

COL 202: DISCRETE MATHEMATICAL STRUCTURES

LECTURE 18

PERFECT MATCHINGS

FEB 09, 2024

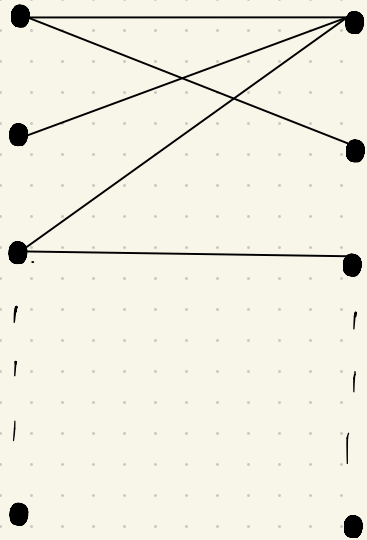
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ROHIT VAISH

GRAPH THEORY

On average, who has more opposite-gender partners: men or women?

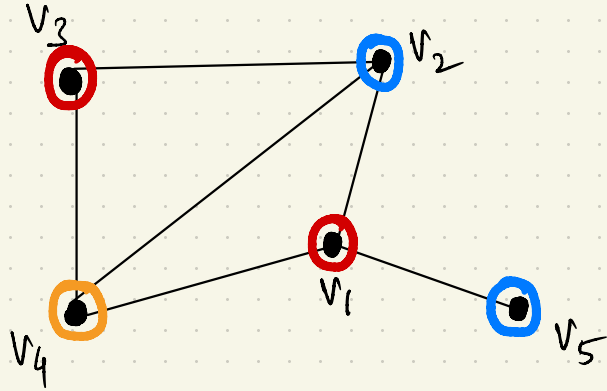
Women Men



$$\frac{\text{Avg. deg. of men}}{\text{Avg. deg. of women}} = \frac{\text{No. of women}}{\text{No. of men}}$$

On average, men have more opposite-gender partners than women if $\# \text{ women} > \# \text{ men}$.

GRAPH COLORING



Properly color a given graph with as few colors as possible.

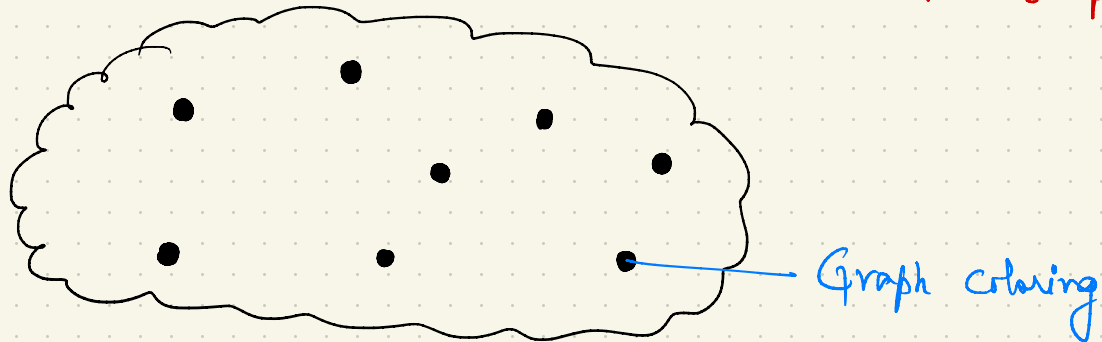
GRAPH COLORING

Given a coloring, checking if it is a **valid** solution: **Easy!**
proper $\leftarrow \rightarrow \leq k$ colors

Given graph G and k , determining if G admits a proper **k-coloring**: **Hard!**

NP-complete

A class of
"equally hard"
problems



BASIC GRAPH COLORING ALGORITHM

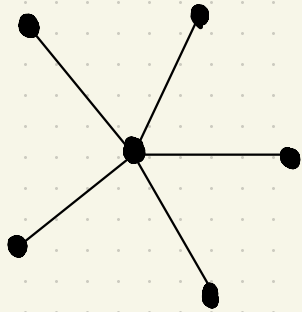
1. Order the vertices v_1, v_2, \dots, v_n
2. Order the colors c_1, c_2, \dots
3. For $i = 1, 2, \dots, n$
 - └ Assign the least-index legal color to v_i

BASIC GRAPH COLORING ALGORITHM

Theorem: If every vertex in G has degree at most d ,
then Basic Algo. uses at most $d+1$ colors.

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"Star" graph

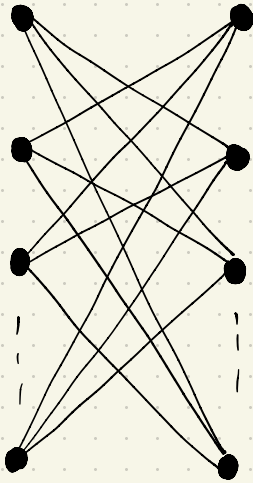
$$\text{degree} = n - 1$$

$$\chi(G) = 2$$

Basic algo will also use two colors
under any ordering of vertices.

