COL866: Special Topics in Algorithms

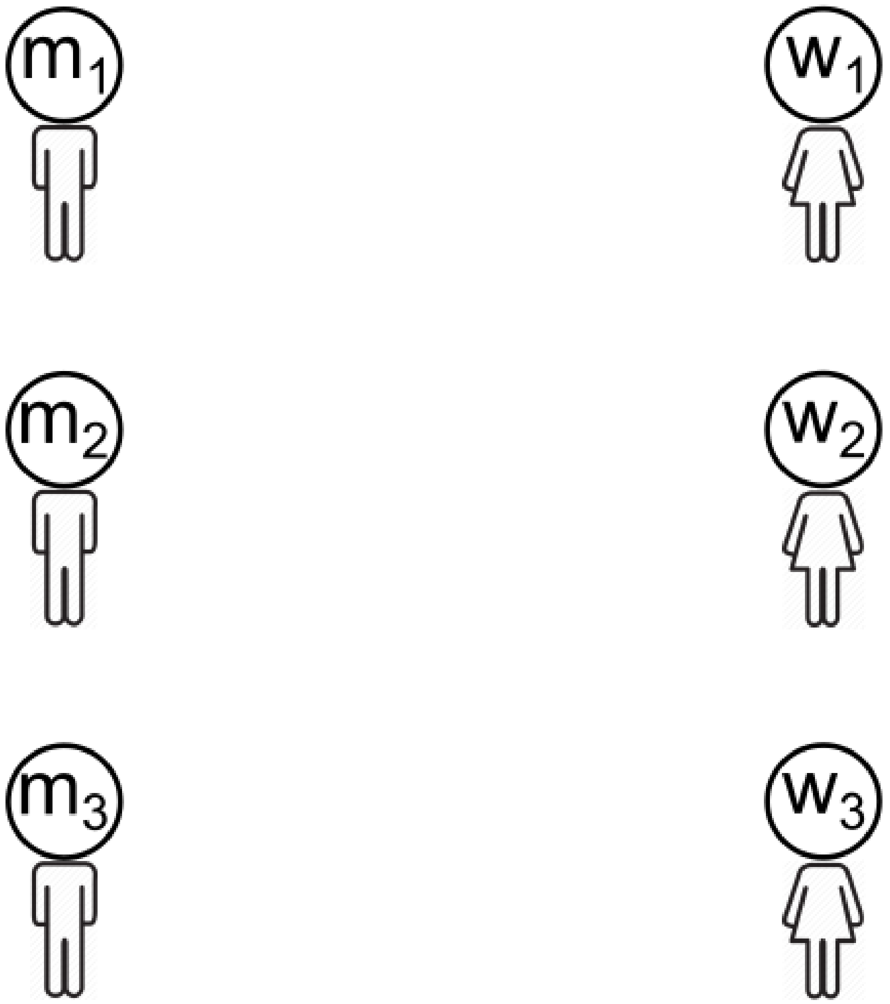
Lecture 2

Structure of Stable Matchings

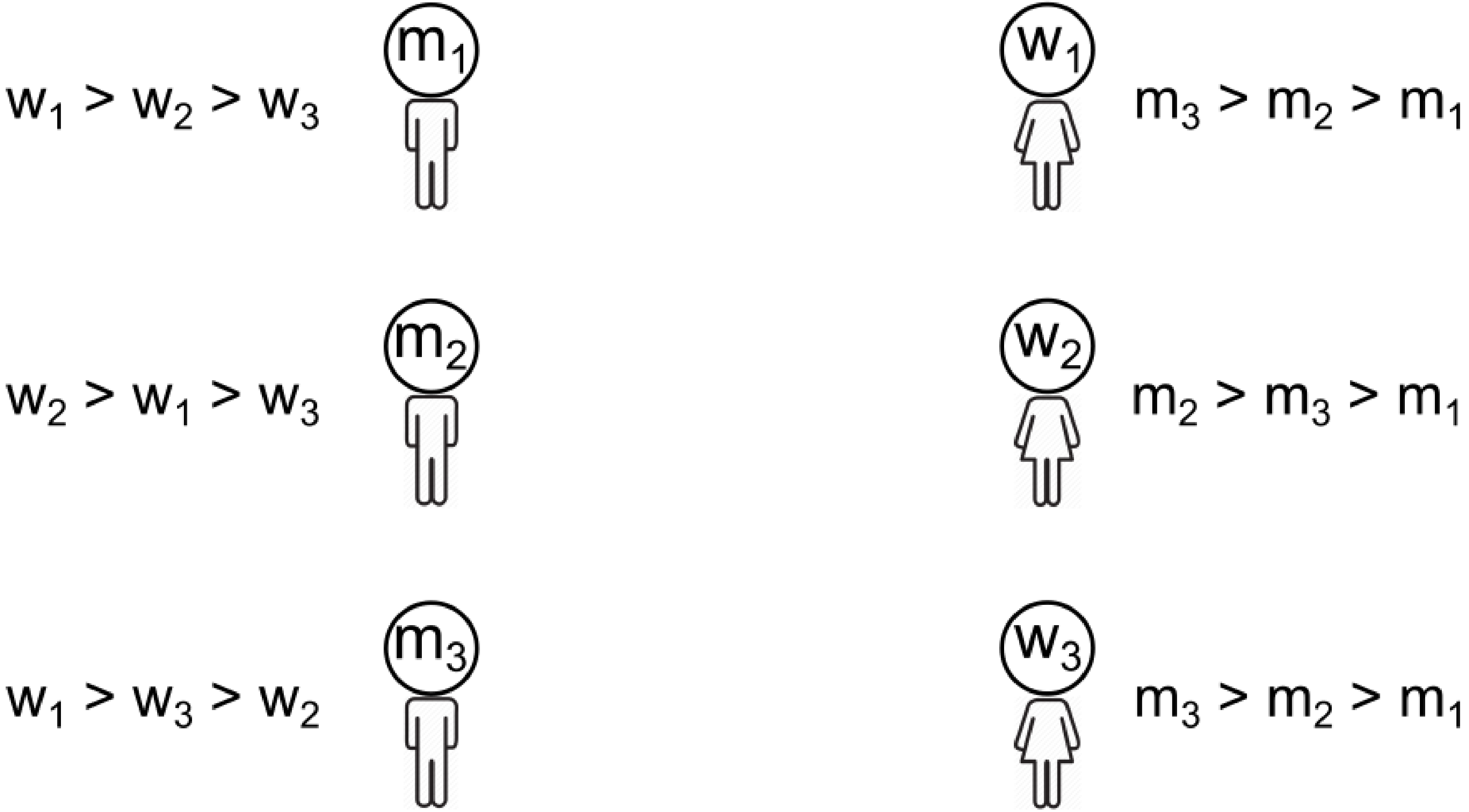
Aug 06, 2022 | Rohit Vaish

Stable Matching Problem

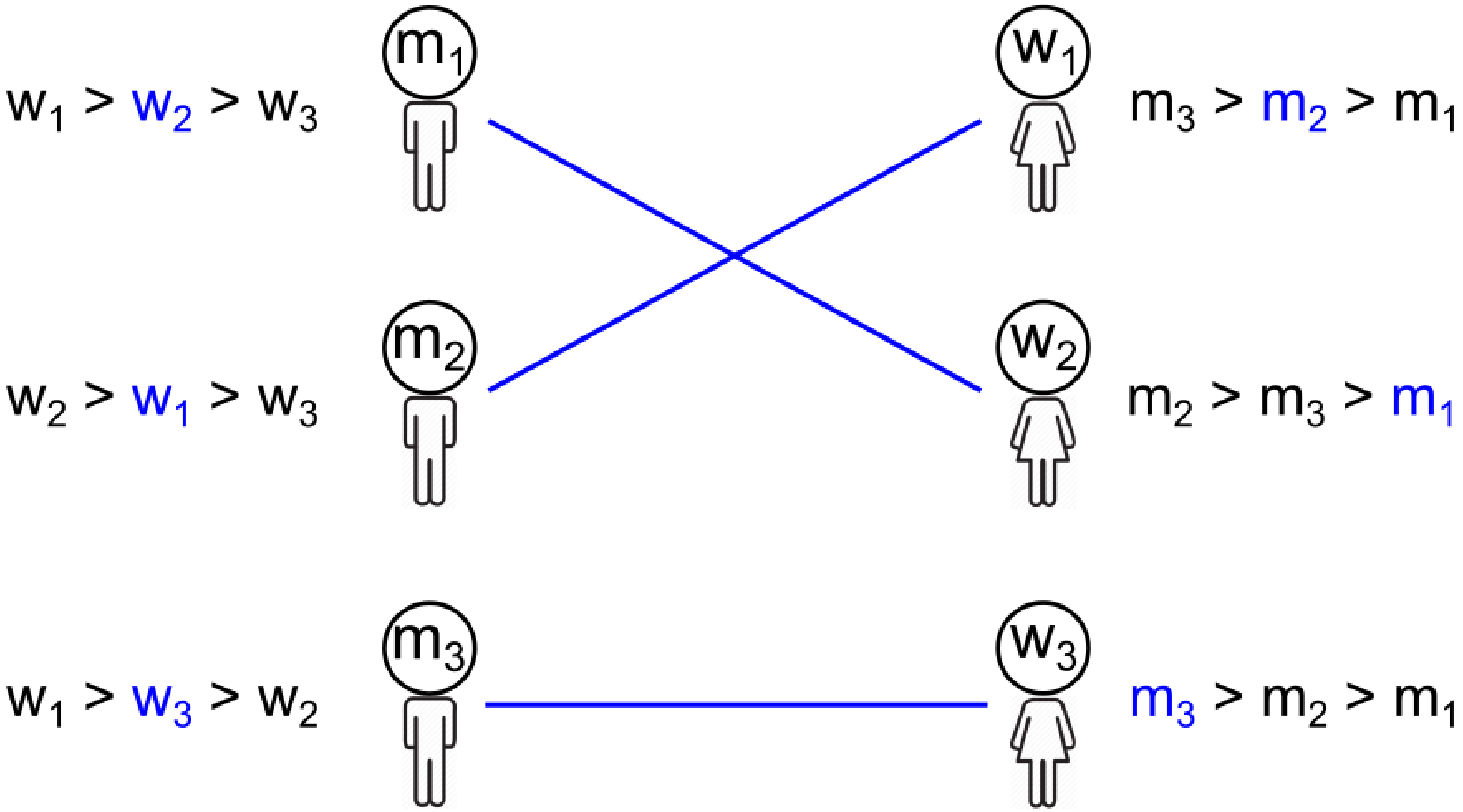
Stable Matching Problem



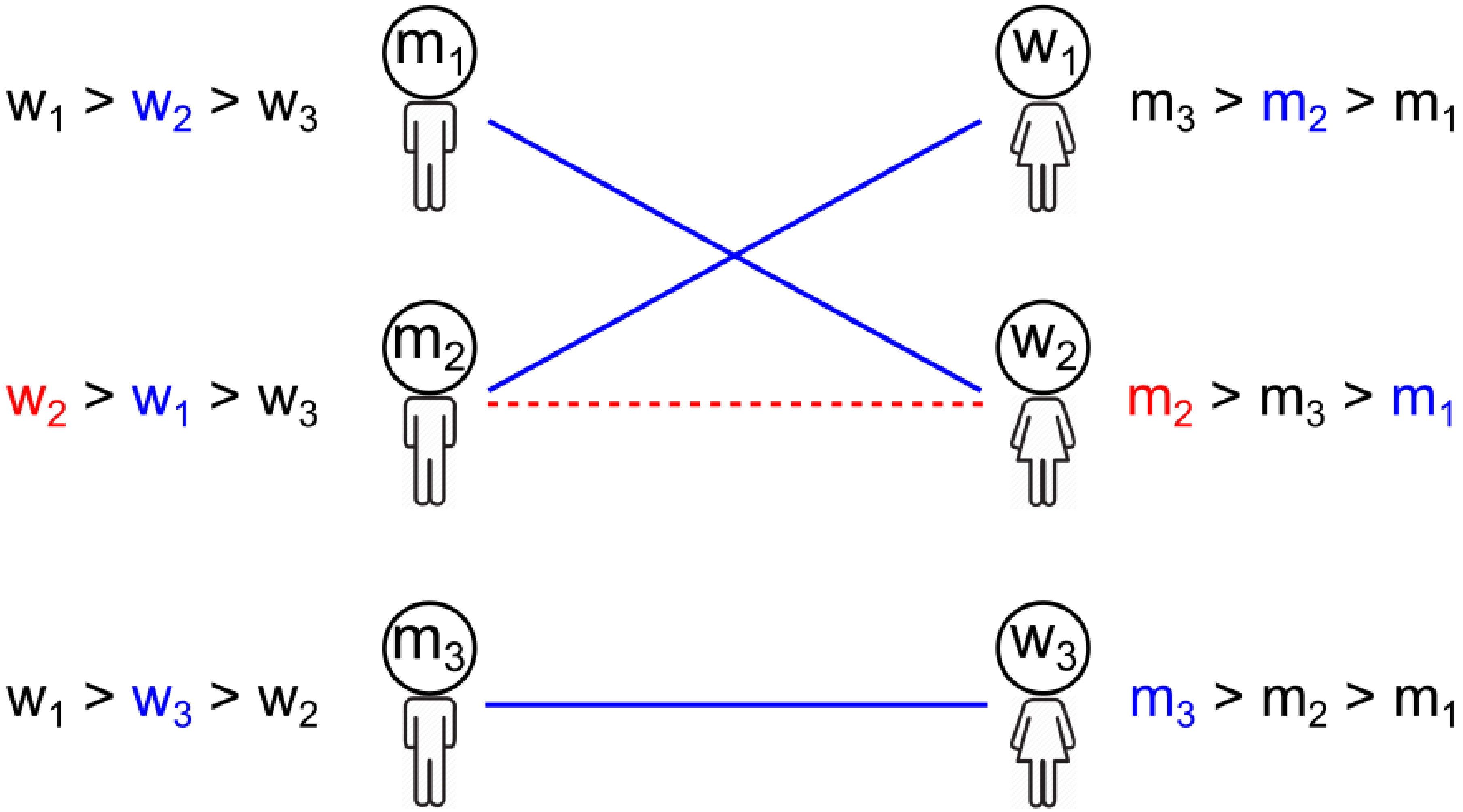
Stable Matching Problem



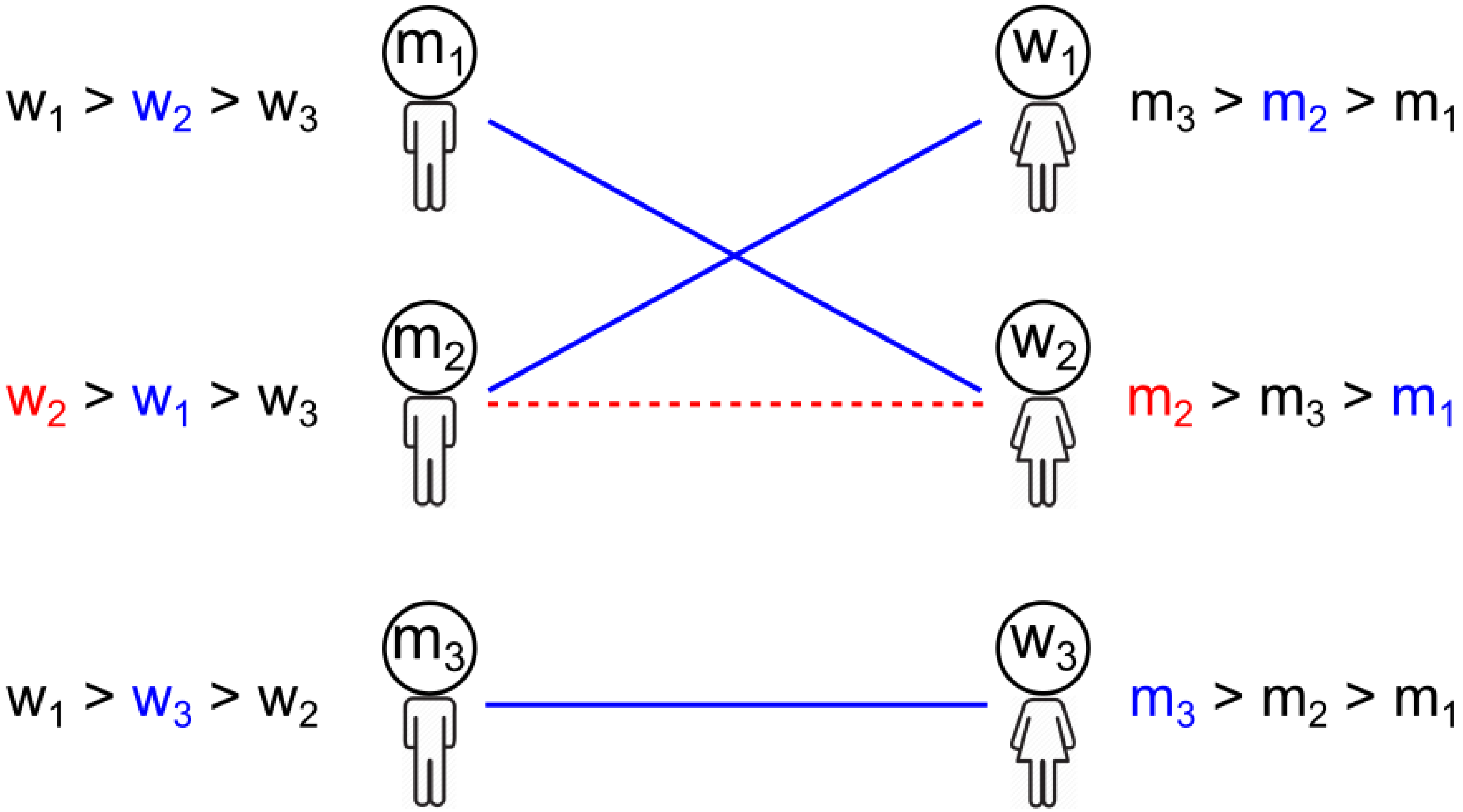
Stable Matching Problem



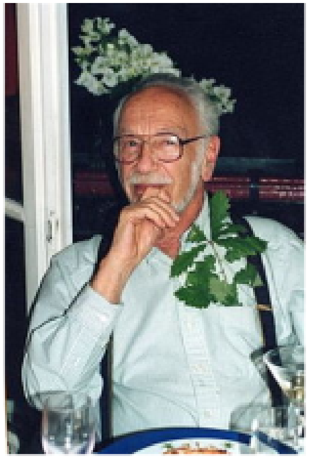
Stable Matching Problem



Stable Matching Problem



A matching is stable if there is no blocking pair.



