

## Tutorial Sheet 2

Announced on: Aug 01 (Thurs)

Problems marked with (★) will not be asked in the tutorial quiz.

1. Consider the following algorithm for checking whether a given number  $n$  is prime.

---

**ALGORITHM 1:** Checking primality

---

**Input:** An integer  $n > 1$ .  
**Output:** Yes/No

```

1 if  $n$  equals 2 or 3 then
2   | return Yes
3 for  $i = 2$  to  $\lfloor \sqrt{n} \rfloor$  do
4   | if  $i$  divides  $n$  then
5     | | return No
6 return Yes

```

---

Prove or disprove:

- The algorithm is correct.
  - The algorithm runs in polynomial time.
2. Consider the following algorithm for calculating  $a^b$  where  $a$  and  $b$  are positive integers.

---

**ALGORITHM 2:** FastPower( $a, b$ )

---

**Input:** Positive integers  $a$  and  $b$ .  
**Output:**  $a^b$ .

```

1 if  $b = 1$  then
2   | return  $a$ 
3  $c := a \cdot a$ 
4  $d := \text{FastPower}(c, \lfloor b/2 \rfloor)$ 
5 if  $b$  is odd then
6   | return  $a \cdot d$ 
7 return  $d$ 

```

---

Suppose each multiplication and division operation can be performed in constant time. Determine the asymptotic running time of FastPower as a function of  $b$ .

3. Let  $A$  and  $B$  be two sorted arrays of length  $n$  each. Assume that all elements within and across the two arrays are distinct. Design an  $\mathcal{O}(\log n)$  algorithm to compute the  $n^{\text{th}}$  smallest

Tutorial Sheet 2:

element of the union of  $A$  and  $B$ .

4. Design an  $\mathcal{O}(\log^2 n)$  algorithm that, given a positive integer  $n$ , determines whether  $n$  is of the form  $a^b$  for some positive integers  $a$  and  $b > 1$ . For the purpose of this problem, you may assume exponentiation to be  $\mathcal{O}(1)$  time, i.e., computing  $p^q$  for two integers  $p$  and  $q$  takes constant time. Similarly, you can assume that comparison of two integers (i.e., determining whether  $p$  equals  $q$ ,  $p > q$  or  $p < q$ ) takes constant time.
5. You are given a sorted (from smallest to largest) array  $A$  of  $n$  distinct integers which can be positive, negative, or zero. You want to decide whether or not there is an index  $i$  such that  $A[i] = i$ . Design the fastest algorithm you can for solving this problem.
6. (★) You are given an  $n$ -by- $n$  grid of distinct numbers. A number is a *local minimum* if it is smaller than all its neighbors. A *neighbor* of a number is one immediately above, below, to the left, or to the right. Most numbers have four neighbors; numbers on the side have three; the four corners have two.
  - (a) Prove that a local minimum always exists.
  - (b) Use the divide-and-conquer algorithm design paradigm to compute a local minimum with only  $\mathcal{O}(n)$  comparisons between pairs of numbers. (Note: since there are  $n^2$  numbers in the input, you cannot afford to look at all of them.)
7. (★) You are given a sequence of  $n$  numbers  $a_1, a_2, \dots, a_n$ . Design an  $\mathcal{O}(n)$  algorithm to find  $i, j$  with  $i \leq j$  such that the sum  $a_i + a_{i+1} + \dots + a_j$  is maximum. Note that the numbers may not be positive.