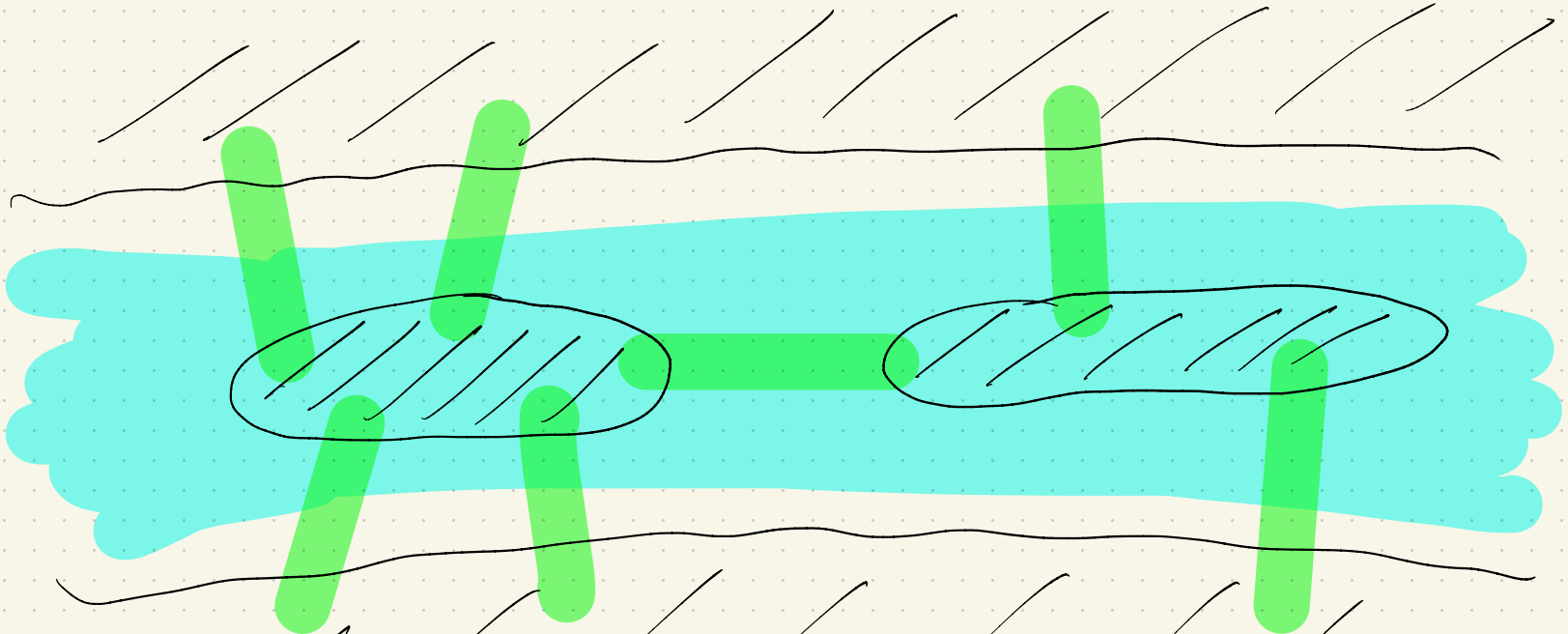
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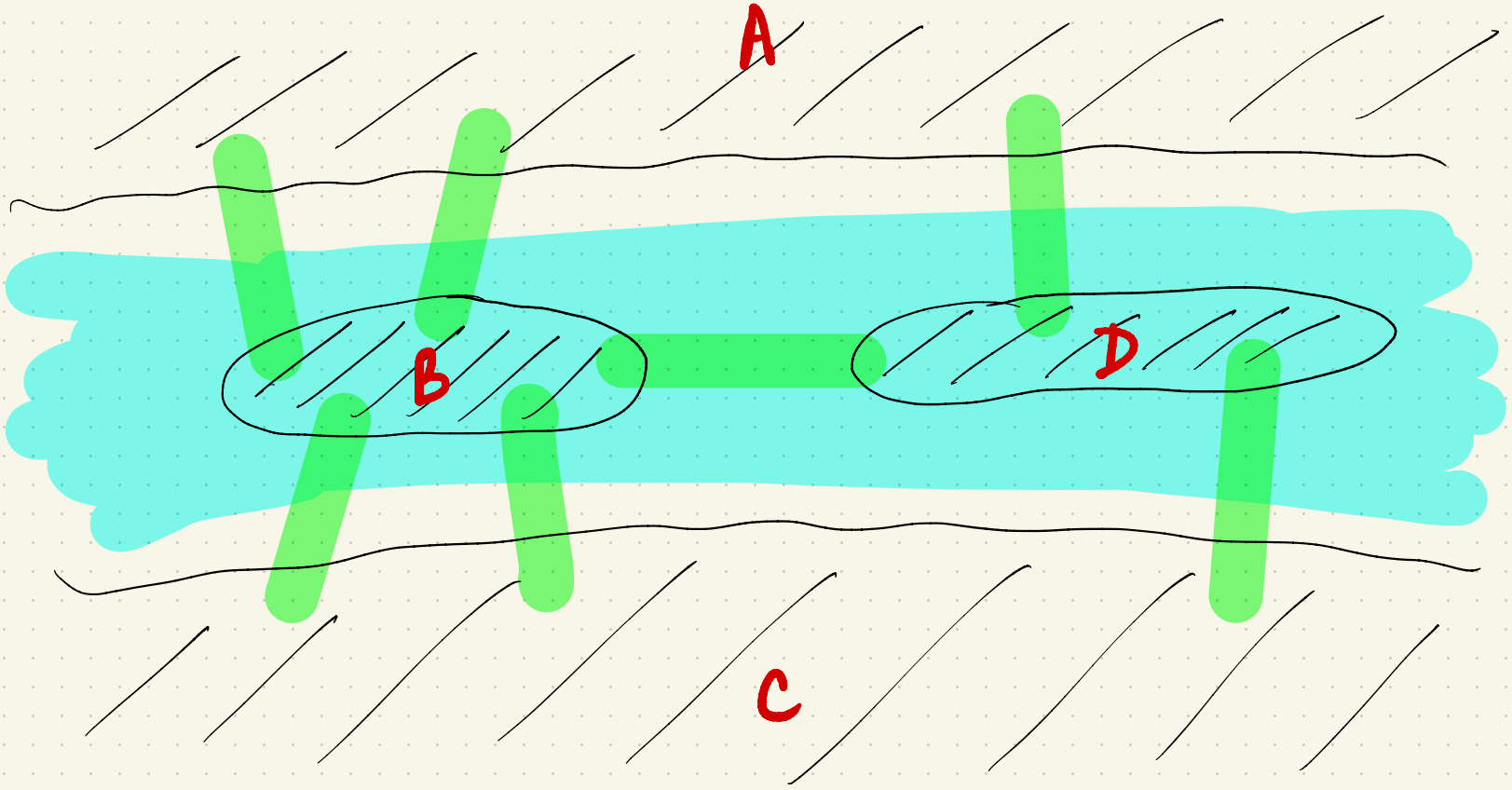
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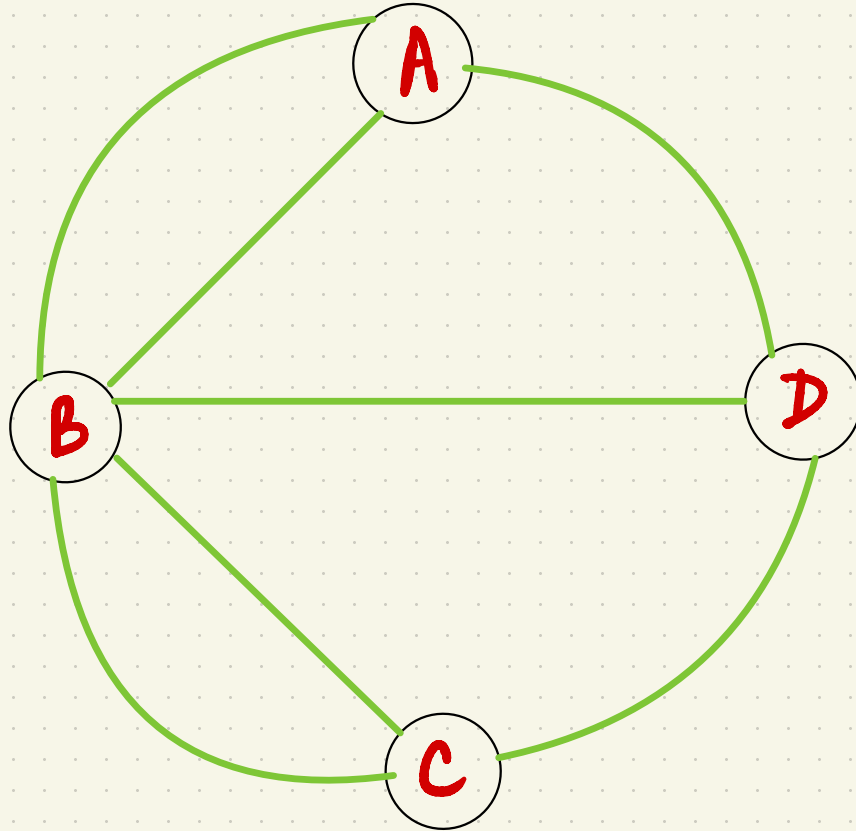
Return to starting point after crossing each bridge exactly once.