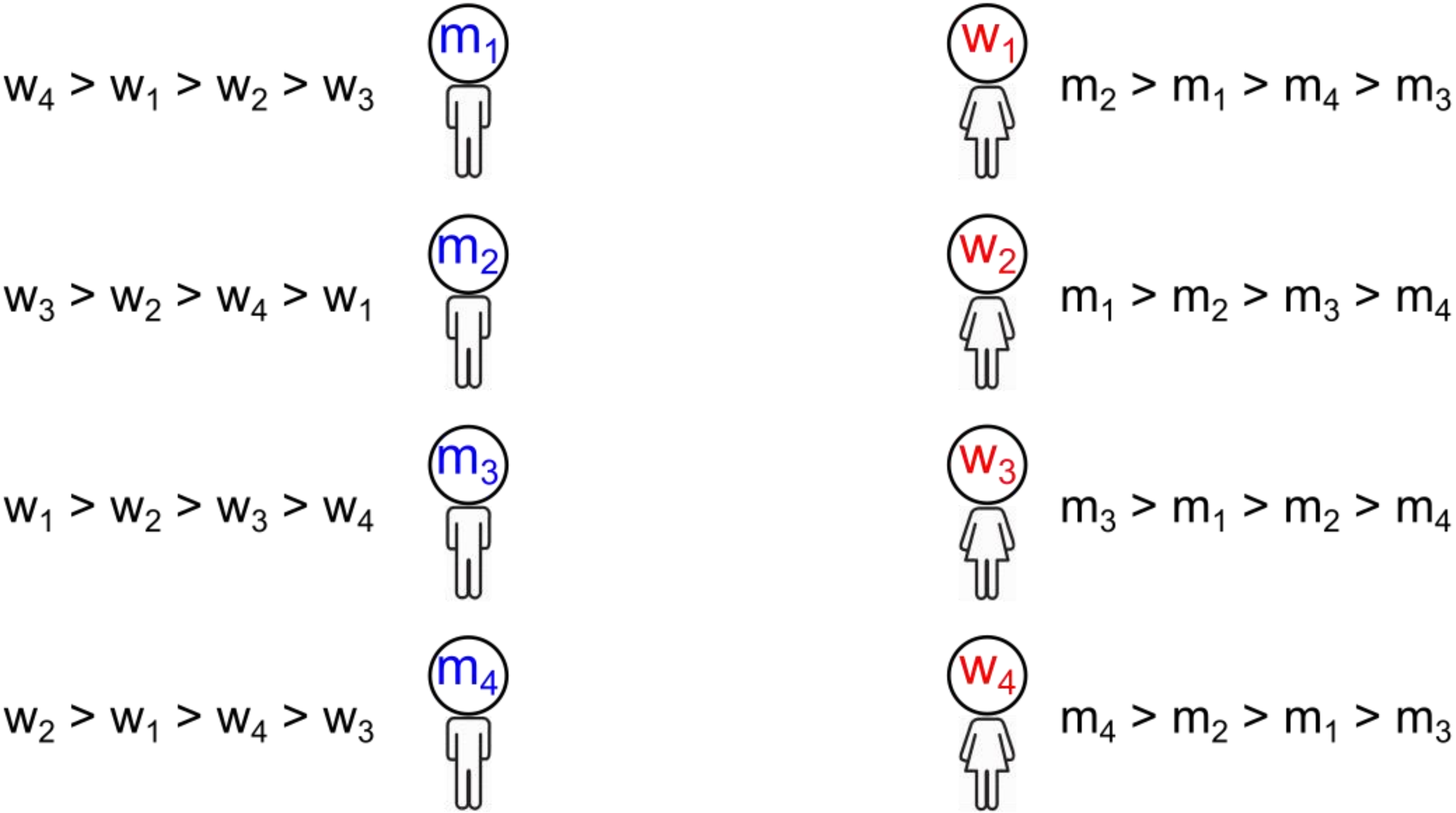
COLLEGE ADMISSIONS AND THE STABILITY OF MARRIAGE

D. GALE* AND L. S. SHAPLEY, Brown University and the RAND Corporation

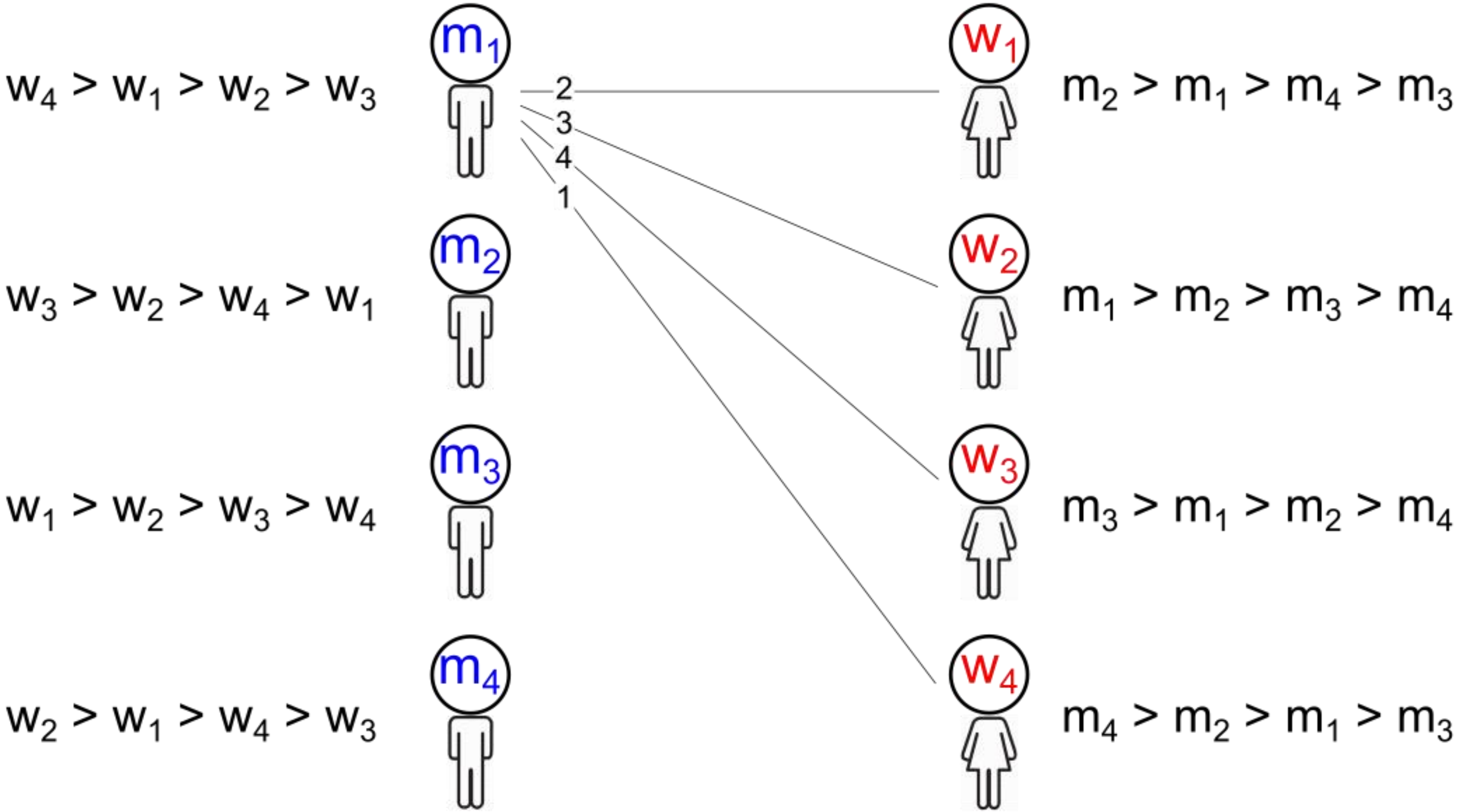
Source: *The American Mathematical Monthly*, Jan., 1962, Vol. 69, No. 1 (Jan., 1962), pp. 9–15

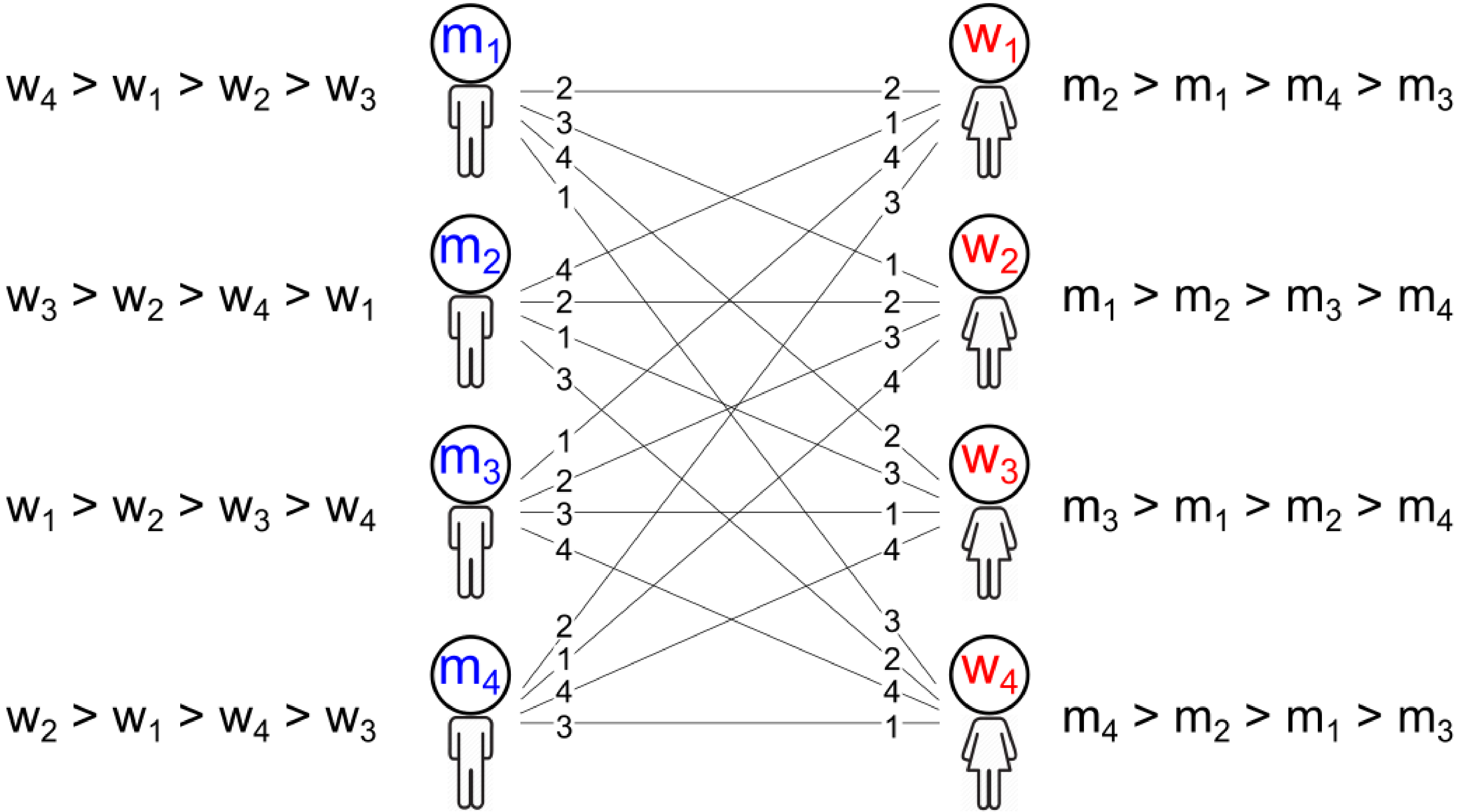
Given any preference profile, a stable matching for that profile always exists and can be computed in polynomial time.

Structure of the Set of Stable Matchings

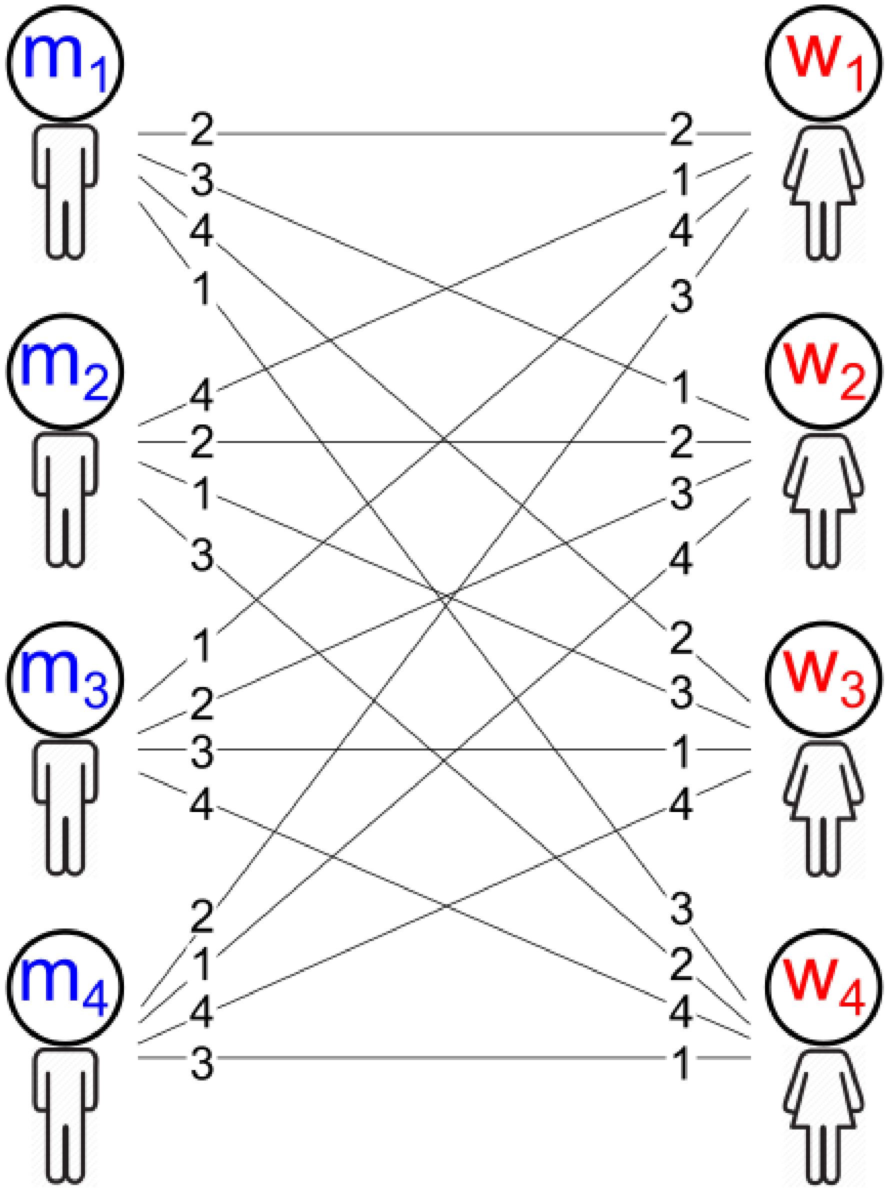


w4 > w1 > w2 > w3
m1
w1
m2 > m1 > m4 > m3
w3 > w2 > w4 > w1
m2
w2
m1 > m2 > m3 > m4
w1 > w2 > w3 > w4
m3
w3
m3 > m1 > m2 > m4
w2 > w1 > w4 > w3
m4
w4
m4 > m2 > m1 > m3

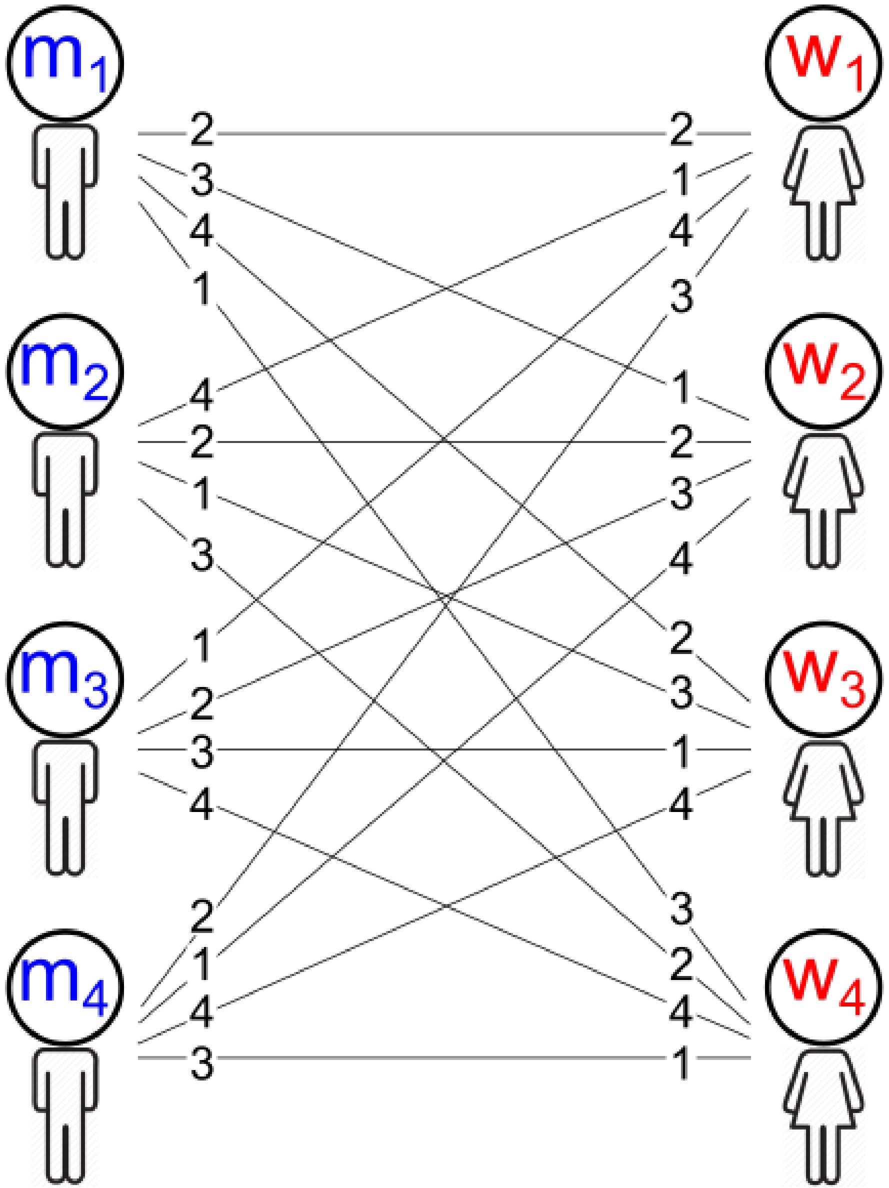

$w_4 > w_1 > w_2 > w_3$
m_1
$w_3 > w_2 > w_4 > w_1$
m_2
$w_1 > w_2 > w_3 > w_4$
m_3
$w_2 > w_1 > w_4 > w_3$
m_4
2
3
4
1
w_1
$m_2 > m_1 > m_4 > m_3$
w_2
$m_1 > m_2 > m_3 > m_4$
w_3
$m_3 > m_1 > m_2 > m_4$
w_4
$m_4 > m_2 > m_1 > m_3$

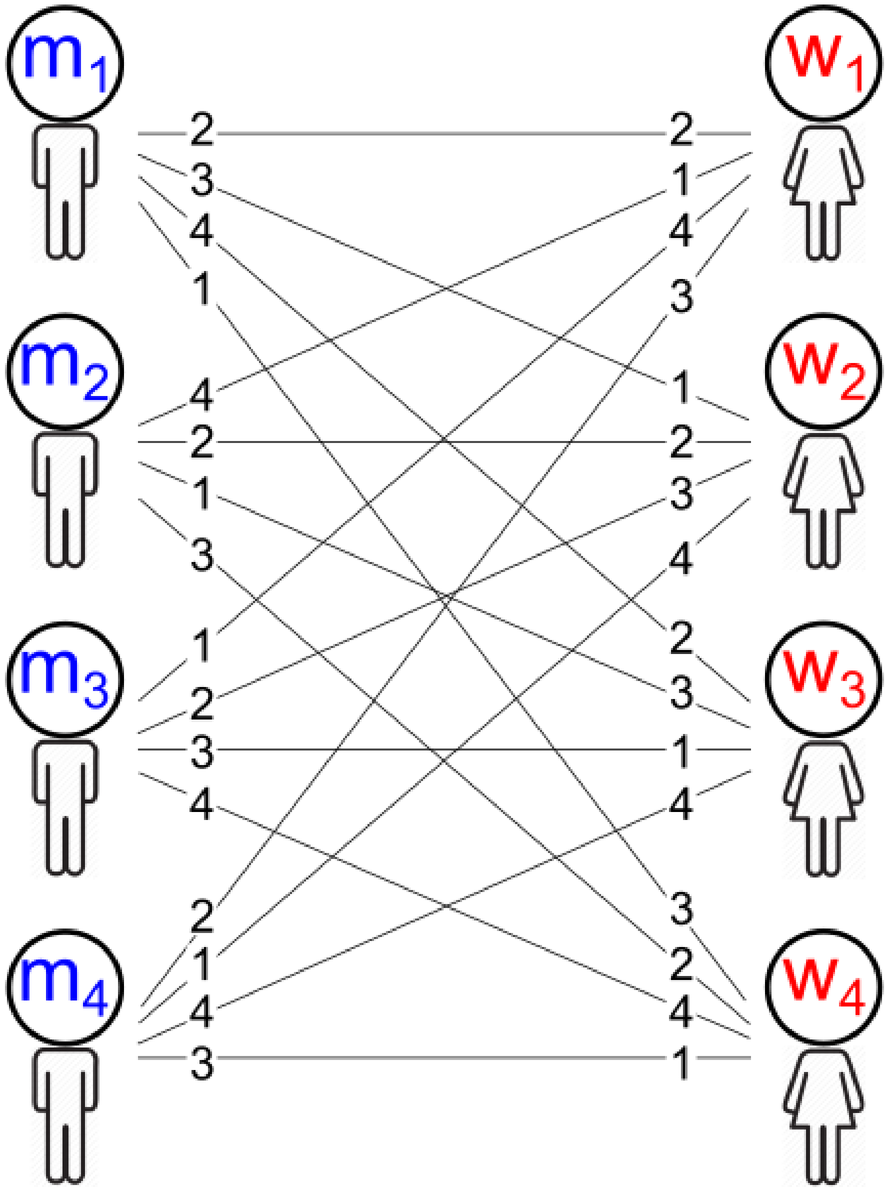


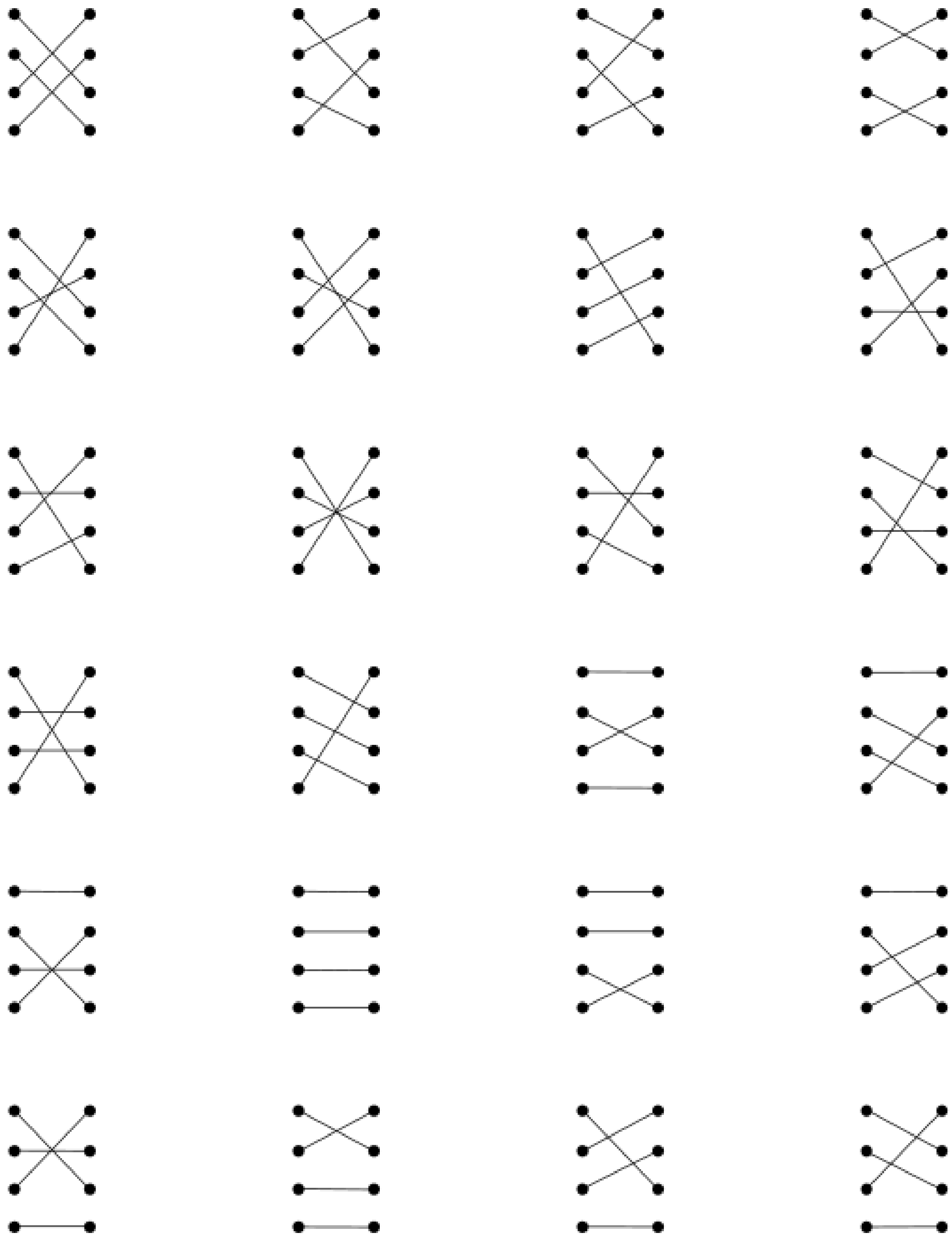
w4 > w1 > w2 > w3
w3 > w2 > w4 > w1
w1 > w2 > w3 > w4
w2 > w1 > w4 > w3
m1
m2
m3
m4
w1
w2
w3
w4
m2 > m1 > m4 > m3
m1 > m2 > m3 > m4
m3 > m1 > m2 > m4
m4 > m2 > m1 > m3