BASIC GRAPH COLORING ALGORITHM

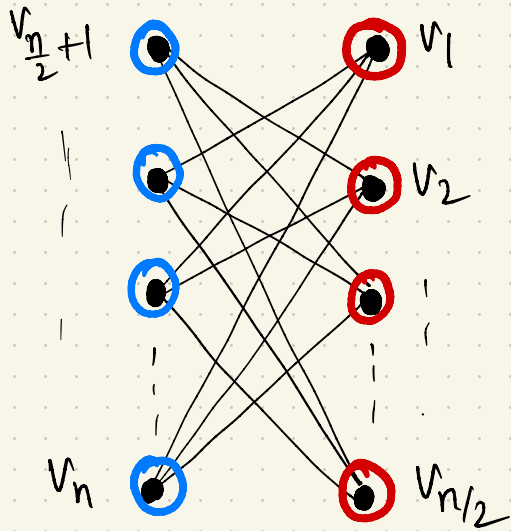
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all edges except the horizontal ones

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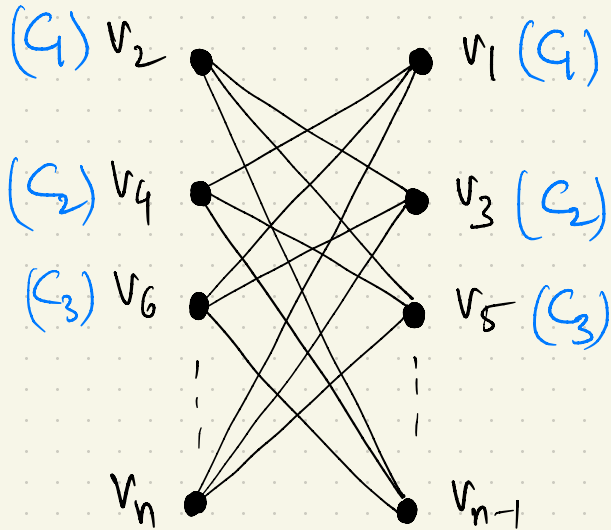


A good ordering \rightarrow 2 colors

all edges except the horizontal ones

BASIC GRAPH COLORING ALGORITHM

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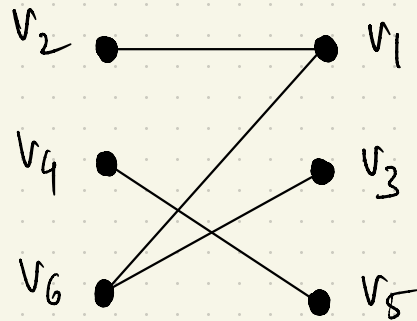


A bad ordering $\rightarrow n/2$ colors

Ordering matters!

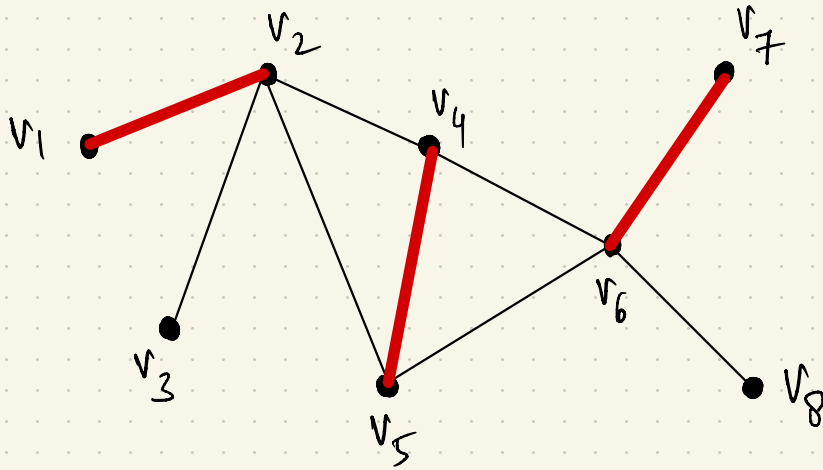
all edges except the horizontal ones

A graph $G = (V, E)$ is **bipartite** if V can be partitioned into two sets V_L and V_R such that all edges in E connect a vertex in V_L with a vertex in V_R .



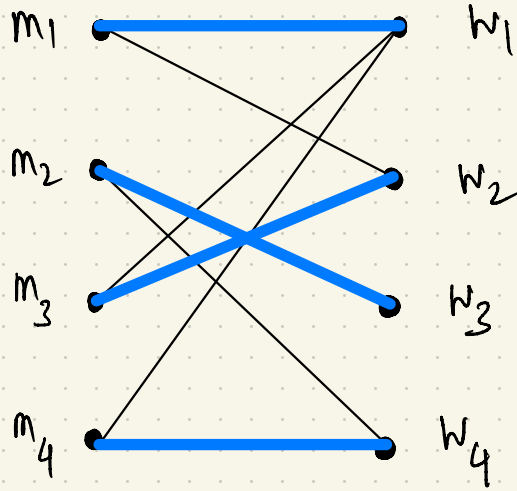
MATCHING

Given a graph $G = (V, E)$ that is not necessarily bipartite
a **matching** is a subgraph of G where every vertex
has degree one.