Is it even possible?
"NO!" — Euler

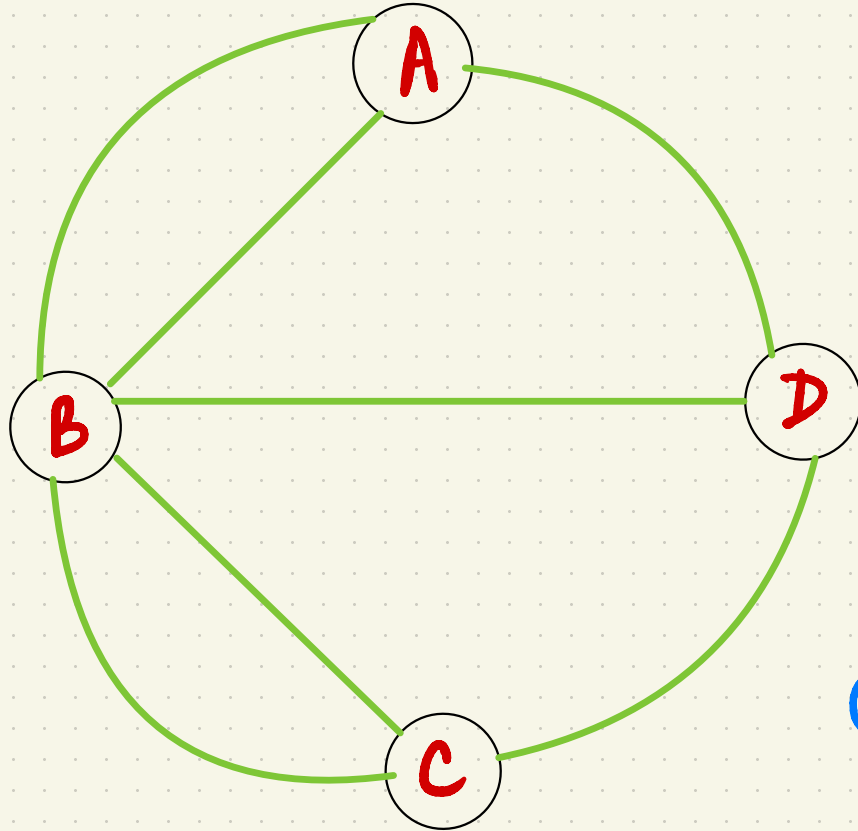
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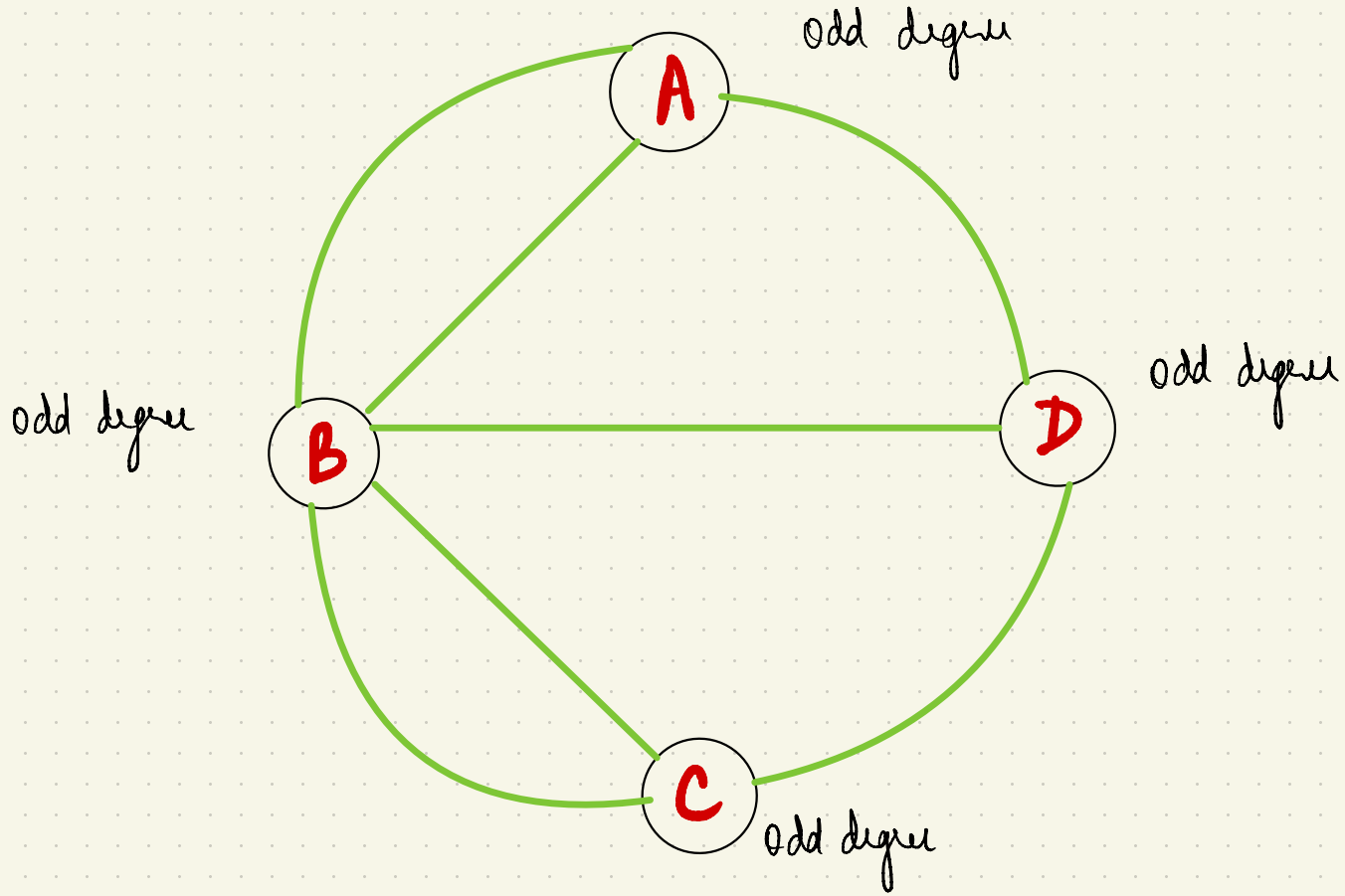