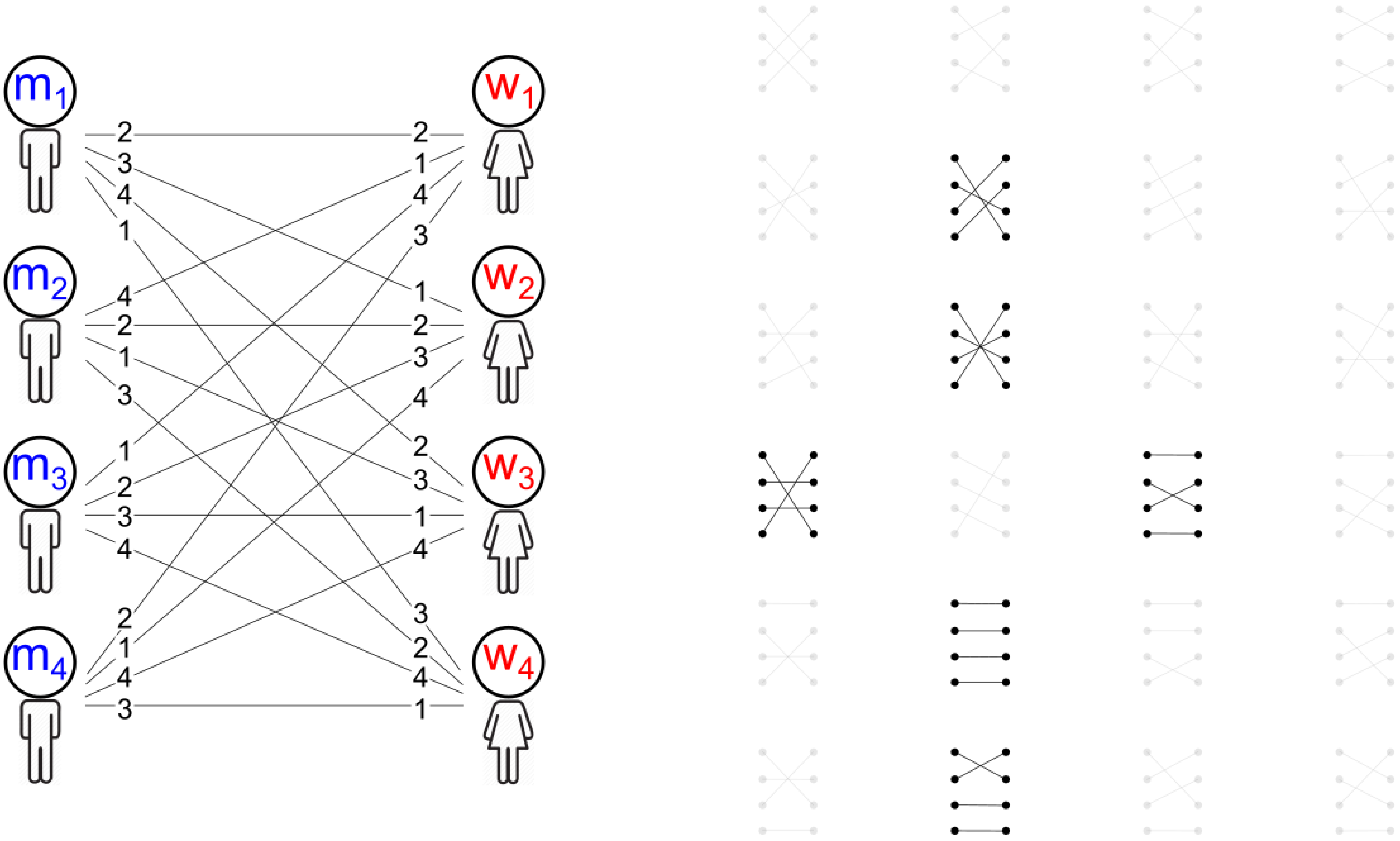


m1
2 3 4 1
m2
4 2 1 3
m3
1 2 3 4
m4
2 1 4 3
w1
2 1 4 3
w2
1 2 3 4
w3
2 3 1 4
w4
3 2 4 1

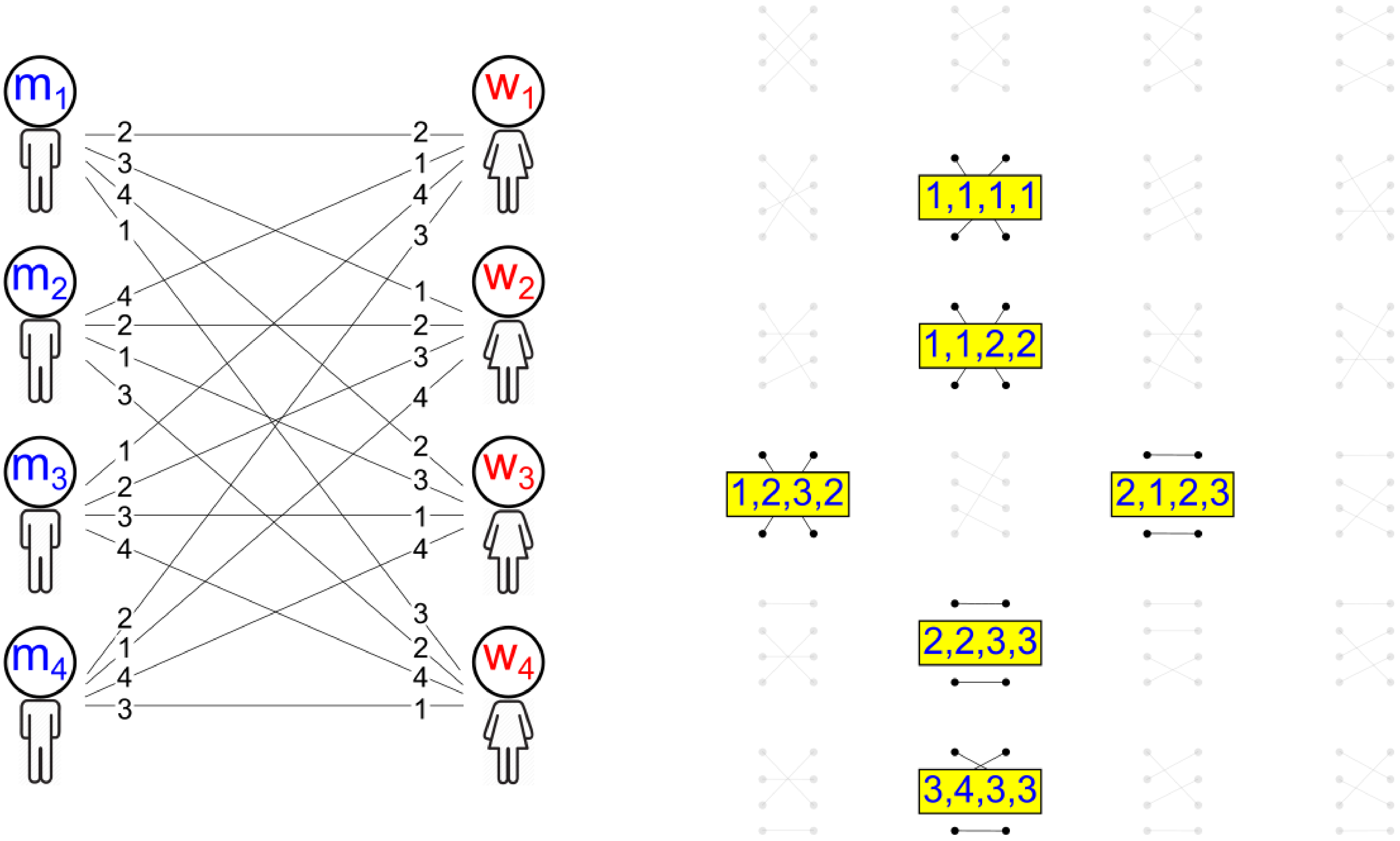

m1
m2
m3
m4
w1
w2
w3
w4
2 3 4 1
4 2 1 3
1 2 3 4
2 1 4 3
2 1 4 3
1 2 3 4
2 3 1 4
3 2 4 1