MATCHING

A matching is *perfect* if it has size $|V|/2$.



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When does a bipartite graph have a perfect matching?

BIPARTITE PERFECT MATCHING

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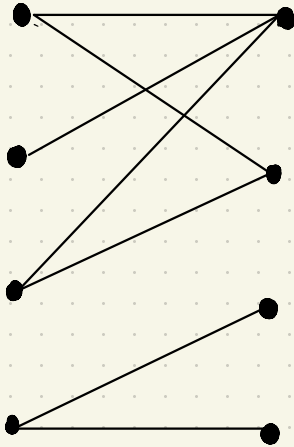
obviously NOT

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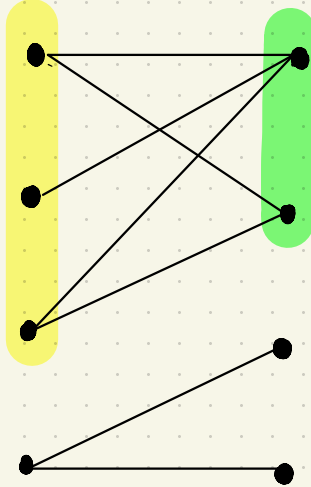
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BIPARTITE PERFECT MATCHING

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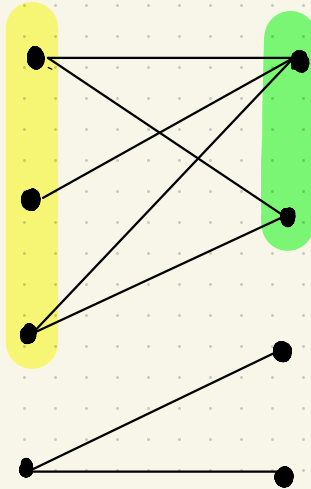
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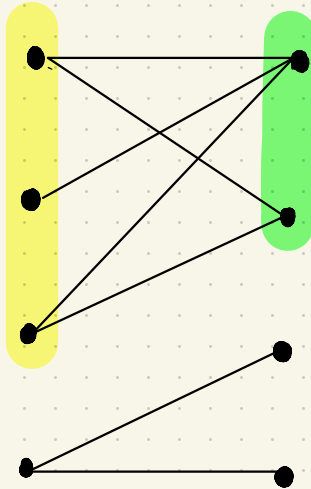


No perfect matching because
the yellow vertices do not have
"sufficiently many" neighbors

BIPARTITE PERFECT MATCHING

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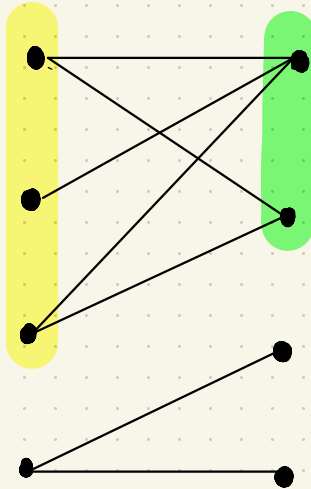
A bare minimum condition

Every subset of left vertices $S \subseteq V_L$
has at least $|S|$ neighbors on the right.

BIPARTITE PERFECT MATCHING

obviously NOT

When does a bipartite graph have a perfect matching?



necessary

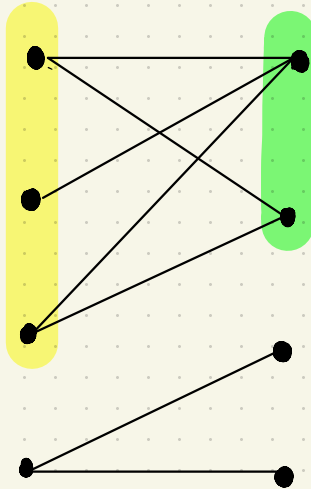
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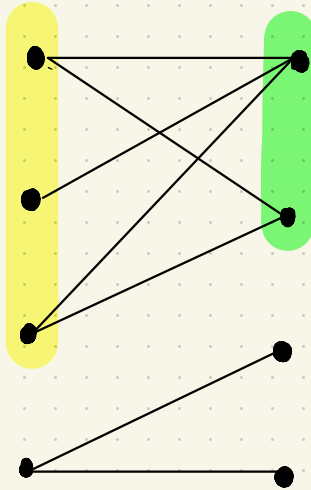
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BIPARTITE PERFECT MATCHING

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When does a bipartite graph have a perfect matching?



A necessary condition

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Is it sufficient?