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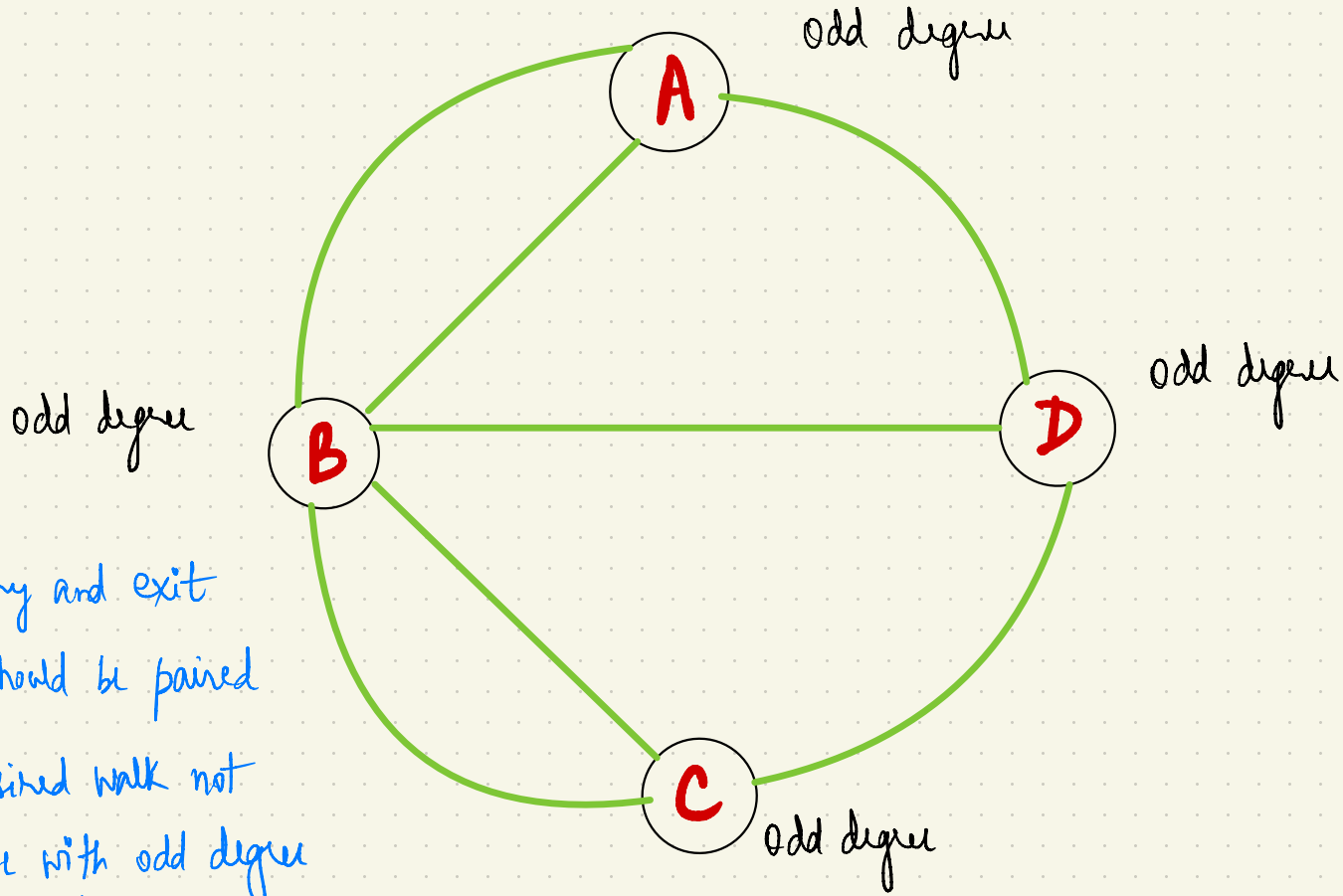


Birth of
Graph theory!
(1736)