m1
m2
m3
m4
w1
w2
w3
w4

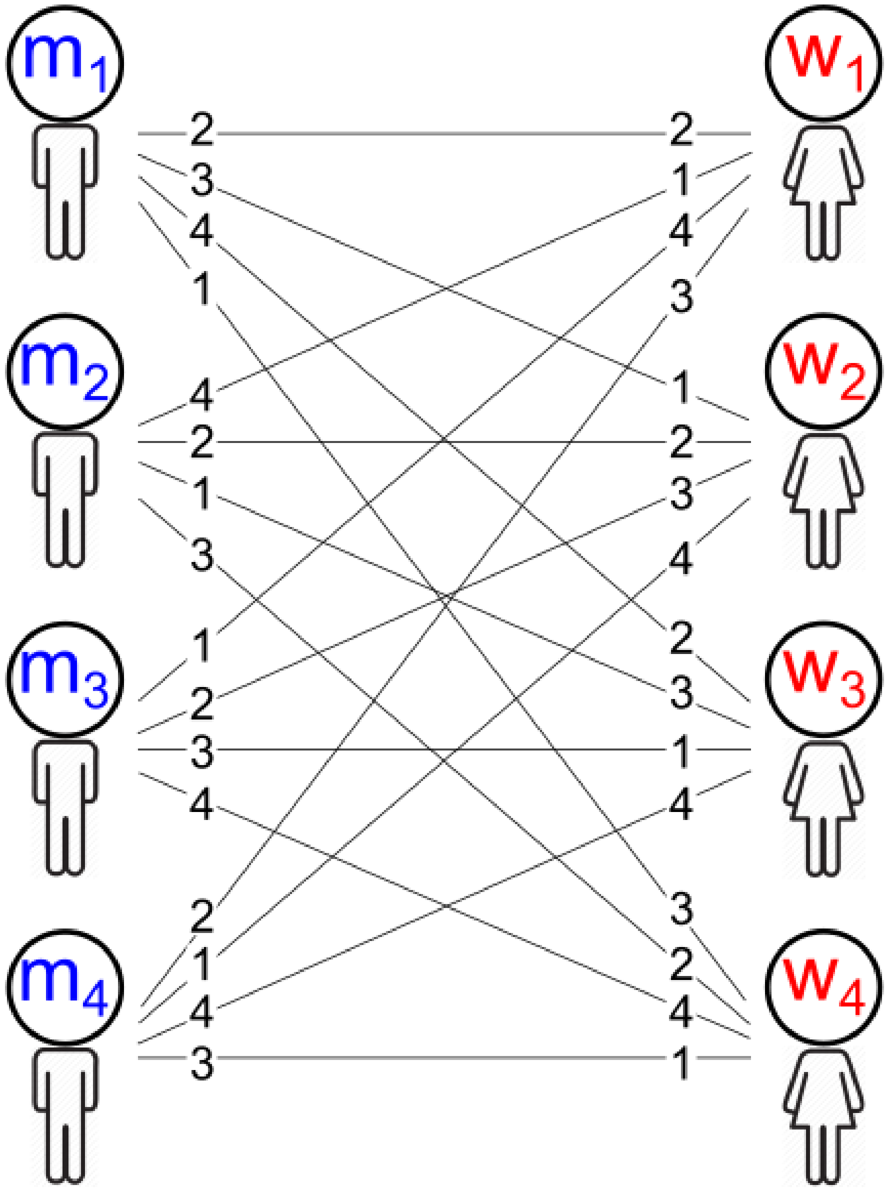




m1
m2
m3
m4
w1
w2
w3
w4



m1
m2
m3
m4
w1
w2
w3
w4
1,1,1,1
1,1,2,2
1,2,3,2
2,1,2,3
2,2,3,3
3,4,3,3



1,1,1,1

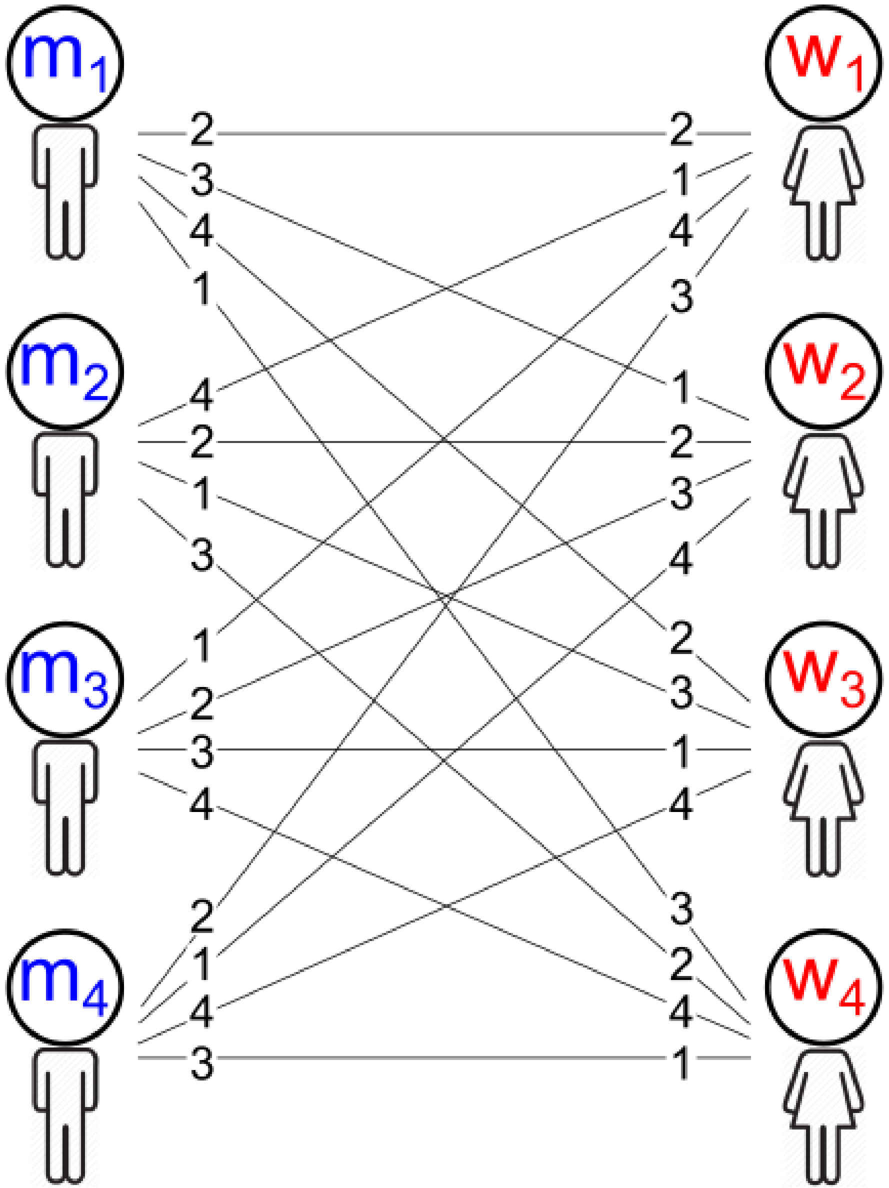
1,1,2,2

1,2,3,2

2,1,2,3

2,2,3,3

3,4,3,3



Men-optimal 1,1,1,1

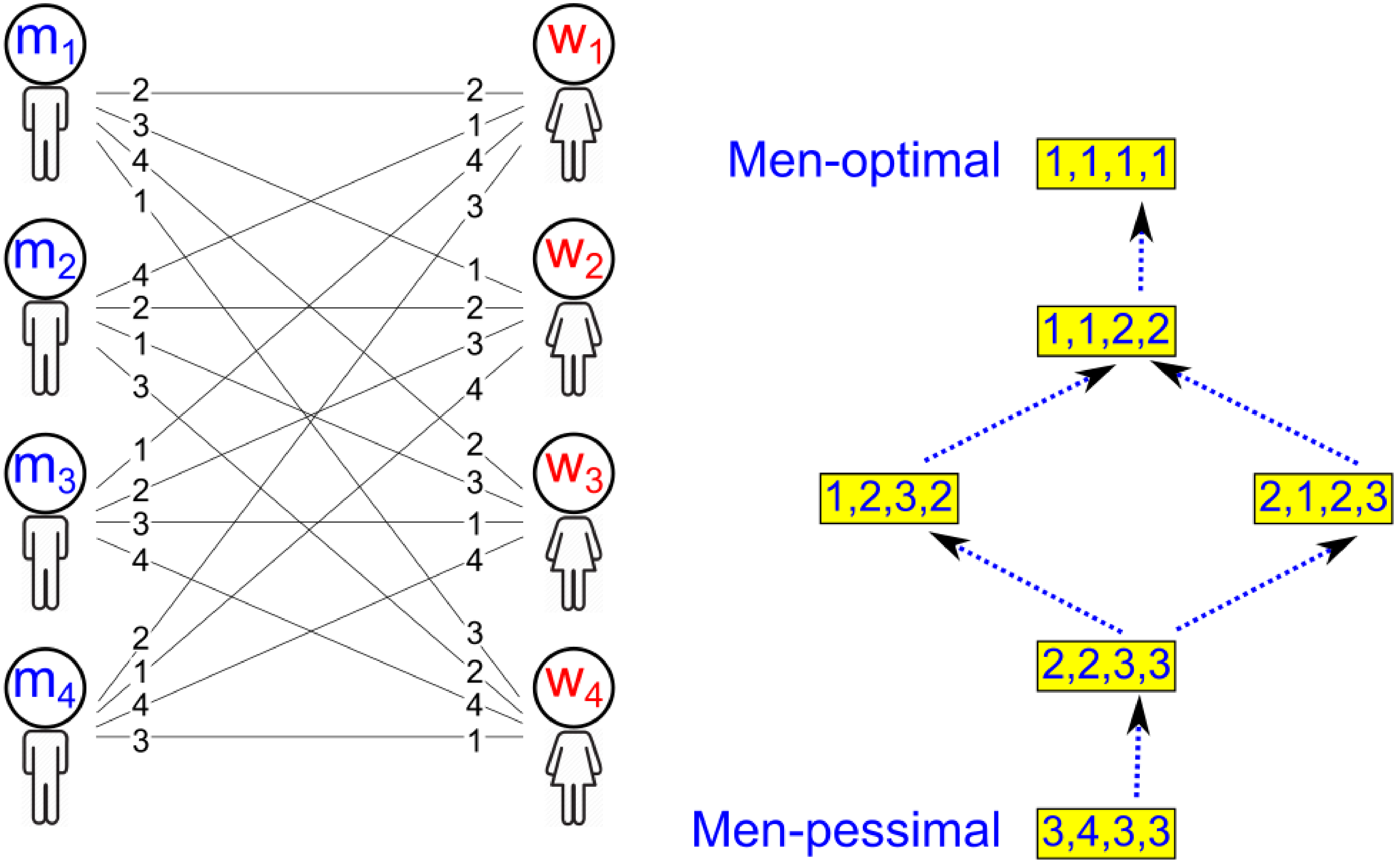
1,1,2,2

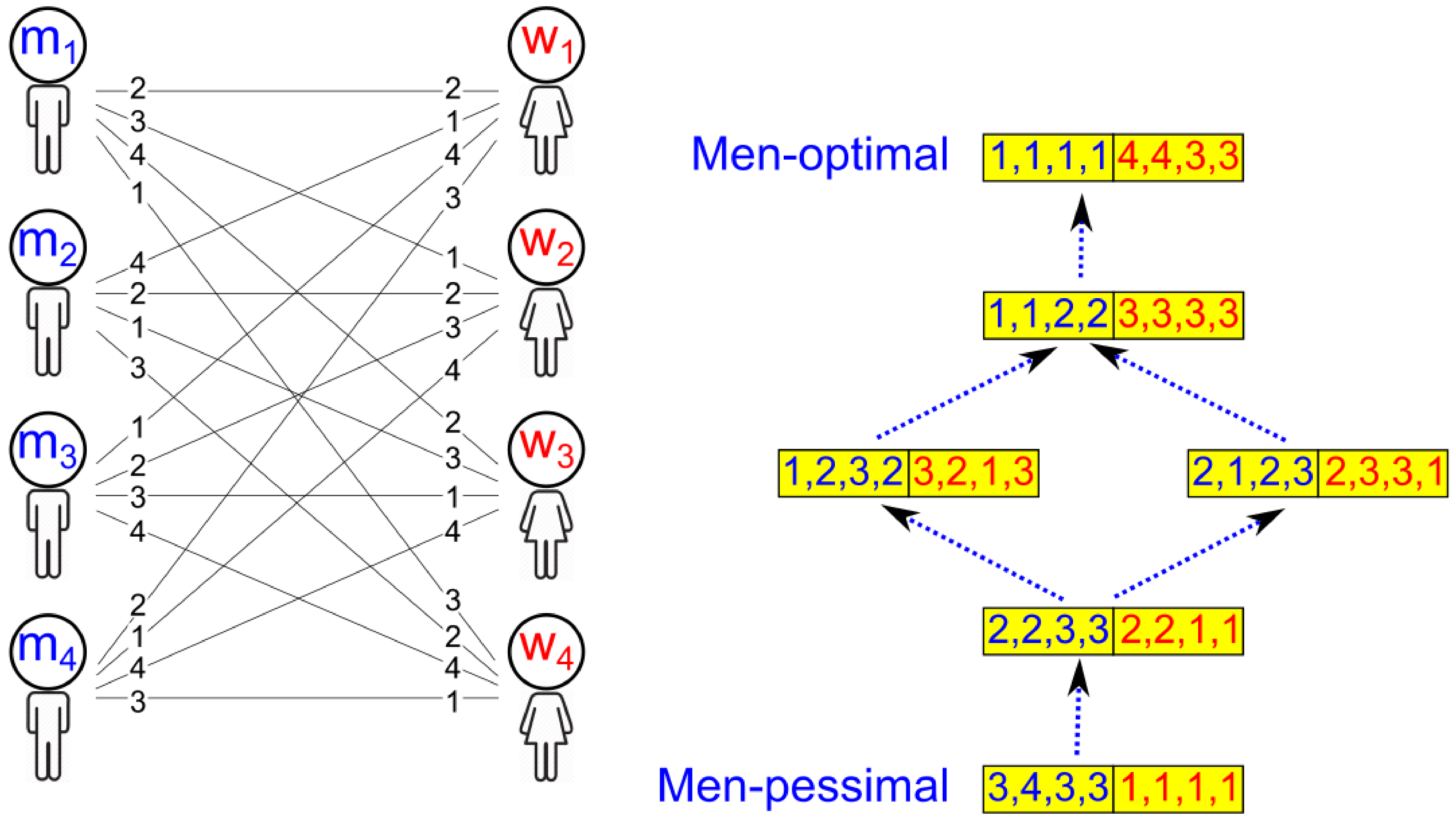
1,2,3,2

2,1,2,3

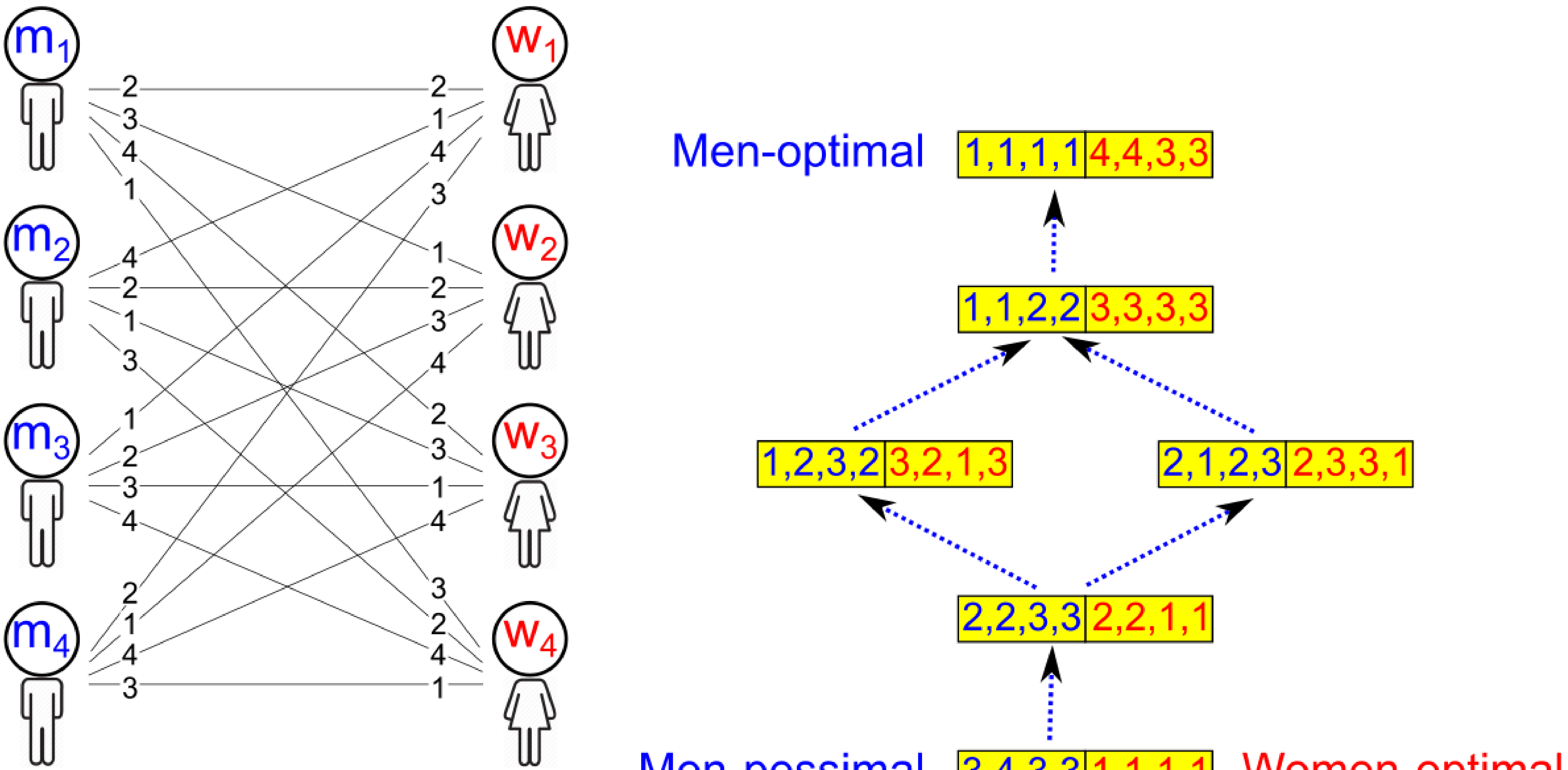
2,2,3,3

Men-pessimal 3,4,3,3

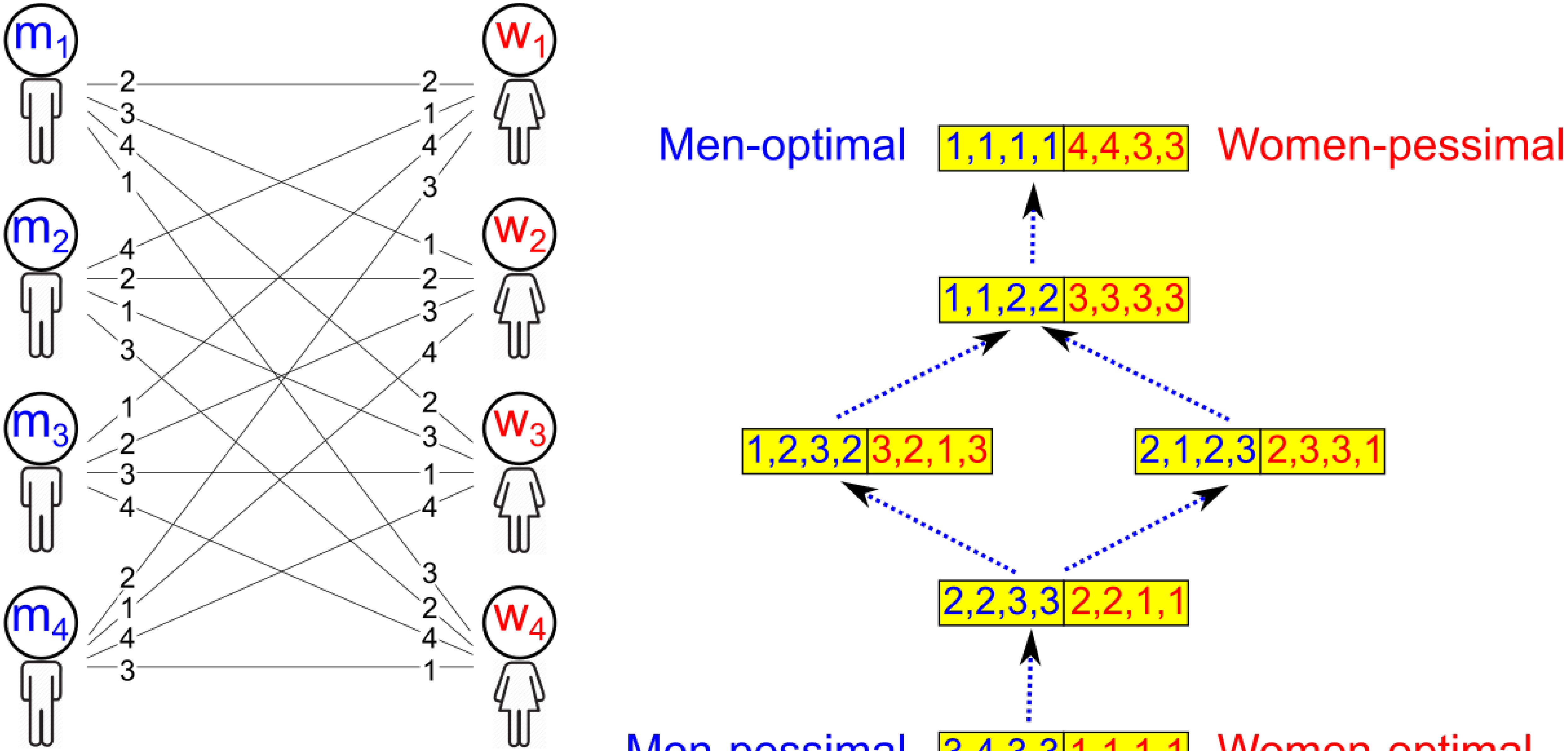

m1
2 3 4 1
m2
4 2 1 3
m3
1 2 3 4
m4
2 1 4 3
w1
2 1 4 3
w2
1 2 3 4
w3
2 3 1 4
w4
3 2 4 1
Men-optimal
1,1,1,1
1,1,2,2
1,2,3,2
2,1,2,3
2,2,3,3
Men-pessimal
3,4,3,3



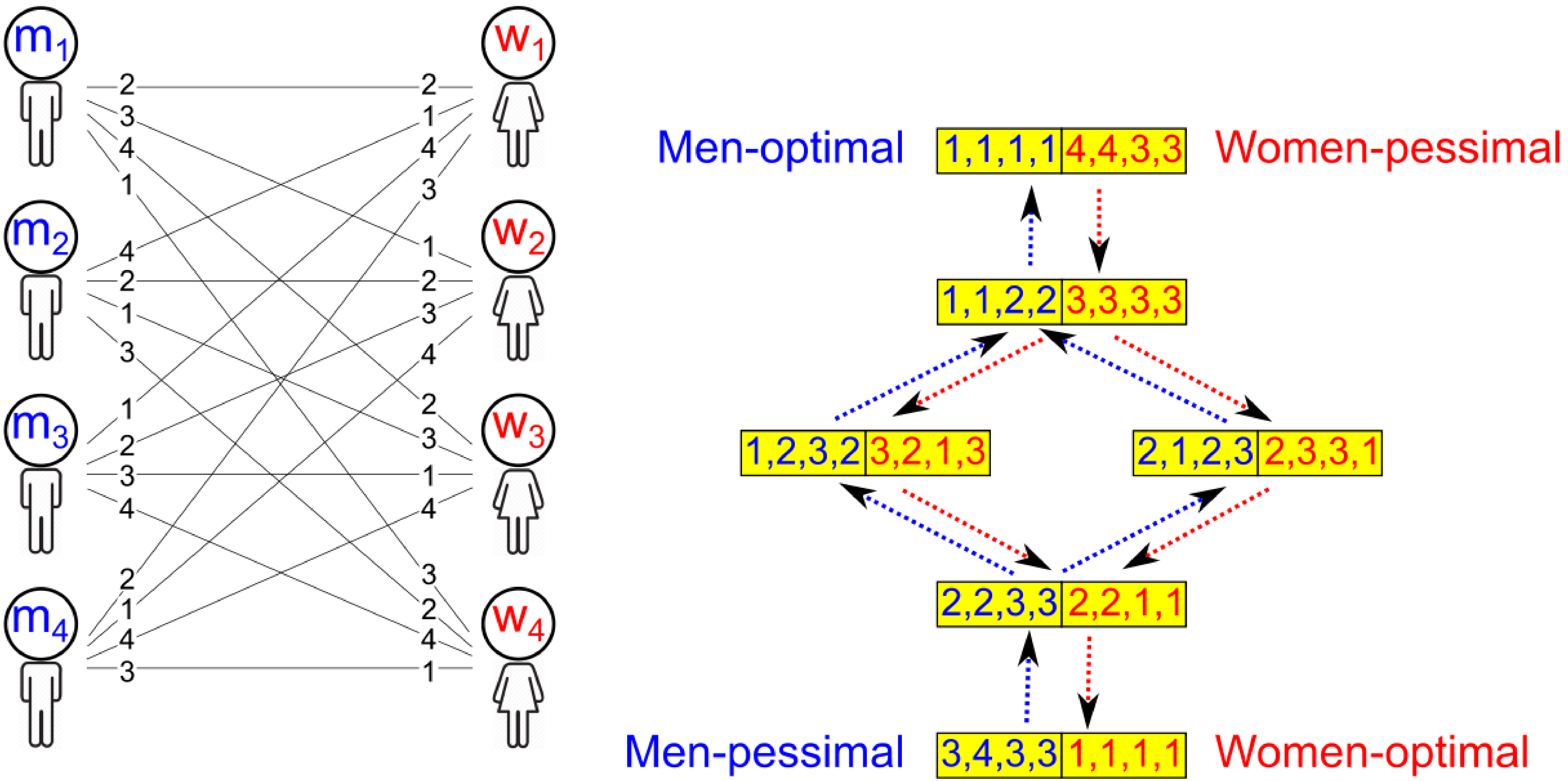
m1
m2
m3
m4
w1
w2
w3
w4
Men-optimal
1,1,1,1 4,4,3,3
1,1,2,2 3,3,3,3
1,2,3,2 3,2,1,3
2,1,2,3 2,3,3,1
2,2,3,3 2,2,1,1
Men-pessimal
3,4,3,3 1,1,1,1



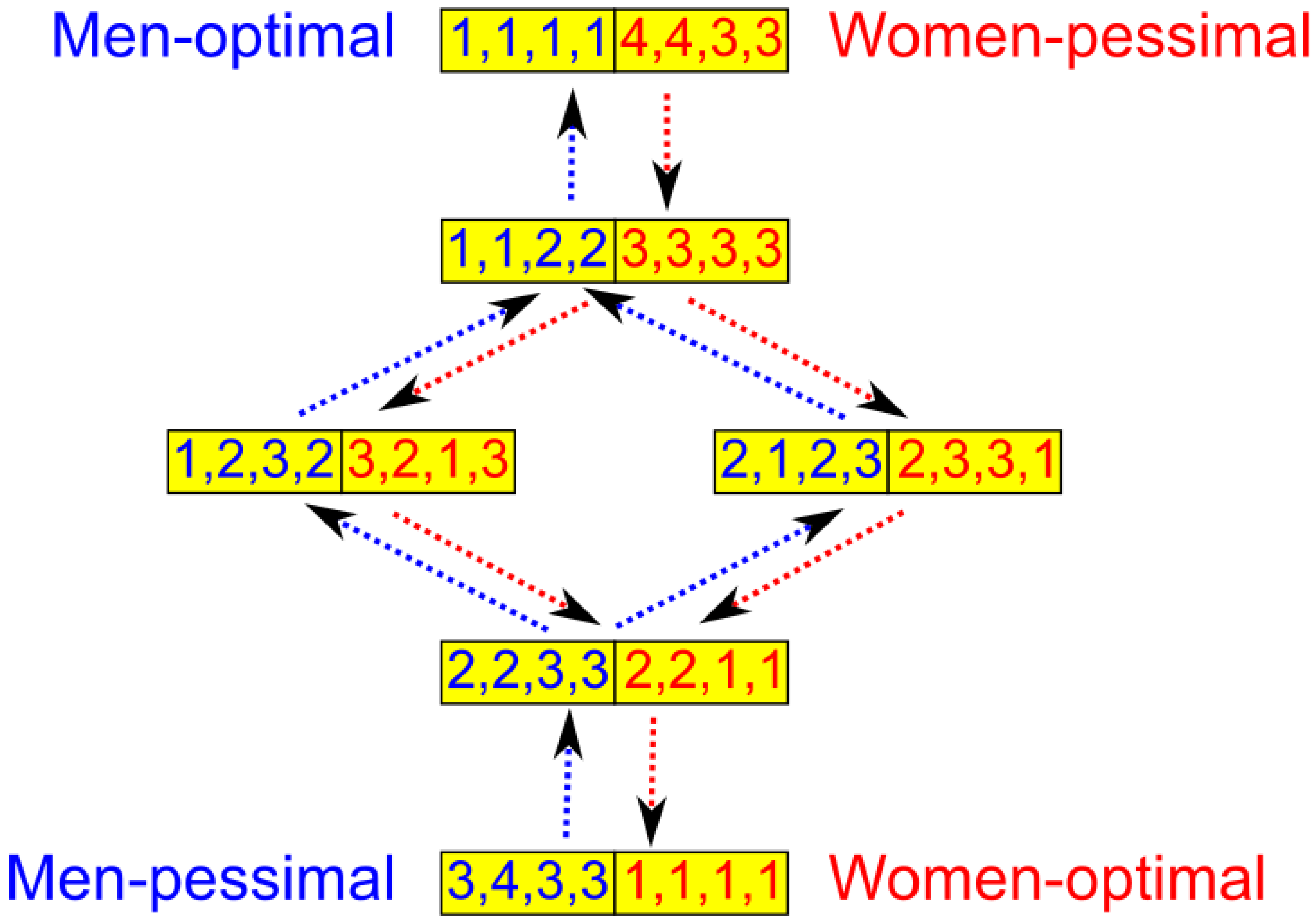
m1
m2
m3
m4
w1
w2
w3
w4
Men-optimal
1,1,1,1 4,4,3,3
1,1,2,2 3,3,3,3
1,2,3,2 3,2,1,3
2,1,2,3 2,3,3,1
2,2,3,3 2,2,1,1
Men-pessimal
3,4,3,3 1,1,1,1
Women-optimal



m_1
m_2
m_3
m_4
w_1
w_2
w_3
w_4
2 3 4 1
4 2 1 3
1 2 3 4
2 1 4 3
2 1 4 3
1 2 3 4
2 3 1 4
3 2 4 1
Men-optimal
1,1,1,1 4,4,3,3
Women-pessimal
1,1,2,2 3,3,3,3
1,2,3,2 3,2,1,3
2,1,2,3 2,3,3,1
2,2,3,3 2,2,1,1
Men-pessimal
3,4,3,3 1,1,1,1
Women-optimal