SEVEN BRIDGES OF KÖNIGSBERG

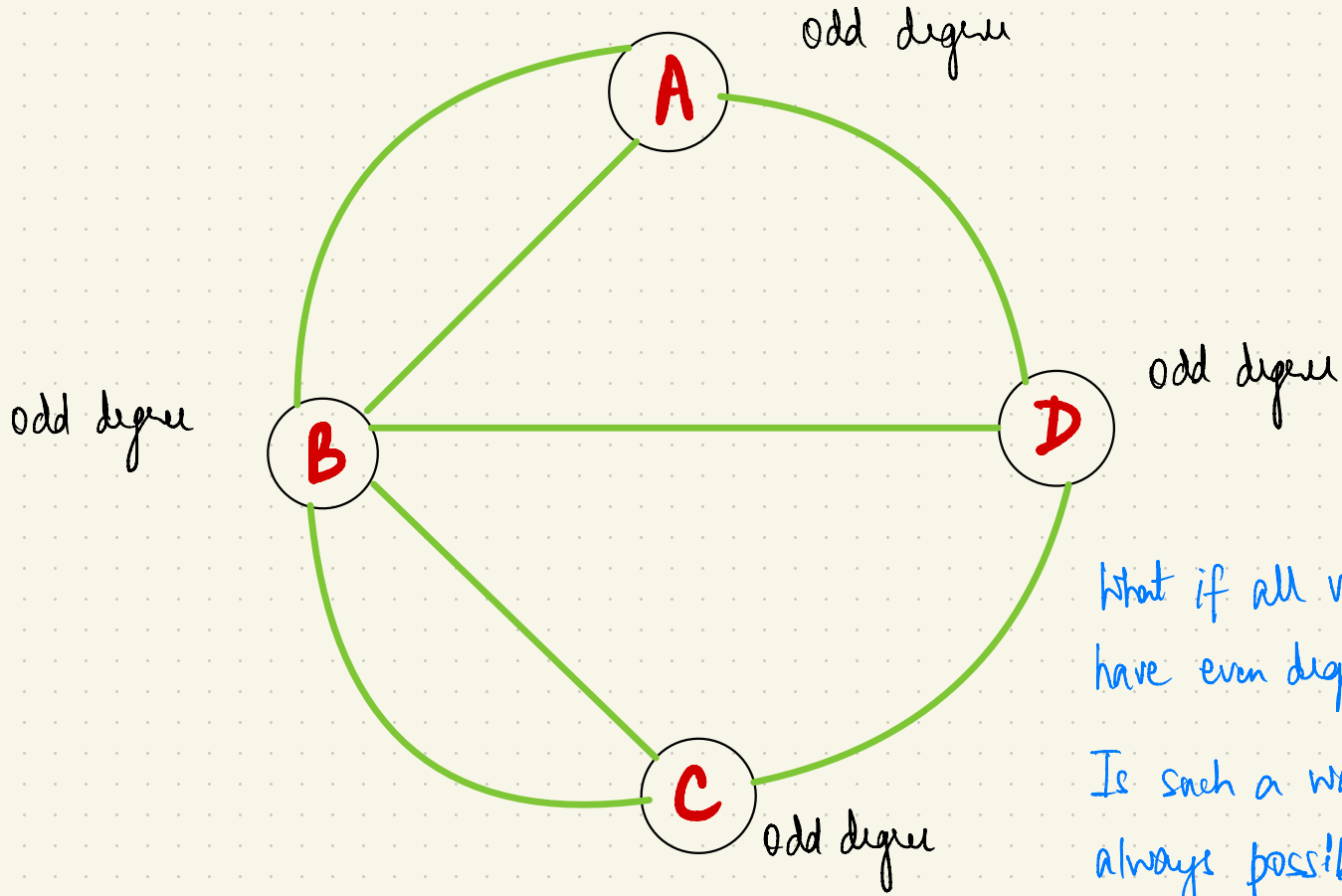


SEVEN BRIDGES OF KÖNIGSBERG



The entry and exit
edges should be paired
⇒ Desired walk not
possible with odd degree
vertices

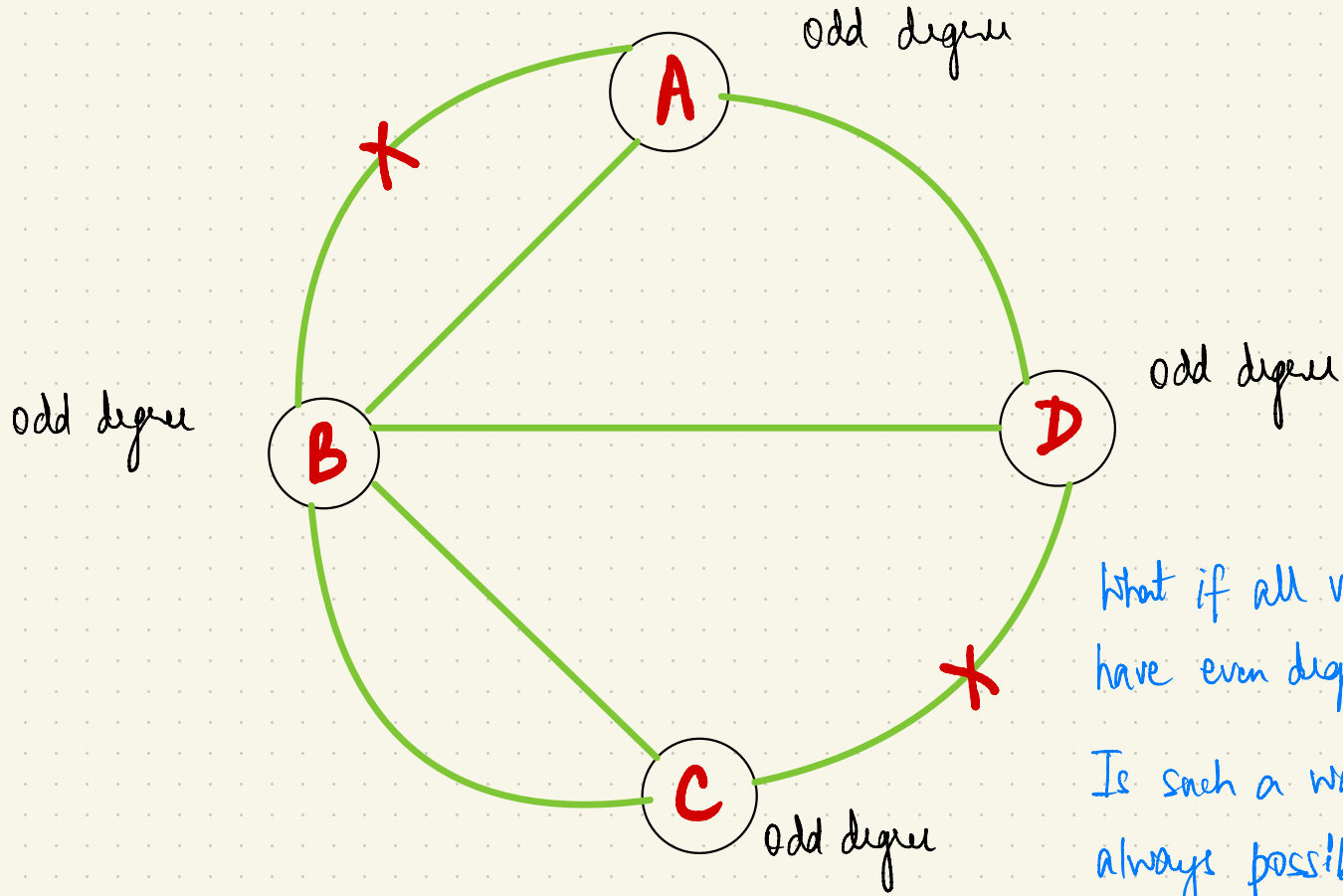
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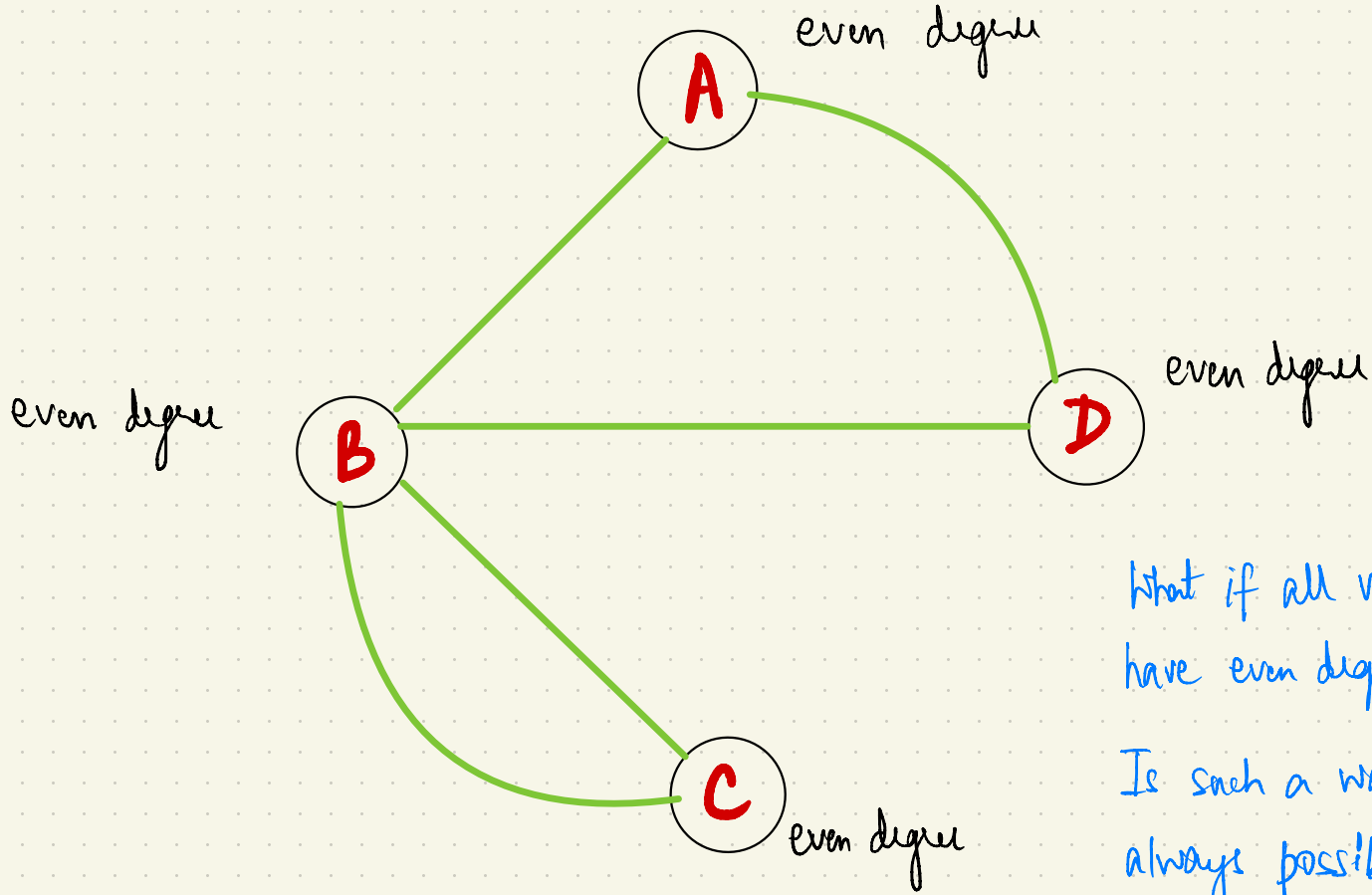
What if all vertices
have even degree?

Is such a walk
always possible?