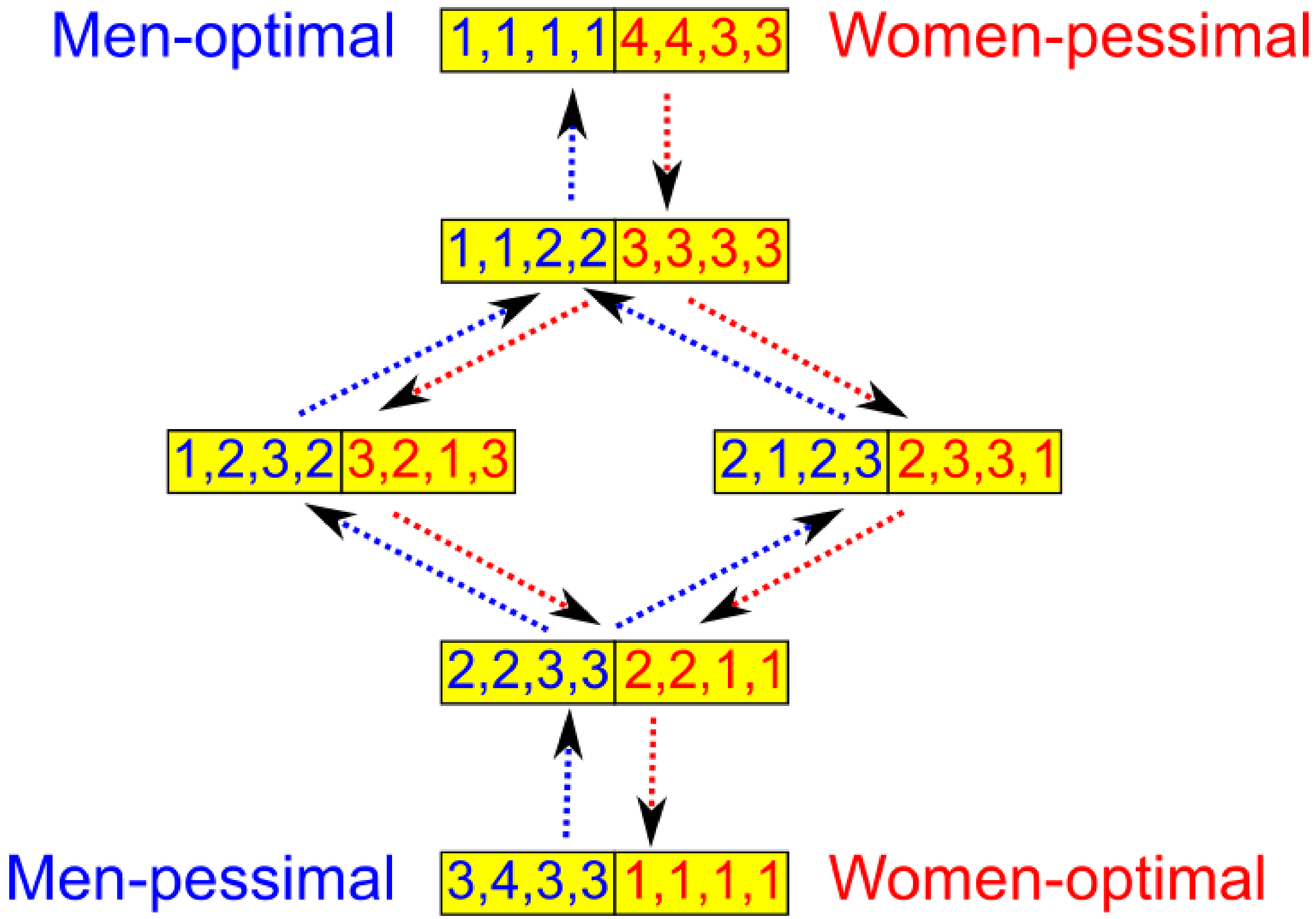


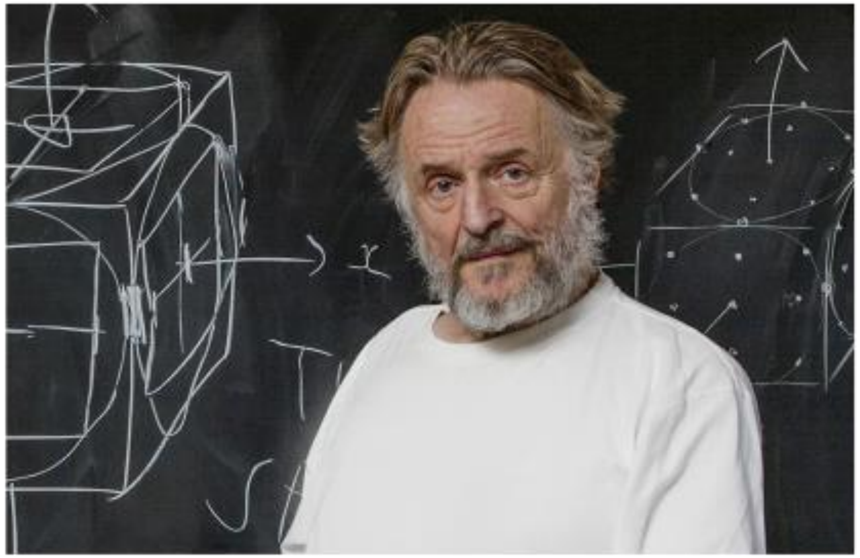
m1
m2
m3
m4
w1
w2
w3
w4
Men-optimal
1,1,1,1 4,4,3,3
Women-pessimal
1,1,2,2 3,3,3,3
1,2,3,2 3,2,1,3
2,1,2,3 2,3,3,1
2,2,3,3 2,2,1,1
Men-pessimal
3,4,3,3 1,1,1,1
Women-optimal



Men-optimal
1,1,1,1
4,4,3,3
Women-pessimal
1,1,2,2
3,3,3,3
1,2,3,2
3,2,1,3
2,1,2,3
2,3,3,1
2,2,3,3
2,2,1,1
Men-pessimal
3,4,3,3
1,1,1,1
Women-optimal

The Lattice of Stable Matchings





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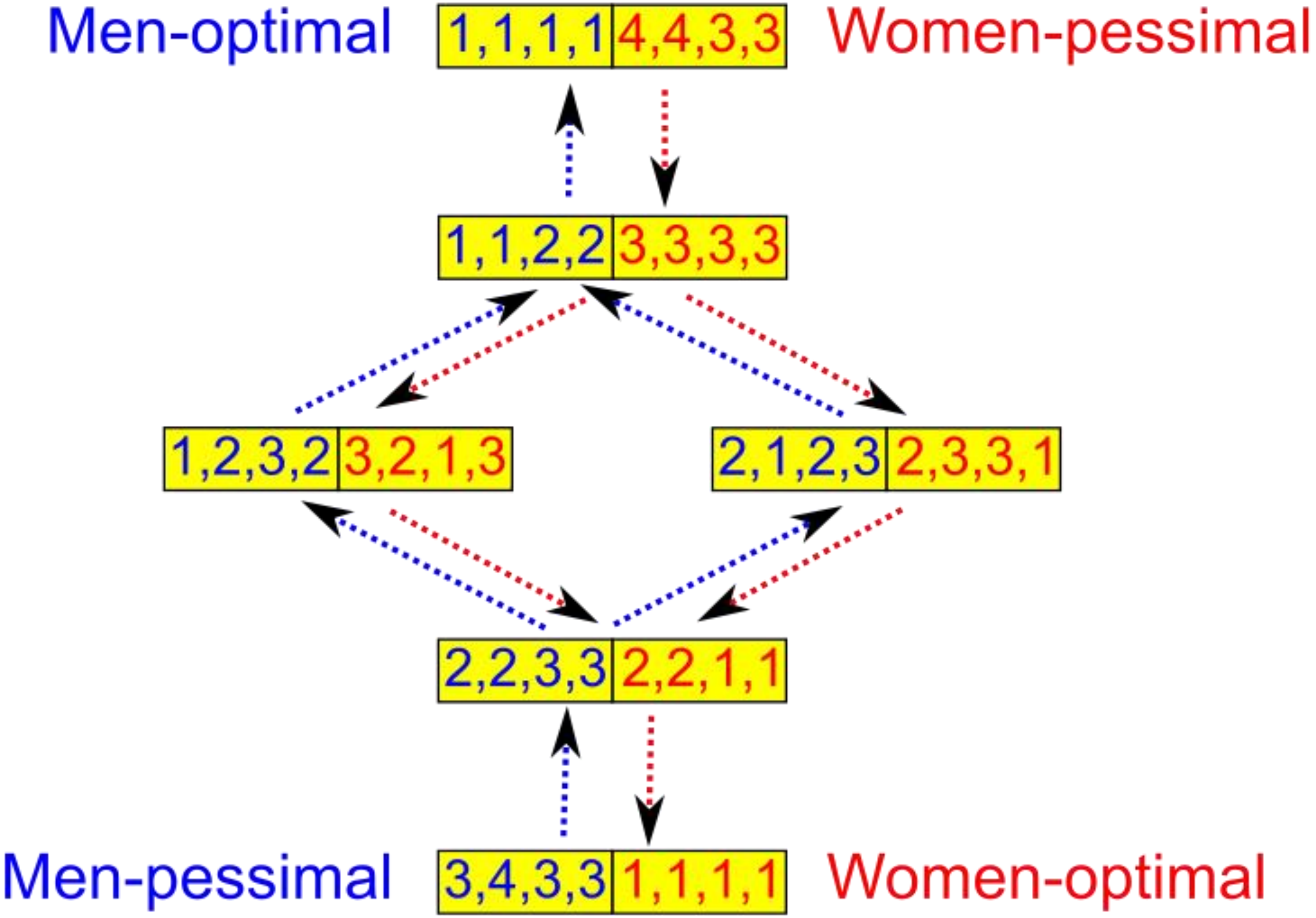
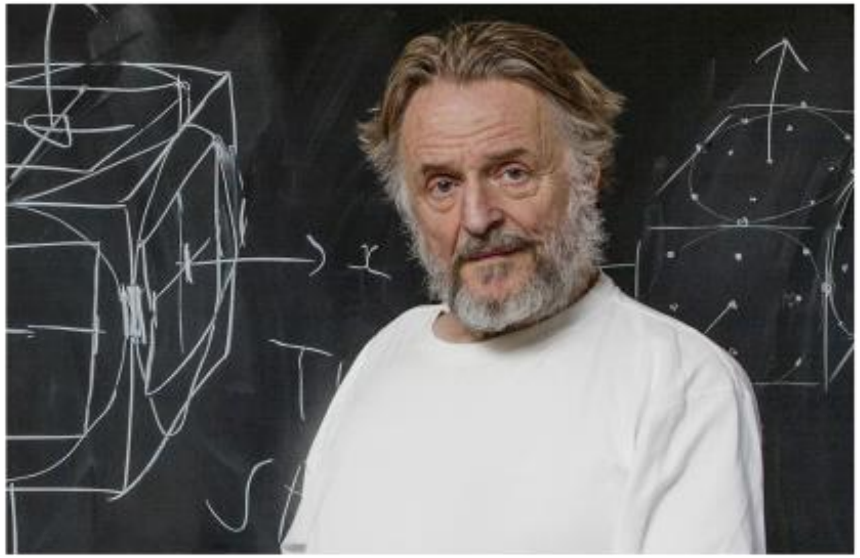
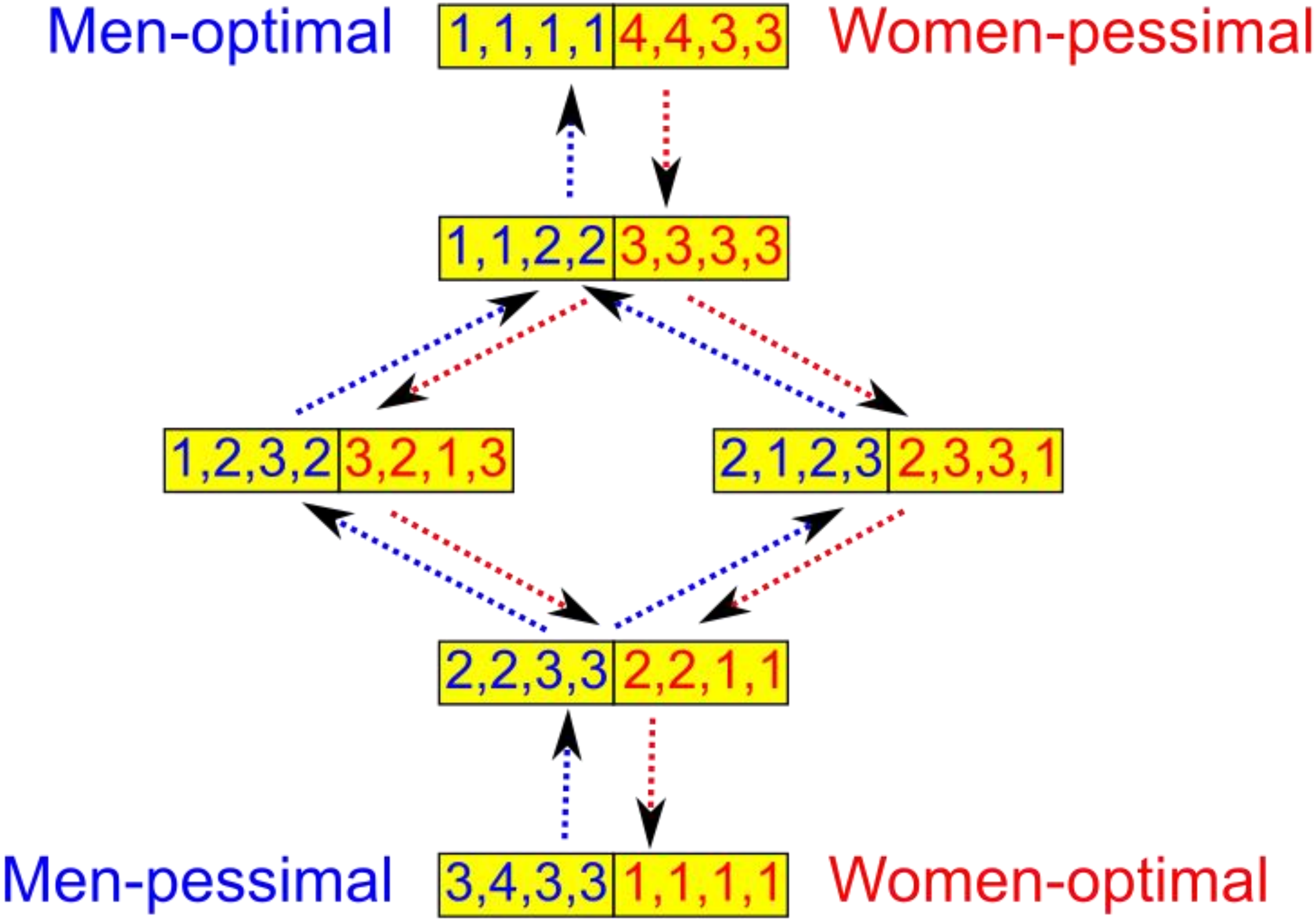




John H. Conway

The Lattice of Stable Matchings



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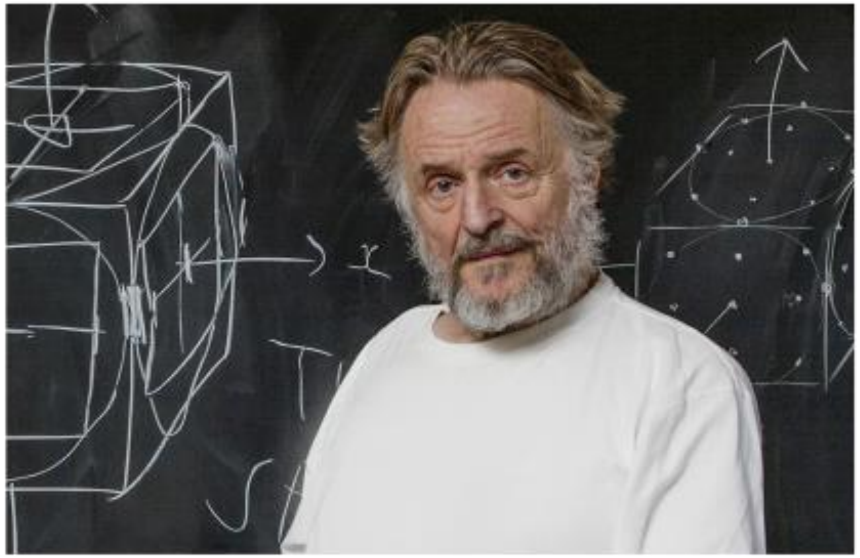
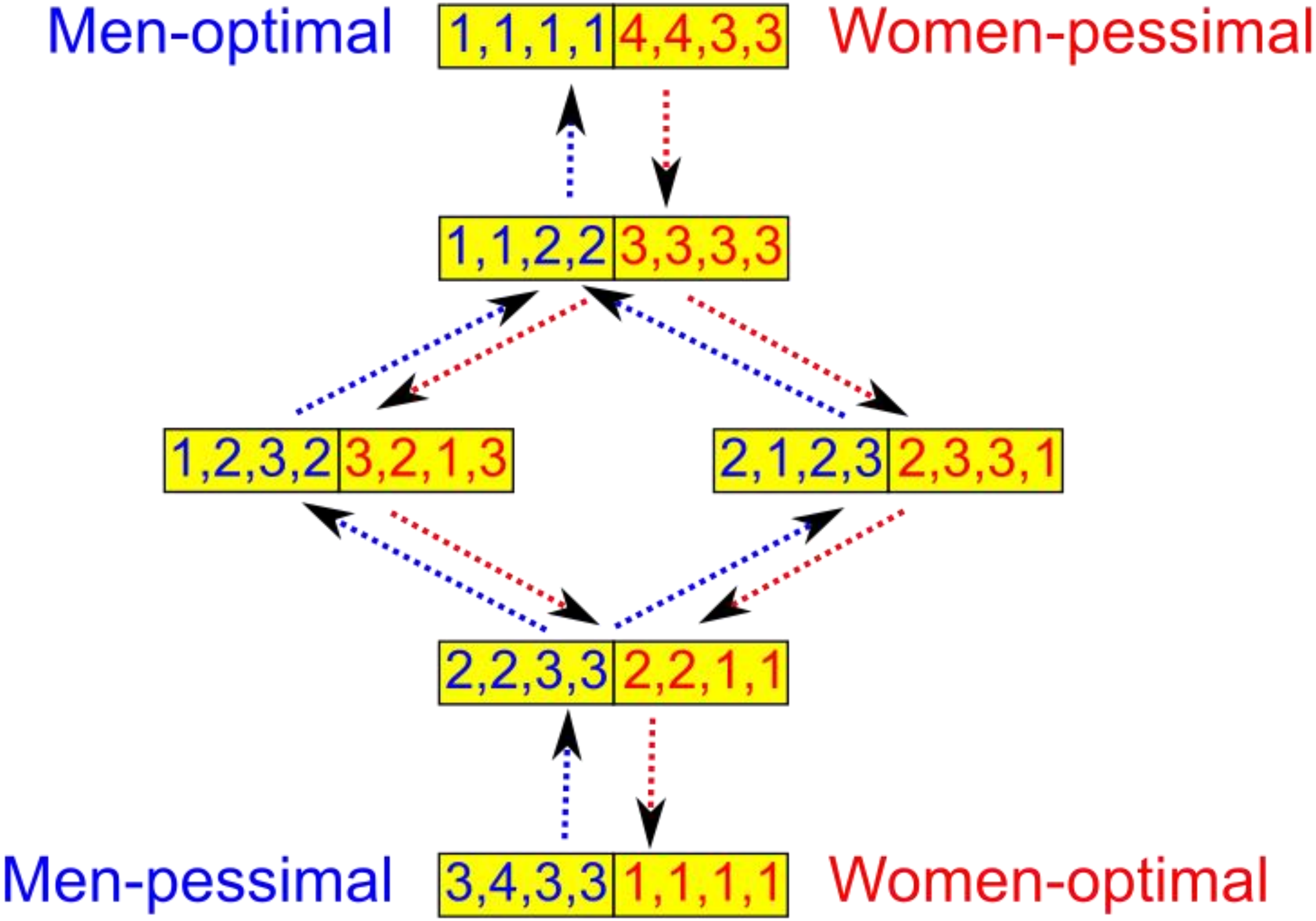