SEVEN BRIDGES OF KÖNIGSBERG



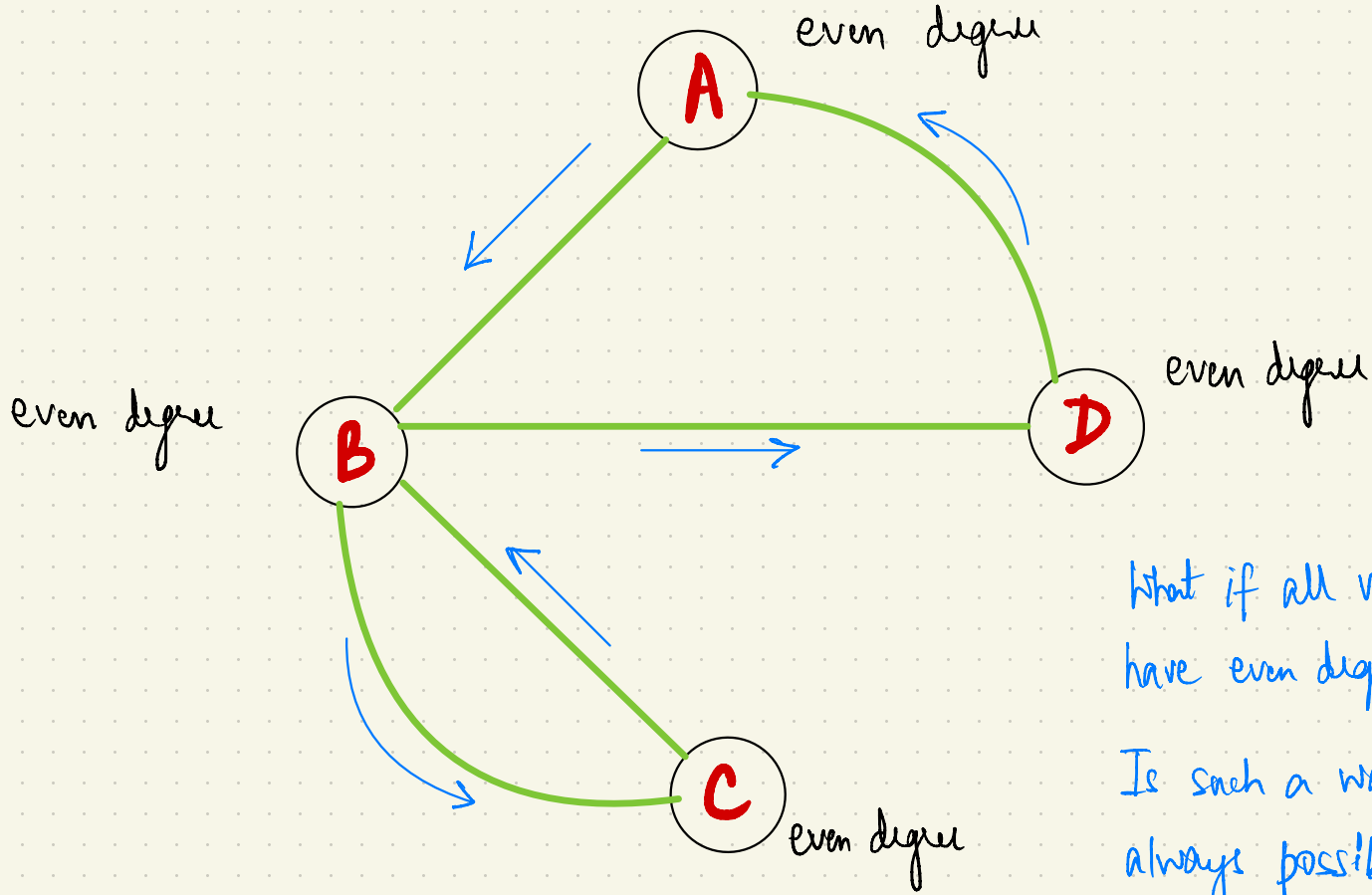
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Is such a walk
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EULER TOUR

A *tour* is a closed walk that visits every vertex at least once.

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A *tour* is a closed walk that visits every vertex at least once.

An *Euler tour* is a *tour* that traverses every *edge* exactly once.

EULER TOUR

A **tour** is a closed walk that visits every vertex at least once.

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a closed walk that visits every vertex **at least** once and every edge **exactly** once.

When does a graph admit an Euler tour?

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Theorem: A connected graph has an Euler tour if and only if every vertex has an even degree.

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* "local" information gives us "global" information
↓ ↓
individual degrees ability to traverse the graph

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* Do not need to know how the nodes are connected

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* Do not need to know **how** the nodes are connected

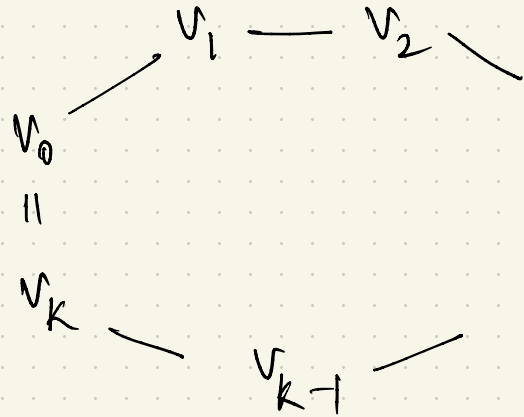
* **Easy** to check if a graph has an Euler tour.
↘ polynomial time

Theorem: A connected graph has an Euler tour if and only if every vertex has an even degree.

Proof: (\Rightarrow) Suppose $G=(V, E)$ has an Euler tour.

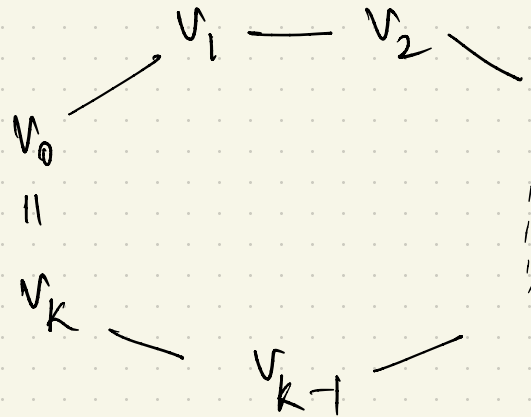
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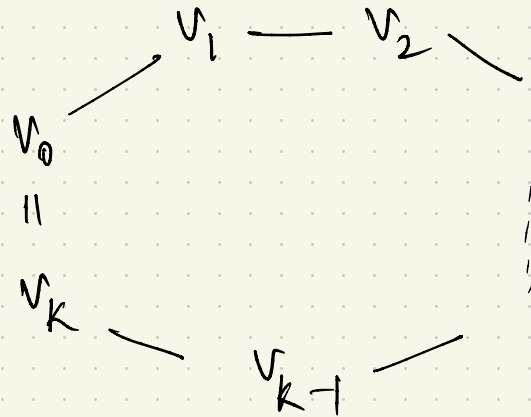
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Every edge in E is traversed exactly once \Rightarrow

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Every edge in E is traversed exactly once \Rightarrow

$$\deg(v) = 2 \times \# \text{ times } v \text{ appears in the tour } v_0, \dots, v_{k-1}, v_k$$

Theorem: A connected graph has an Euler tour if and only if every vertex has an even degree.

Proof: (\Leftarrow) For $G=(V,E)$, assume $\deg(v)$ is even for all $v \in V$.

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Let $W = v_0 - v_1 - \dots - v_k$ be the longest walk that traverses any edge at most once. (well-defined)

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Obs 1: Any edge in E that is incident to v_k is covered by W .

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(Otherwise W may be extended — not possible)

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Obs 1: Any edge in E that is incident to v_k is covered by W .
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Obs 2: $v_0 = v_k$
(Otherwise v_k has odd no. of edges in $W \Rightarrow$ odd degree)

Theorem: A connected graph has an Euler tour if and only if every vertex has an even degree.

Proof: (\Leftarrow) For $G = (V, E)$, assume $\deg(v)$ is even for all $v \in V$.

Let $W = v_0 - v_1 - \dots - v_k$ be the longest walk that traverses any edge at most once.

We now want to show that W is an Euler tour.

Theorem: A connected graph has an Euler tour if and only if every vertex has an even degree.

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Let $W = v_0 - v_1 - \dots - v_k$ be the longest walk that traverses any edge at most once.

Suppose W is not an Euler tour.

Theorem: A connected graph has an Euler tour if and only if every vertex has an even degree.

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Let $W = v_0 - v_1 - \dots - v_k$ be the longest walk that traverses any edge at most once.

Suppose W is not an Euler tour.

\Rightarrow There is an edge, say e , in E such that e is **not** in W

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\rightarrow why?

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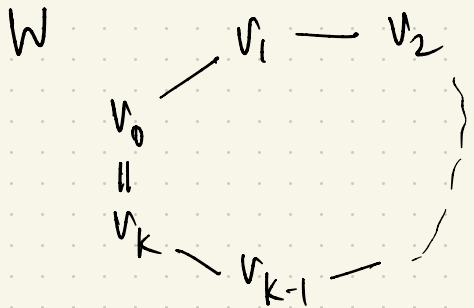
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 \rightarrow **Why?** Connectedness of graph G

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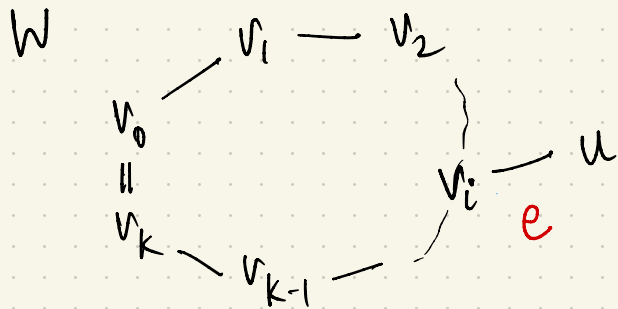
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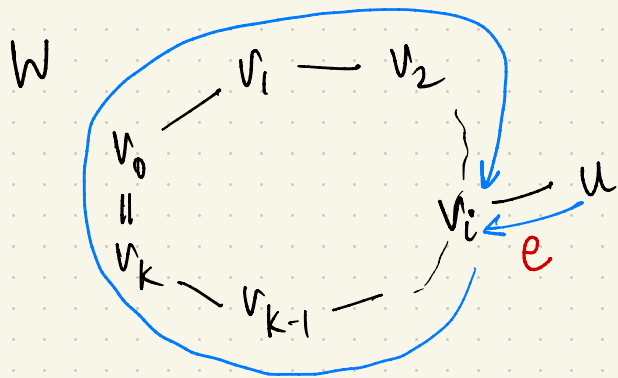
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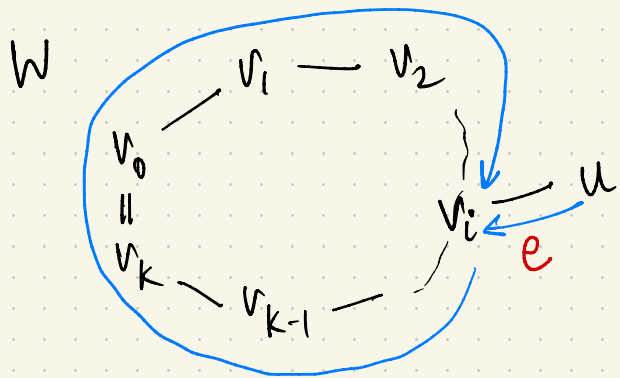
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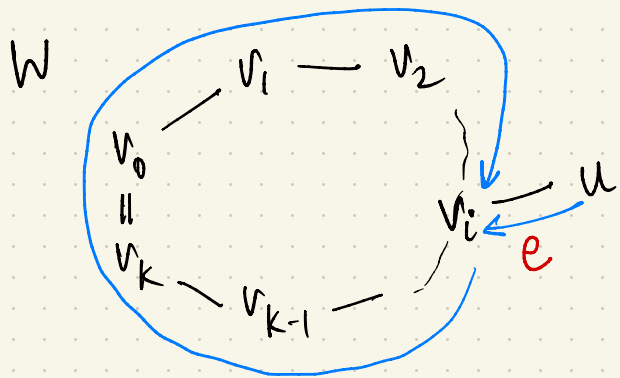
Can make W longer!

$$u - v_i - v_{i+1} - \dots - v_k = v_0 - \dots - v_i$$

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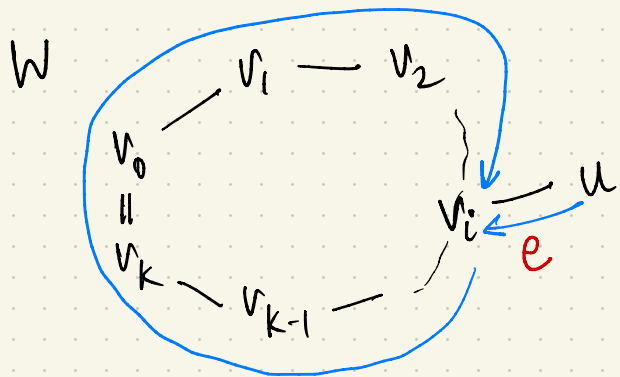
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Contradiction! $\Rightarrow W$ must be an Euler tour.

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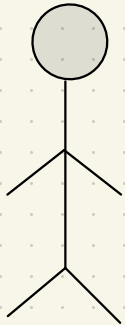
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AN APPLICATION OF EULER TOURS

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Genome Sequencing



= ... TACGAGACAGTACA ...