John H. Conway



Some Observations

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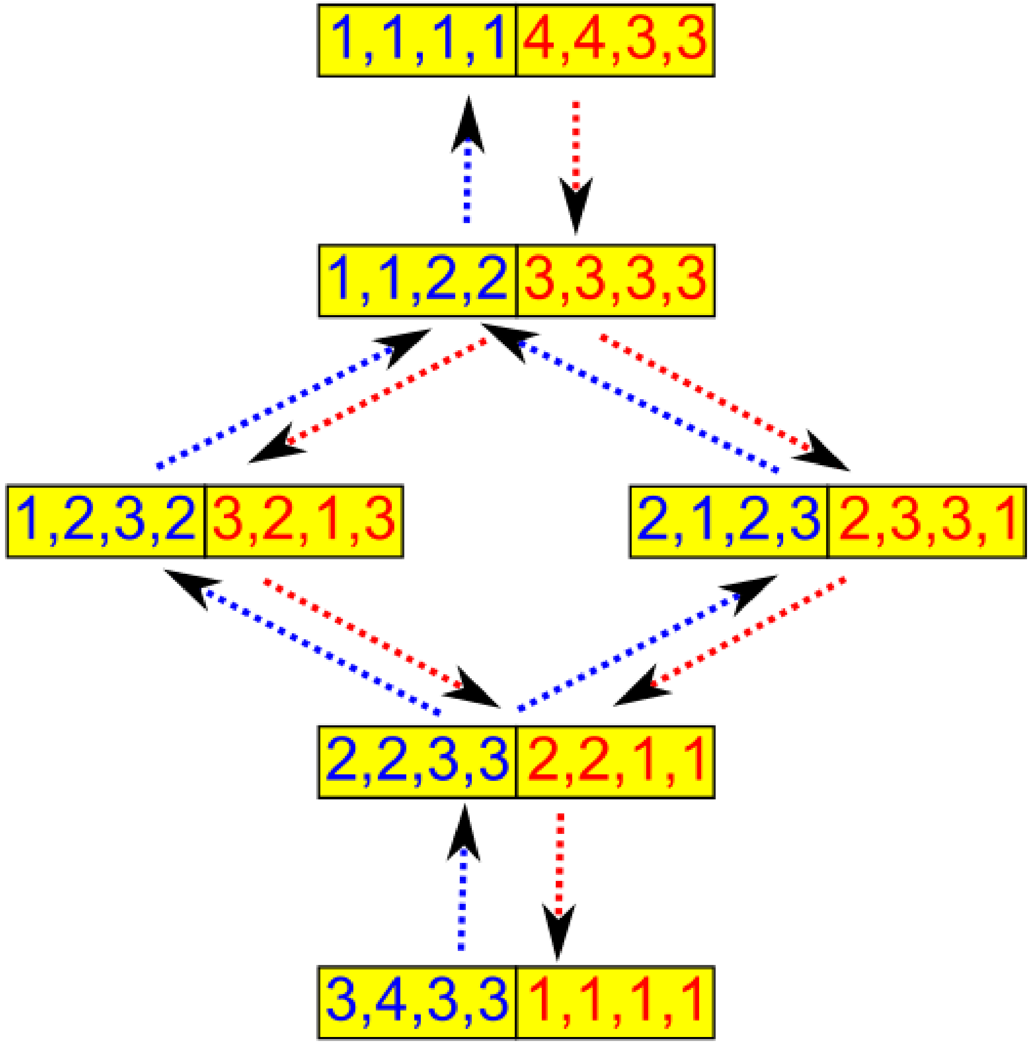
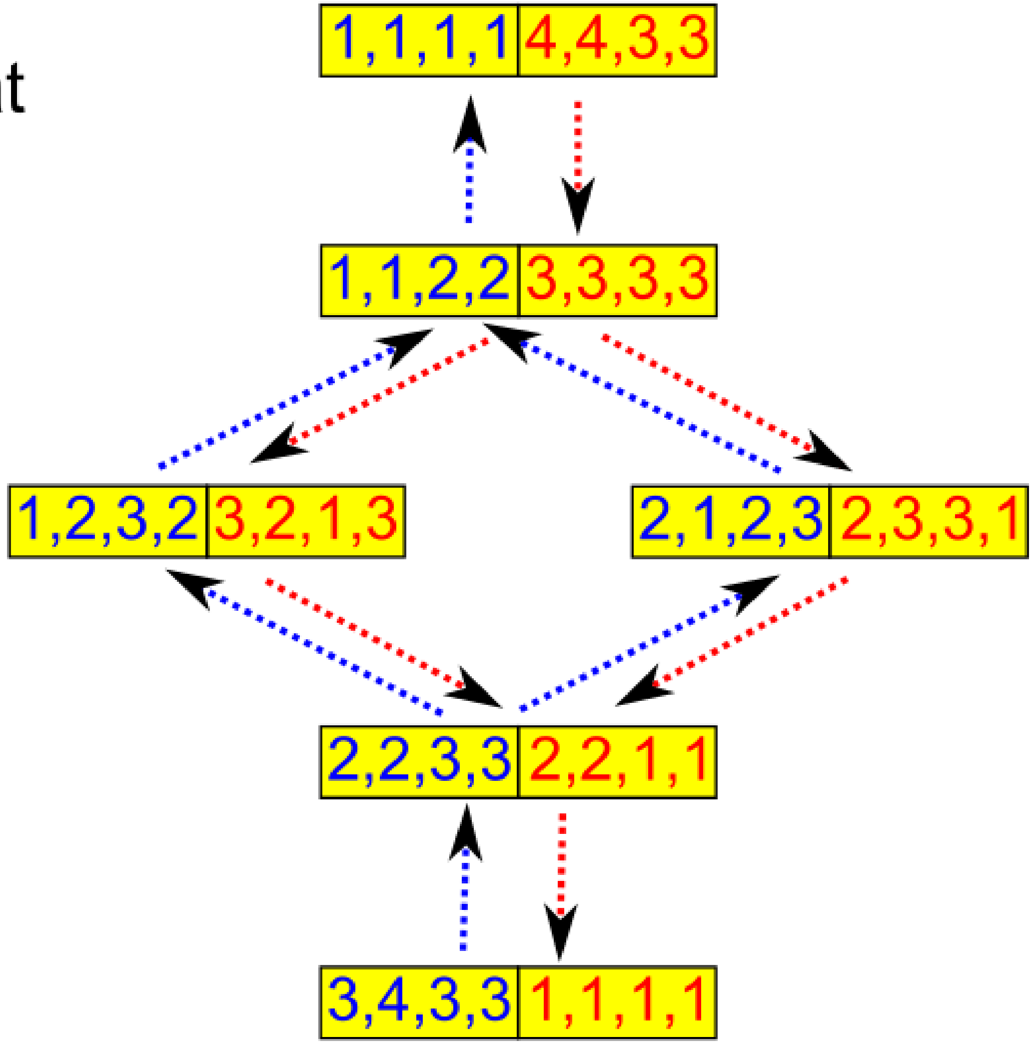


Some Observations

Consensus

There is a stable matching that all men find at least as good as *any other* stable matching, and one that they find at least as bad. (Analogously for the women.)



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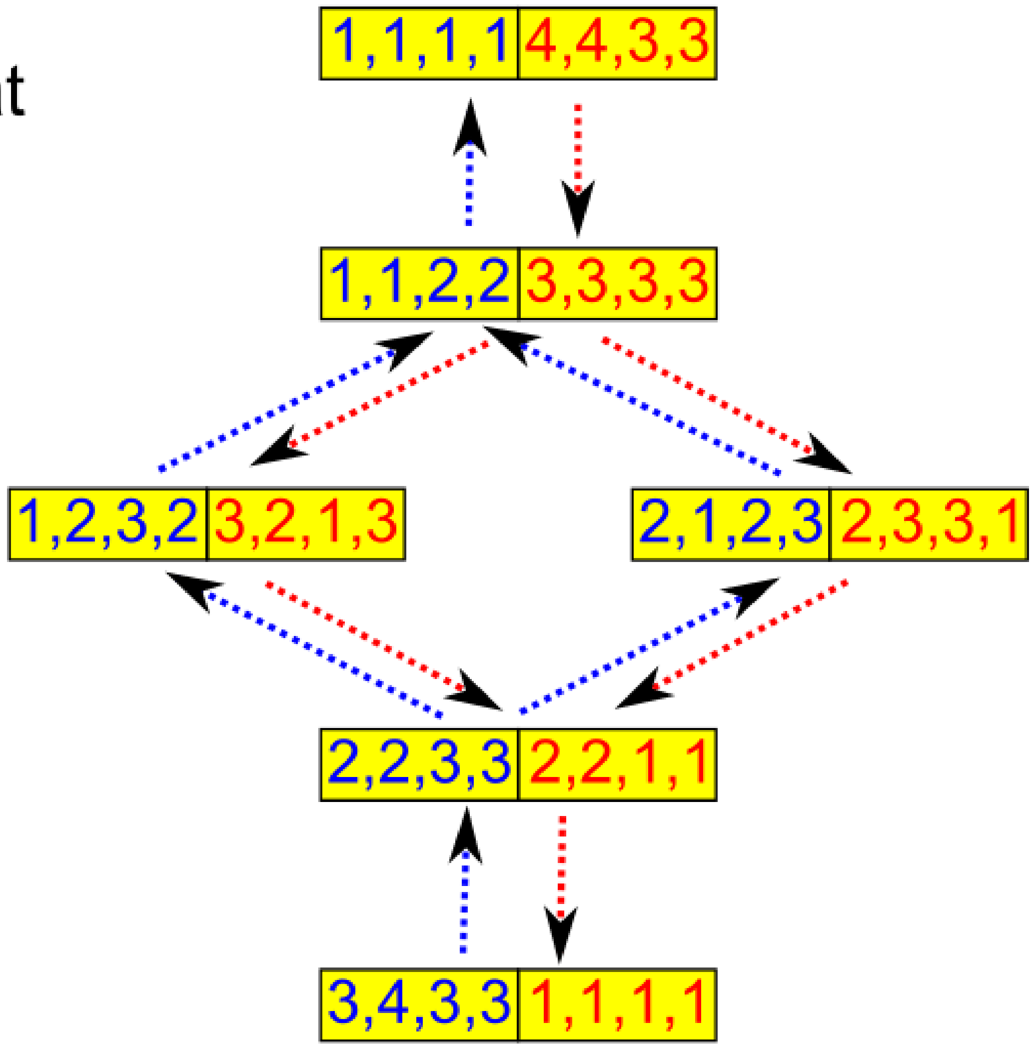
Conflict

For any distinct stable matchings P and Q, if all men find P at least as good as Q, then all women find Q at least as good as P (and vice versa).



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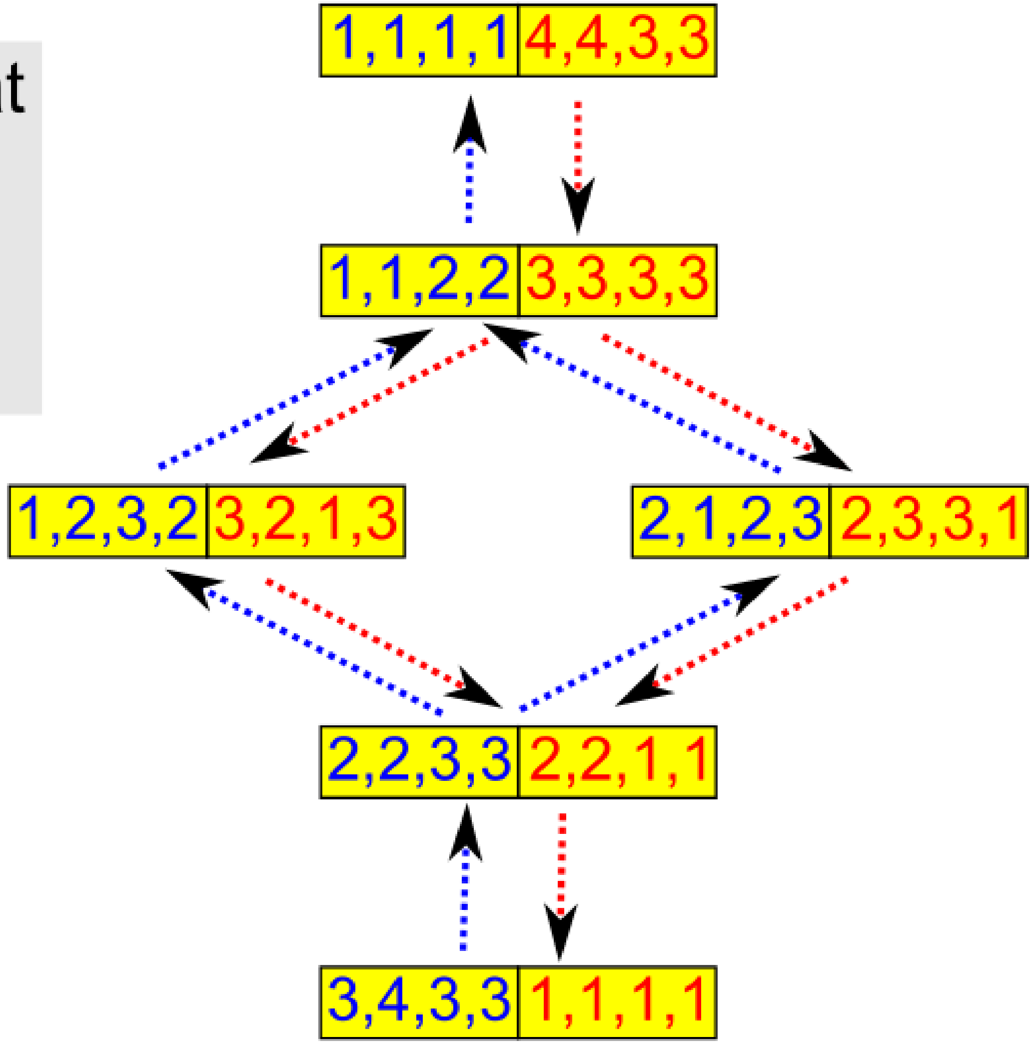

Call a man m and a woman w **achievable** for each other if there is some stable matching in which they are matched with each other.

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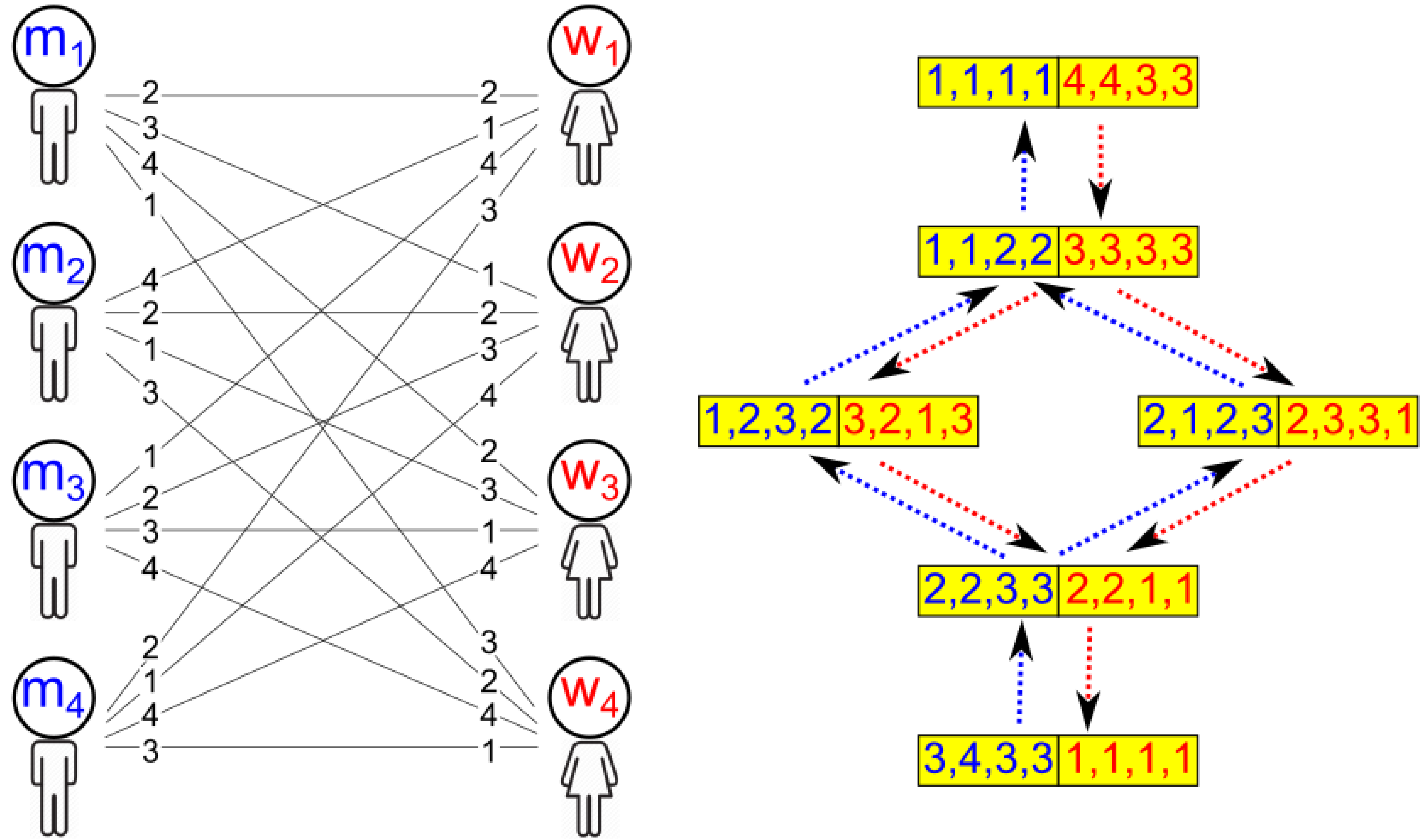

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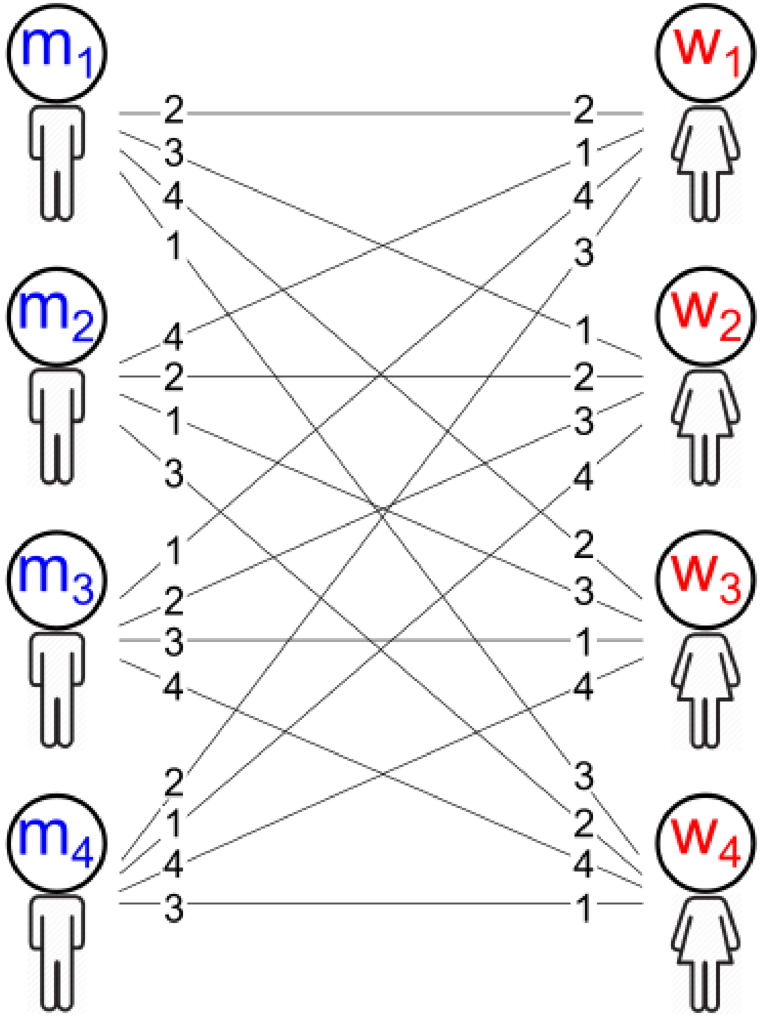

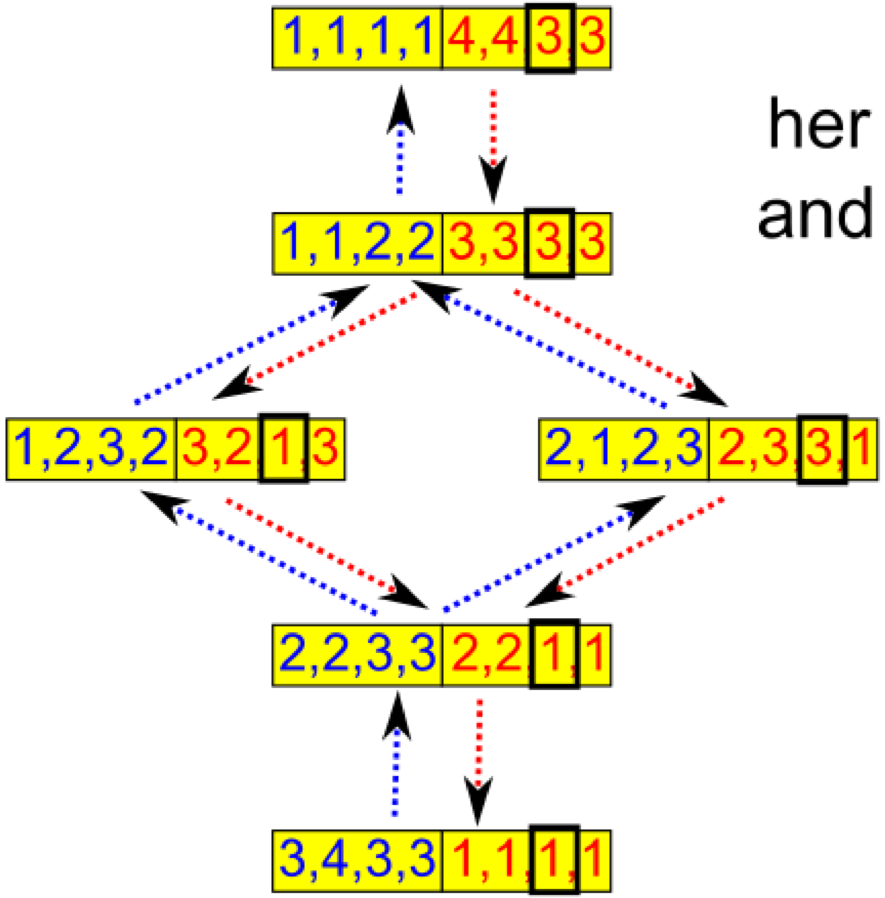

For w_3,
her first choice (man m_3)
and third choice (man m_2)
are **achievable**

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Consider a mapping in which each man is mapped to his favorite achievable woman (**men-optimal**), and another mapping in which each woman is mapped to her favorite achievable man (**women-optimal**).