~ 3 billion letters

AN APPLICATION OF EULER TOURS

Multiple identical copies
of a genome

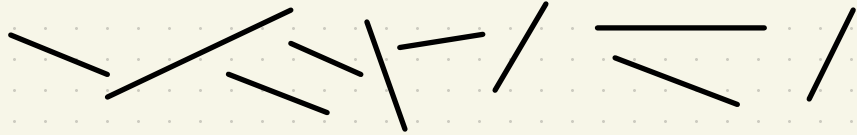


AN APPLICATION OF EULER TOURS

Multiple identical copies
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Shatter the genome
into reads



AN APPLICATION OF EULER TOURS

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Shatter the genome
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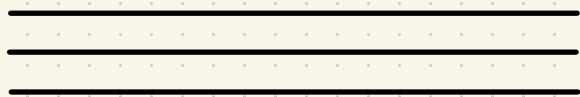
Sequence the heads

AGATATCA

CGATCCAT

AN APPLICATION OF EULER TOURS

Multiple identical copies
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Shatter the genome
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Sequence the heads

AGATATCA

CGATCCAT

Assemble the genome
using overlapping heads

```
AGATCCG
GATCCGA
TAGATCC
-----
TAGATCCGA
```

AN APPLICATION OF EULER TOURS

Multiple identical copies
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Shatter the genome
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Sequence the heads

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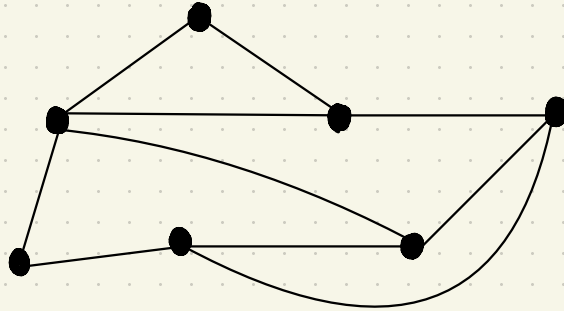
Assemble the genome
using overlapping heads

```
  A G A T C C G
    G A T C C G A
  T A G A T C C
  -----
  T A G A T C C G A
```

Euler
tours!

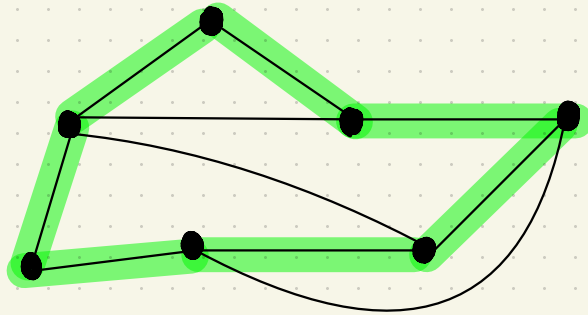
HAMILTONIAN CYCLE

A cycle that visits every vertex exactly once



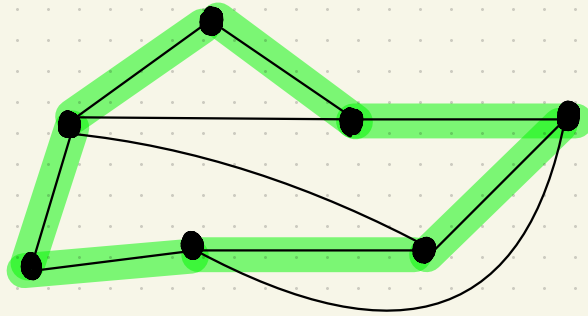
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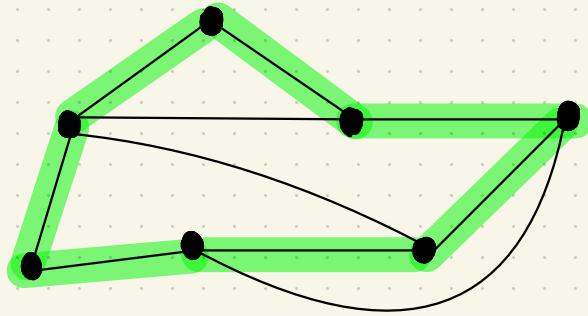
A cycle that visits every vertex exactly once



How to check if a graph has a Hamiltonian cycle?

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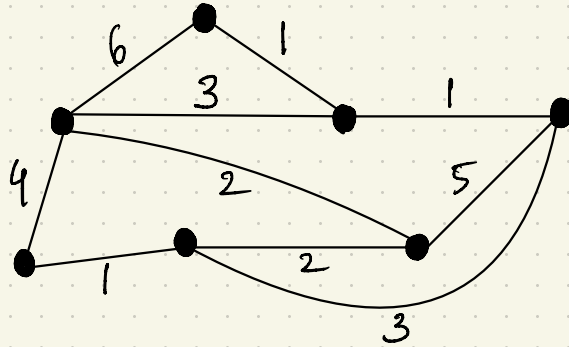
How to check if a graph has a Hamiltonian cycle?

NP-complete



HAMILTONIAN CYCLE

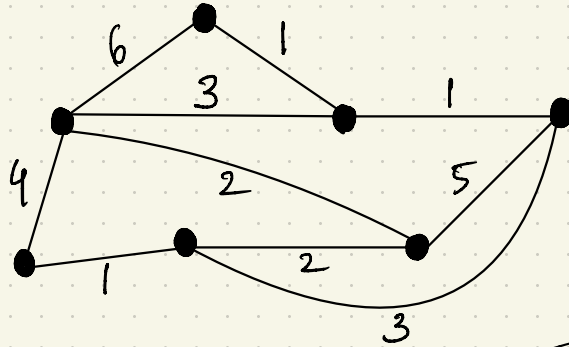
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How to find a minimum-weight Hamiltonian cycle?

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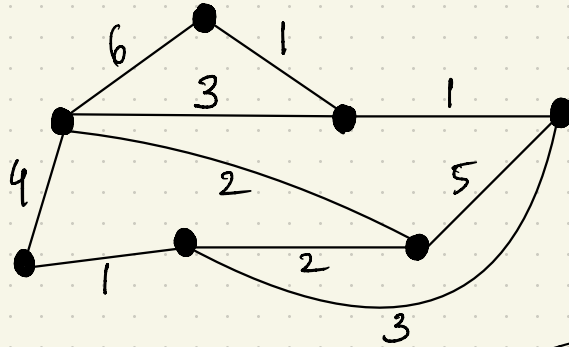


→ e.g., weight ≤ 15

How to find a minimum-weight Hamiltonian cycle?

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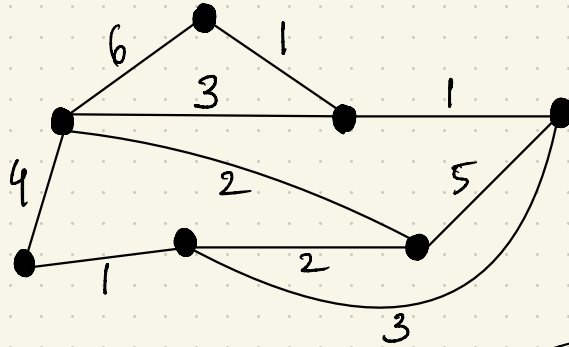
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How to find a minimum-weight Hamiltonian cycle?

NP-complete 😞

HAMILTONIAN CYCLE

A cycle that visits every vertex exactly once



→ e.g., weight ≤ 15

How to find a minimum-weight Hamiltonian cycle?

NP-complete



→ Traveling Salesperson Problem