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Consider a mapping in which each man is mapped to his favorite achievable woman (**men-optimal**), and another mapping in which each woman is mapped to her favorite achievable man (**women-optimal**).

We will show that men/women-optimal mappings are actually *matchings*.

Men-optimal mapping is a matching.

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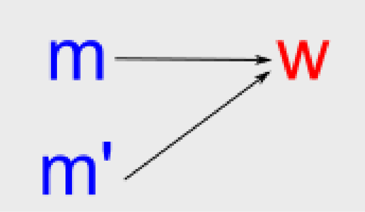
Suppose not. Then two men m and m' must map to the same woman w.



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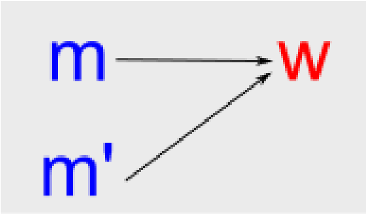

Suppose w prefers m over m'.

Men-optimal mapping is a matching.

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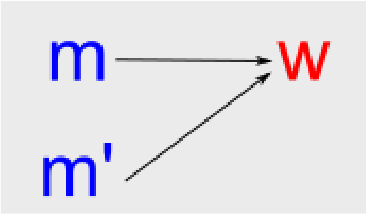

Suppose w prefers m over m'.

There must be a stable matching P where m' and w are matched.

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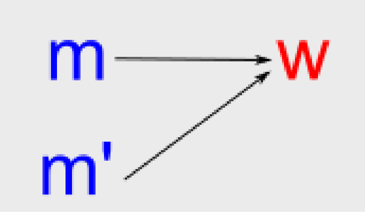
Suppose not. Then two men m and m' must map to the same woman w.



Suppose w prefers m over m'.

There must be a stable matching P where m' and w are matched.

In P, m must be matched to a woman he likes *less* than w (because w is m's favorite achievable woman).

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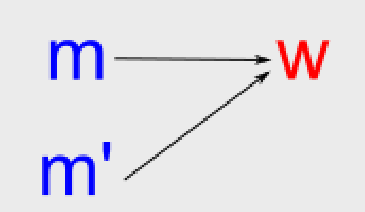

Suppose w prefers m over m'.

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In P, m must be matched to a woman he likes *less* than w (because w is m's favorite achievable woman).

But then, m and w will block P.

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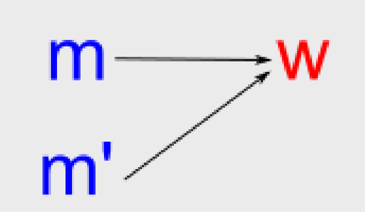
Suppose not. Then two men m and m' must map to the same woman w.



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But then, m and w will block P.

■

Algorithm for computing the men-optimal (or women-optimal) matching?

[Gale and Shapley, 1962]

Given any preference profile, the matching computed by the men-proposing deferred-acceptance algorithm is men-optimal. Similarly, a women-optimal matching is obtained when women propose.

[Gale and Shapley, 1962]

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Suffices to show that in the men-proposing DA algorithm, a man is never rejected by his favorite achievable woman.

[Gale and Shapley, 1962]

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By way of contradiction, suppose man m is the first man to be rejected by his favorite achievable woman w.

Then, w must have received a better proposal from some other man m'.

When m' proposes to w, his past rejections (if any) must all have been from women that are *unachievable* for him.

[Gale and Shapley, 1962]

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Some Observations

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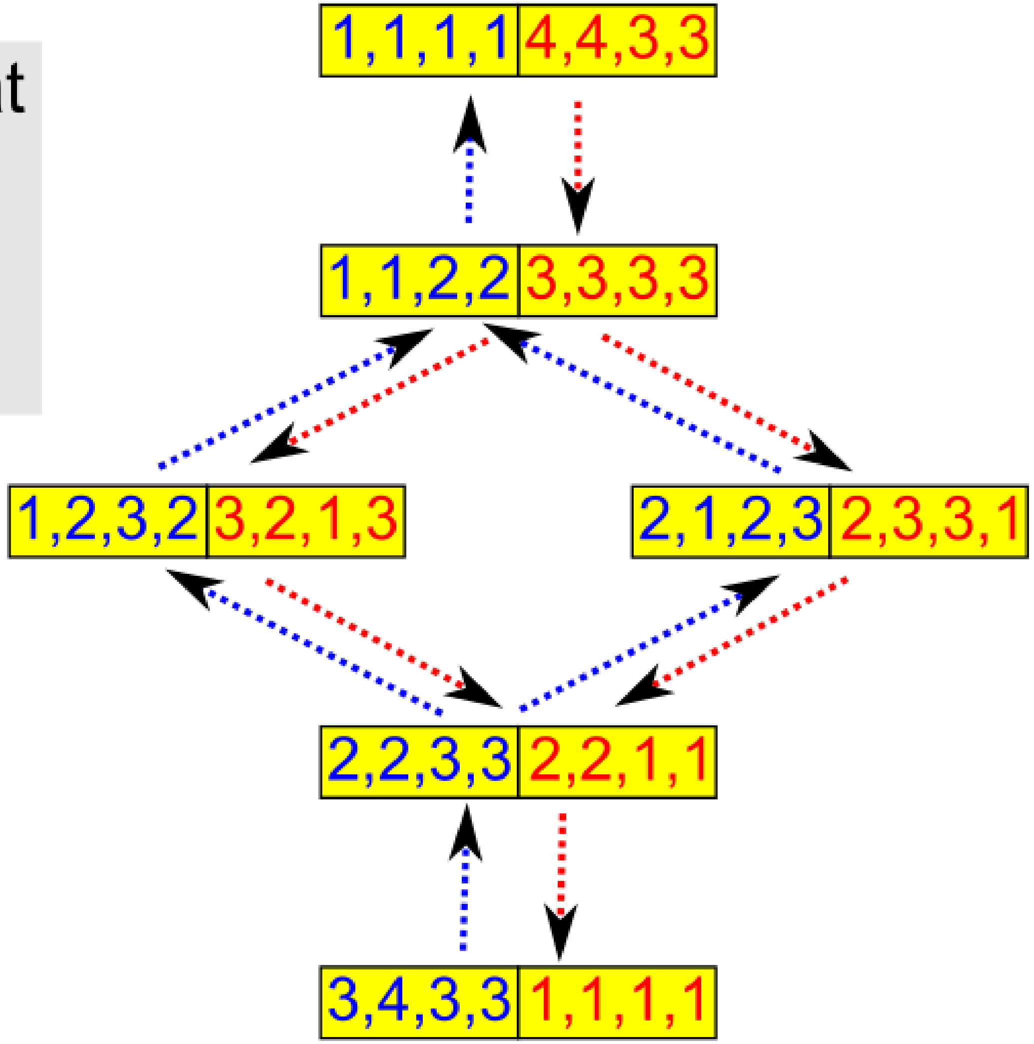
For any distinct stable matchings P and Q, if all men find P at least as good as Q, then all women find Q at least as good as P (and vice versa).



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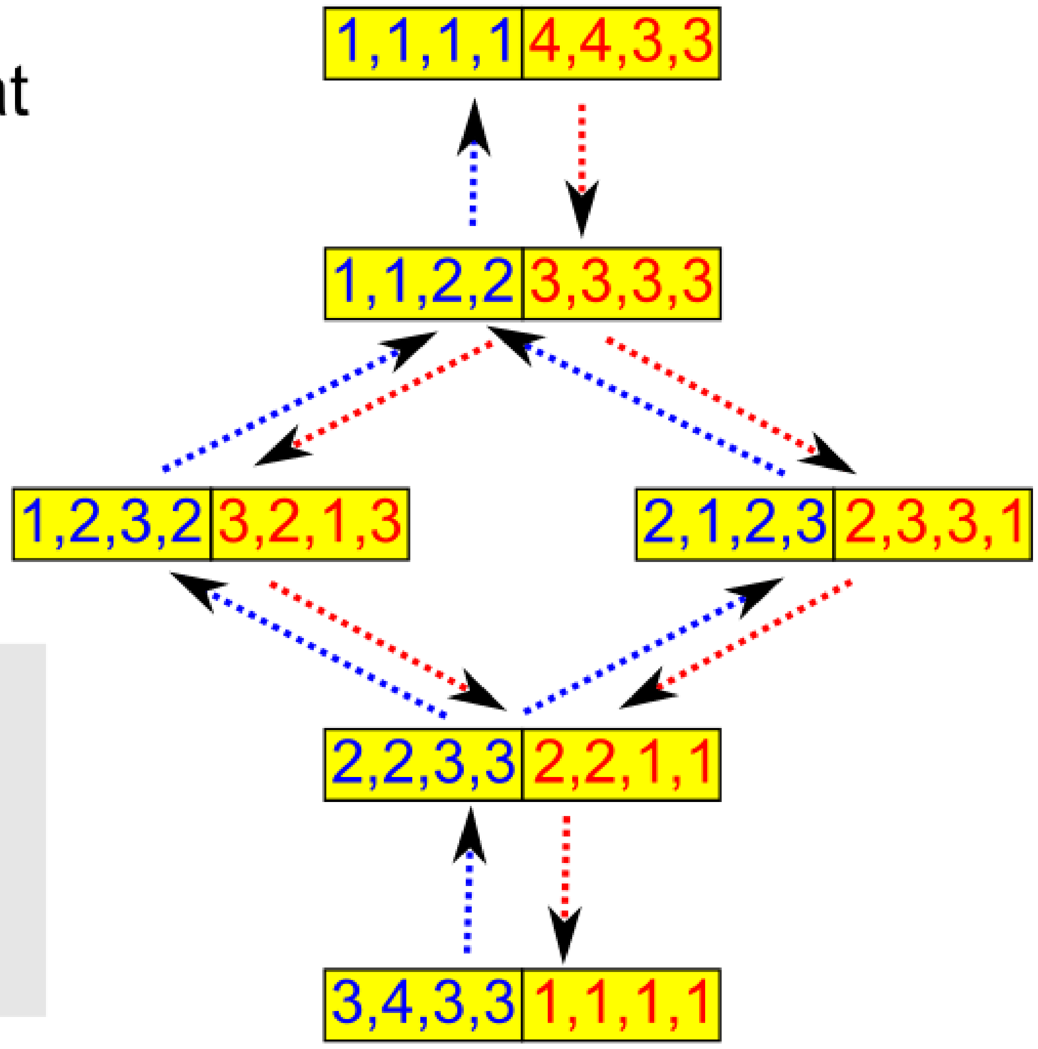

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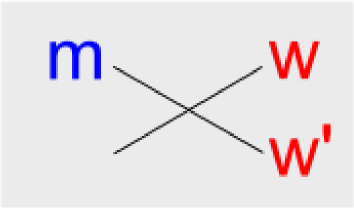

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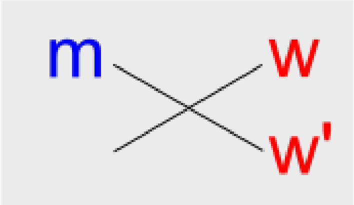

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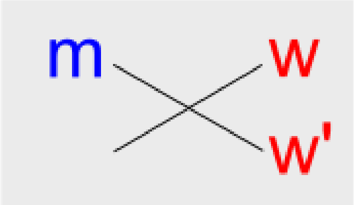
Q: 

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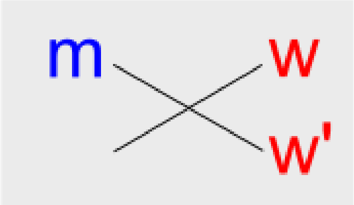

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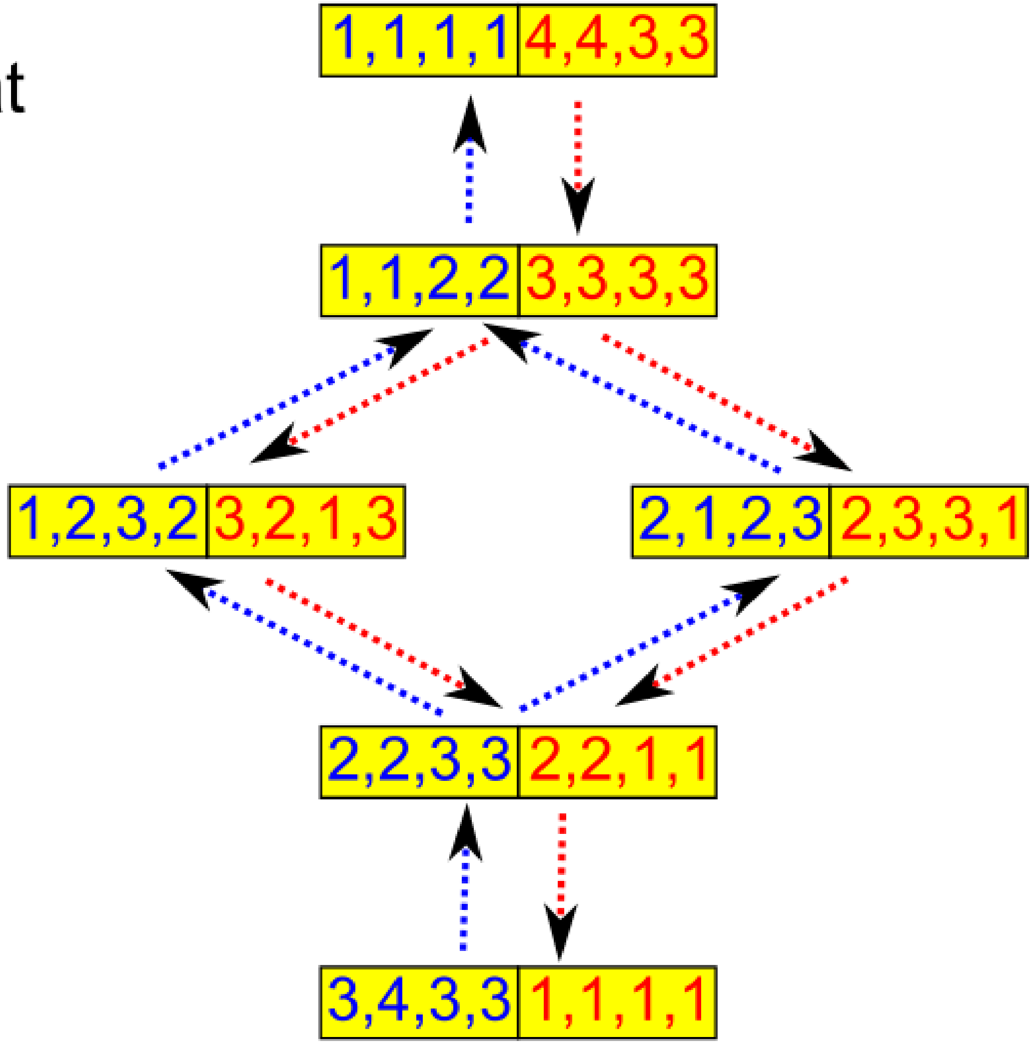

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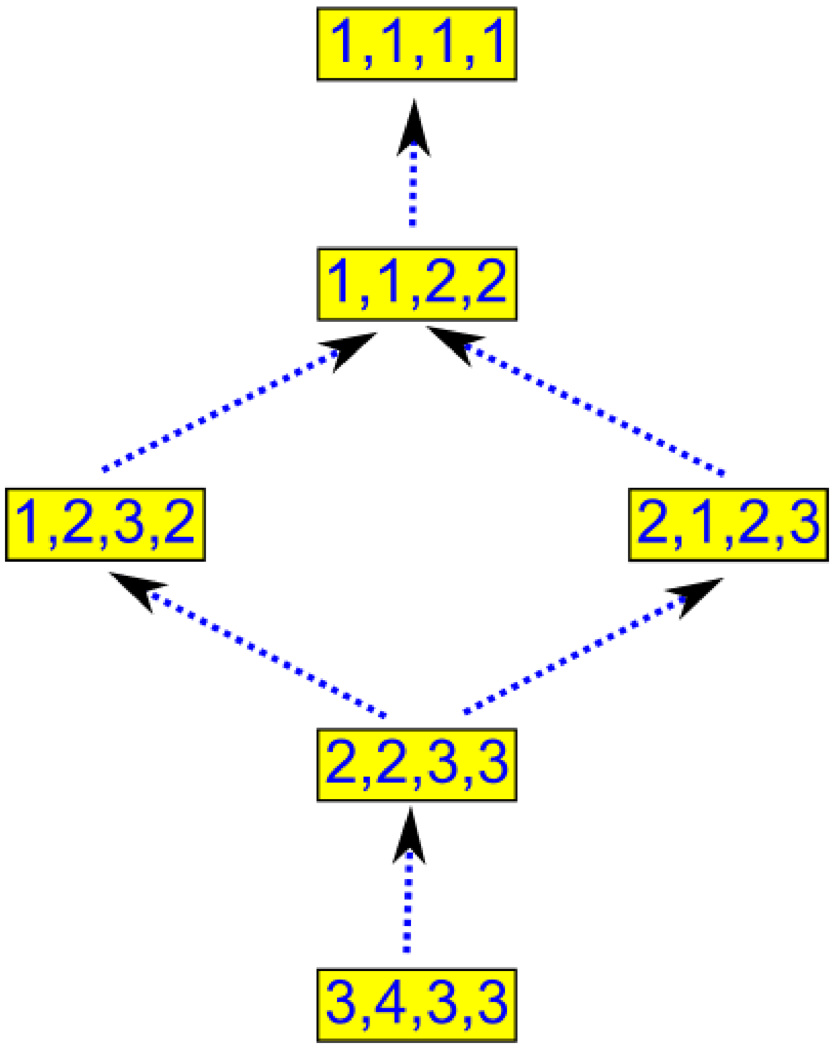

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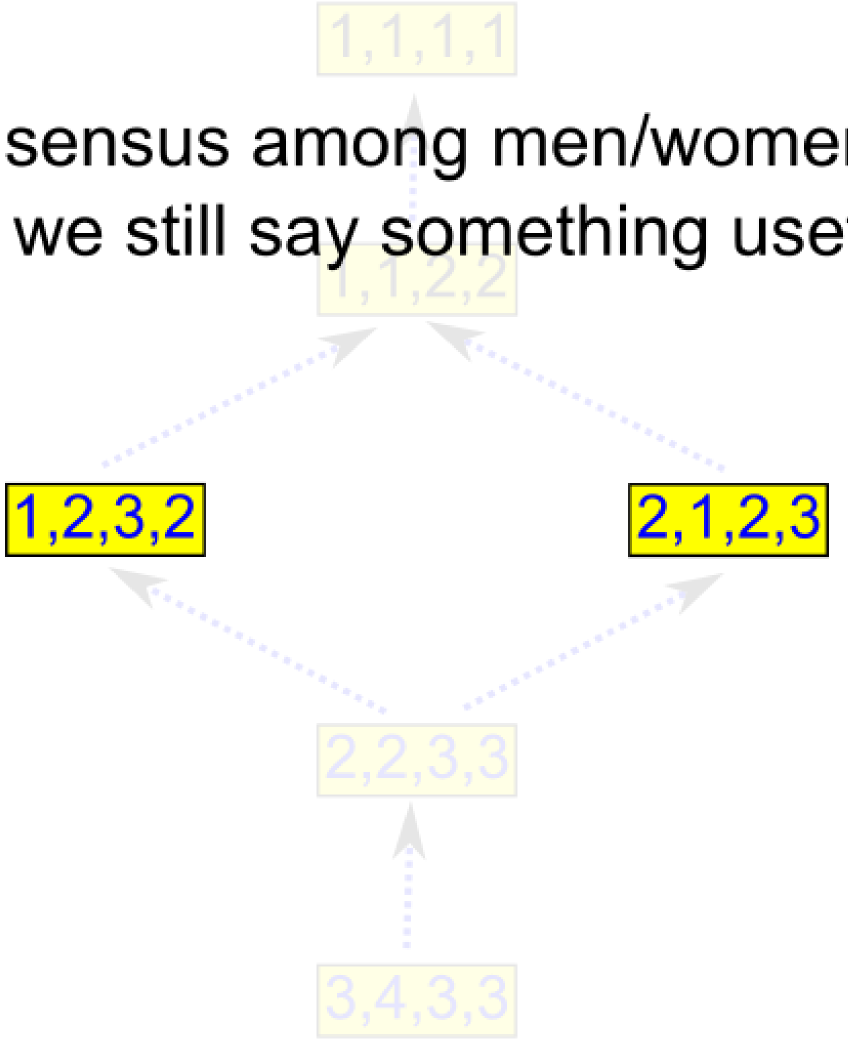
When there isn't a consensus among men/women w.r.t. two matchings, can we still say something useful?



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Let's generalize this idea to arbitrary pairs of stable matchings.

Let P and Q be any pair of stable matchings (not necessarily distinct).

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Suffices to show that for any m and w, $\max_{P,Q}(m) = w \Leftrightarrow \max_{P,Q}(w) = m$.

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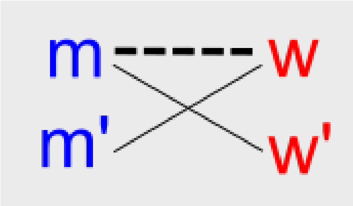
Q: 

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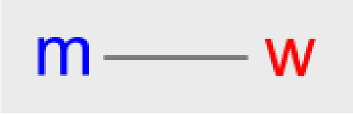

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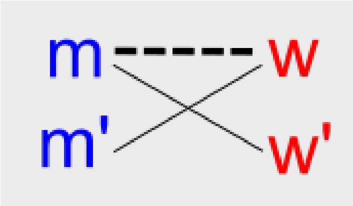

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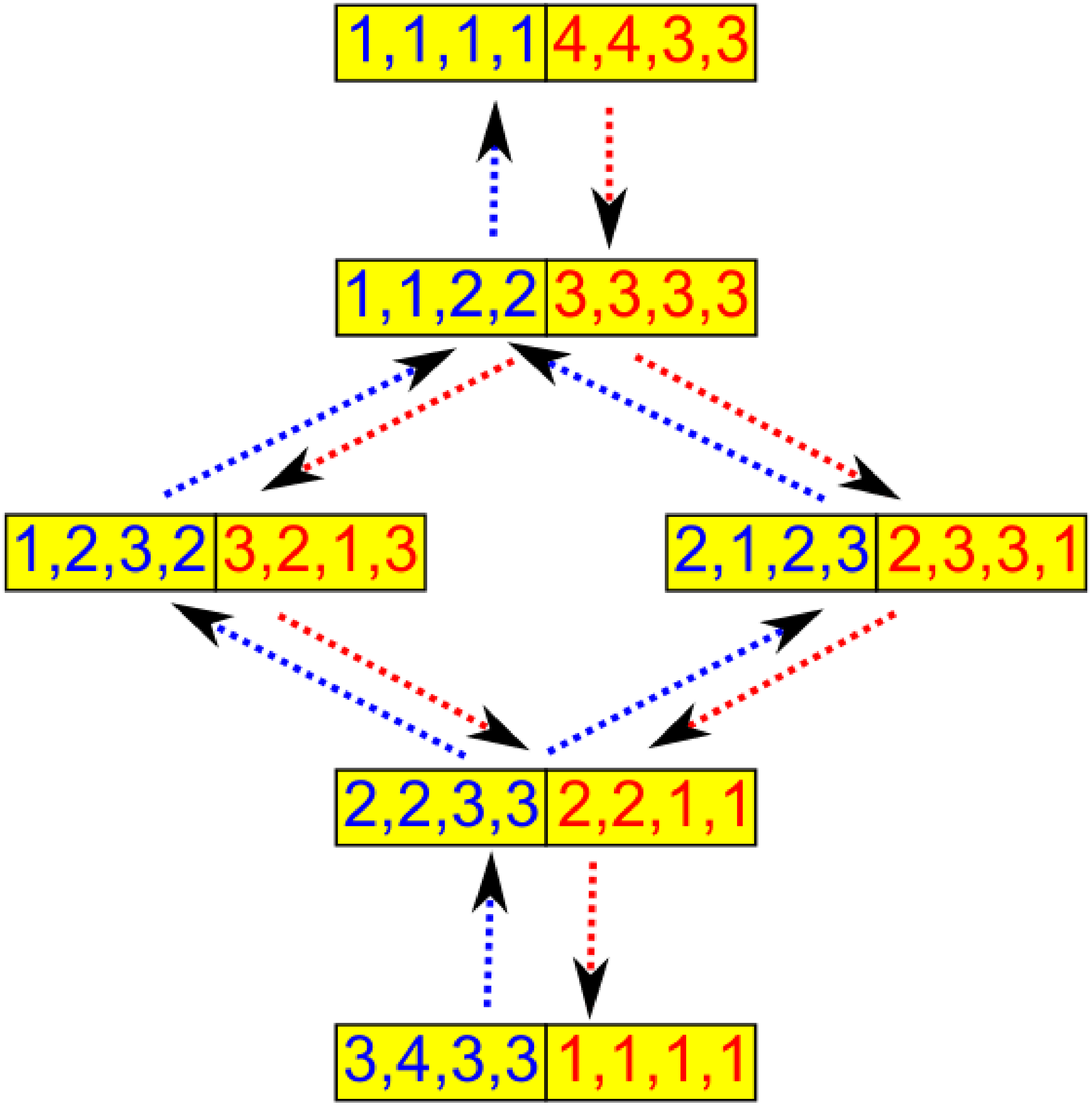
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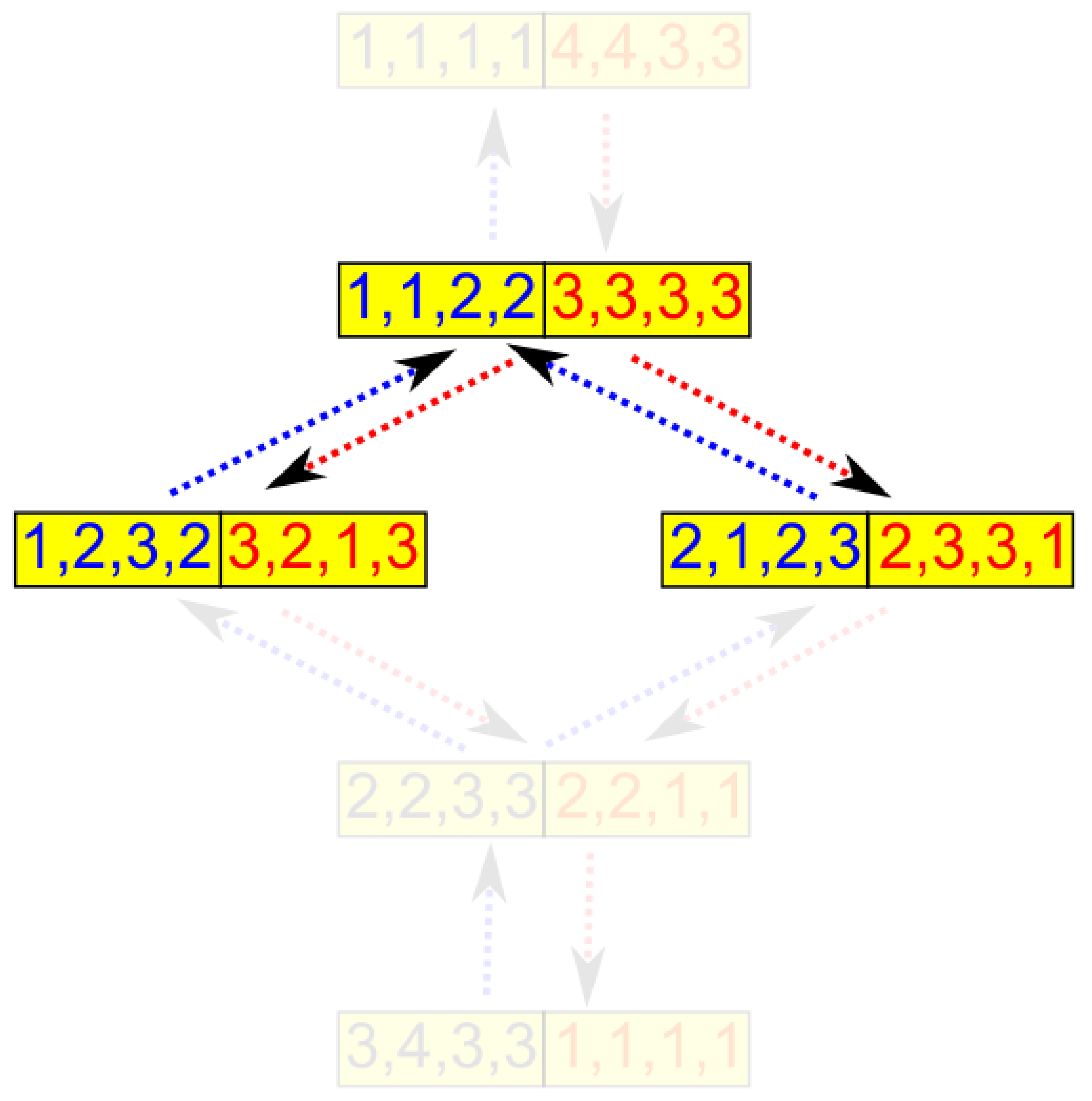
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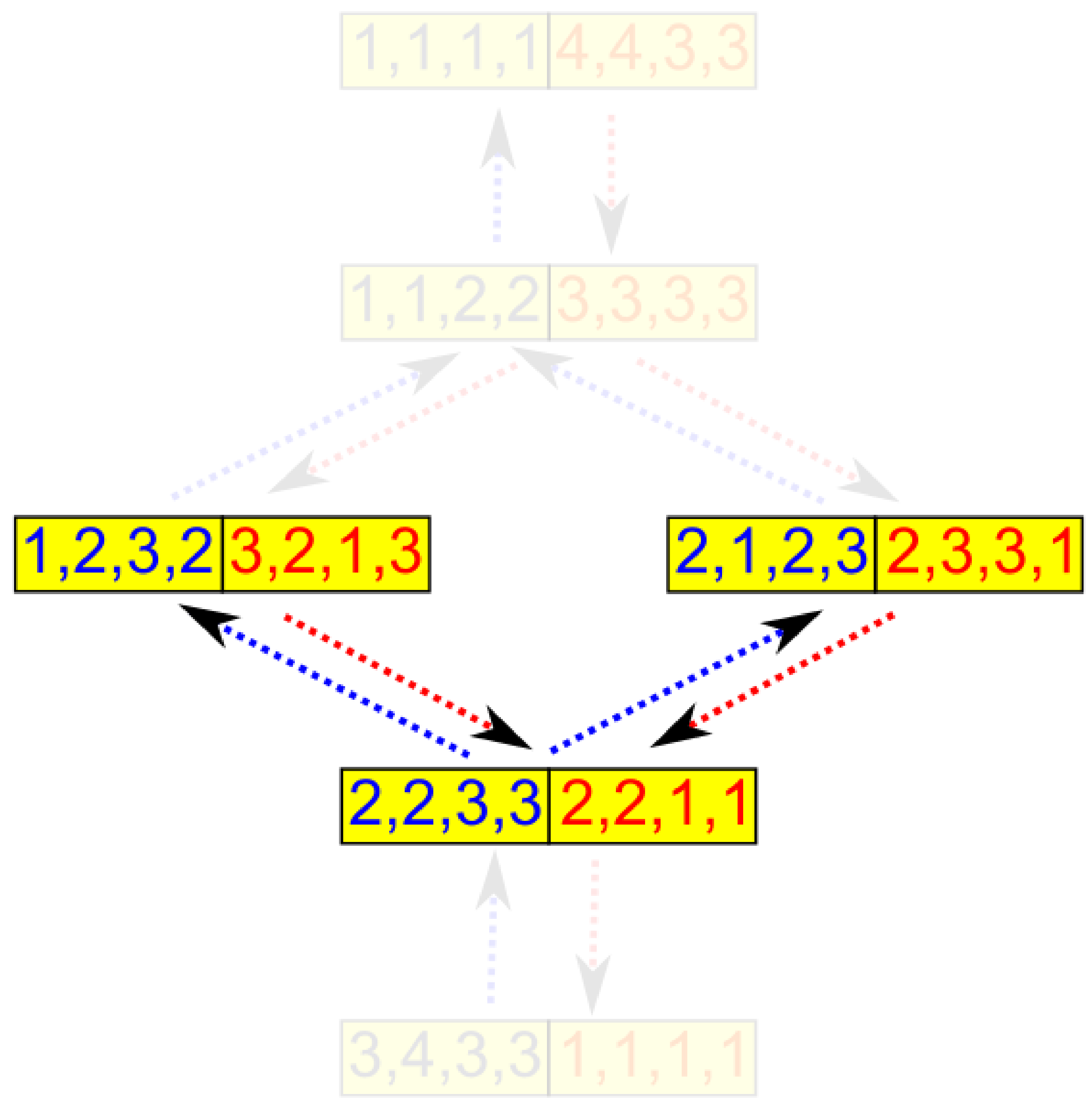
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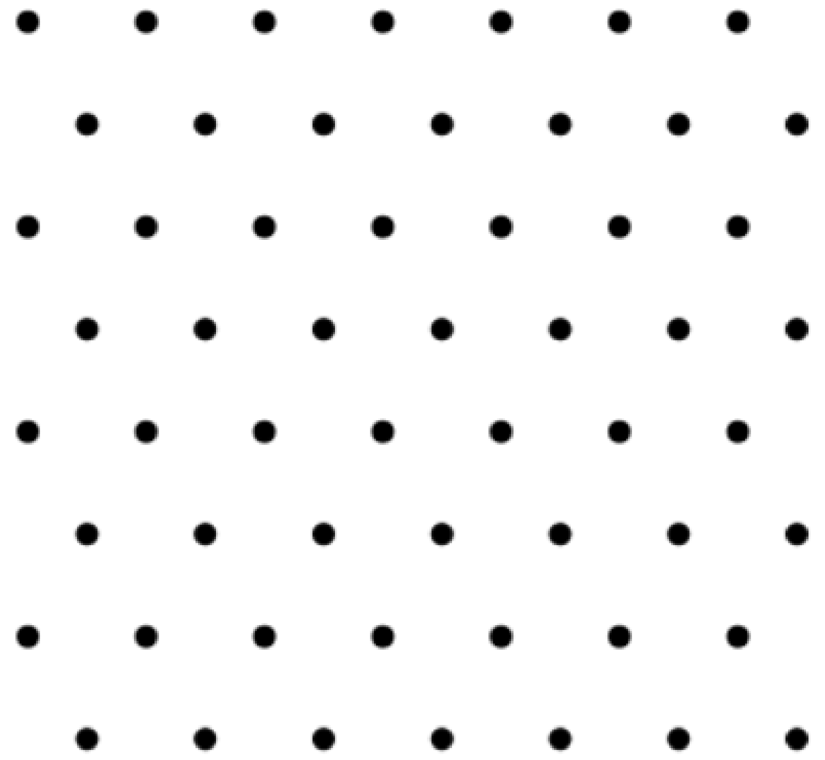

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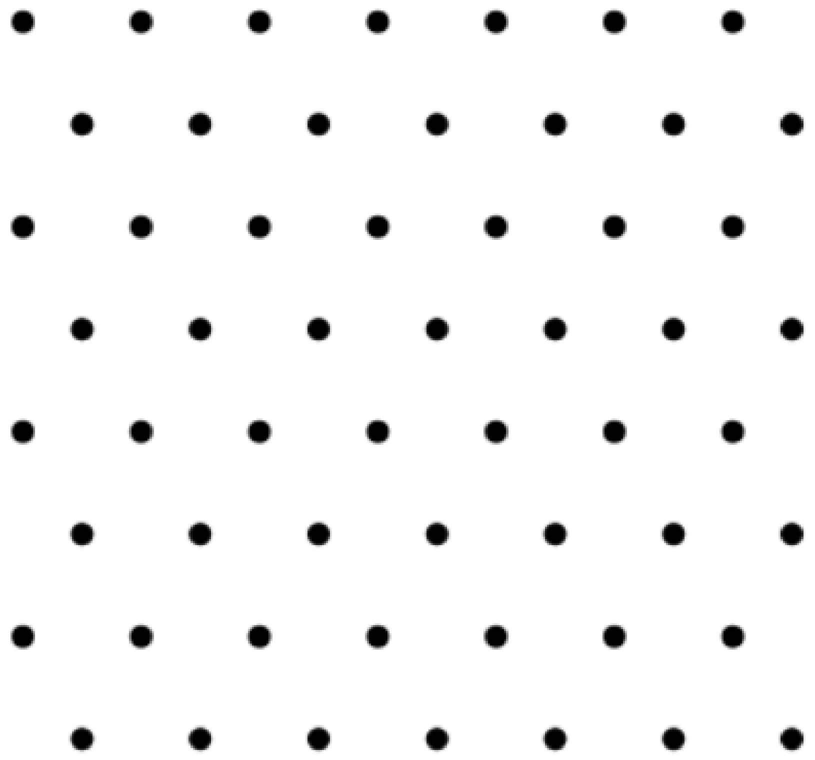

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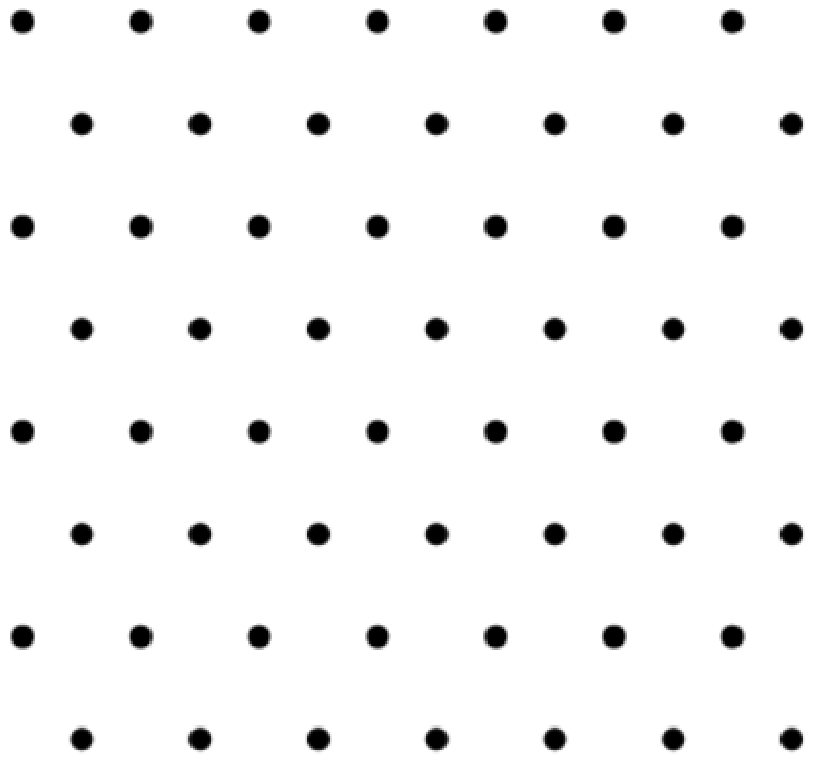

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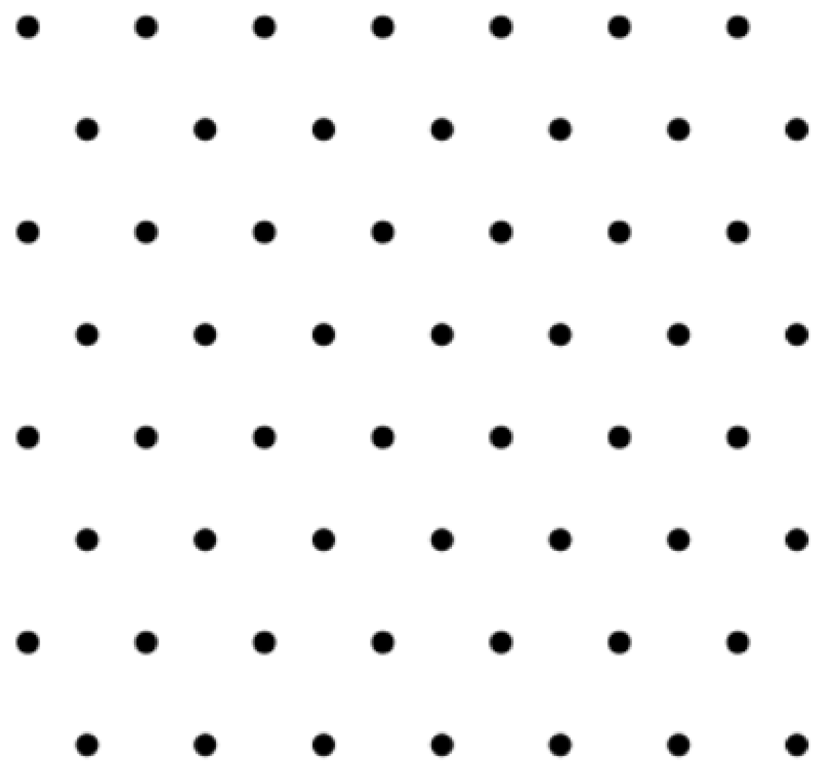

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The Rural Hospitals Theorem

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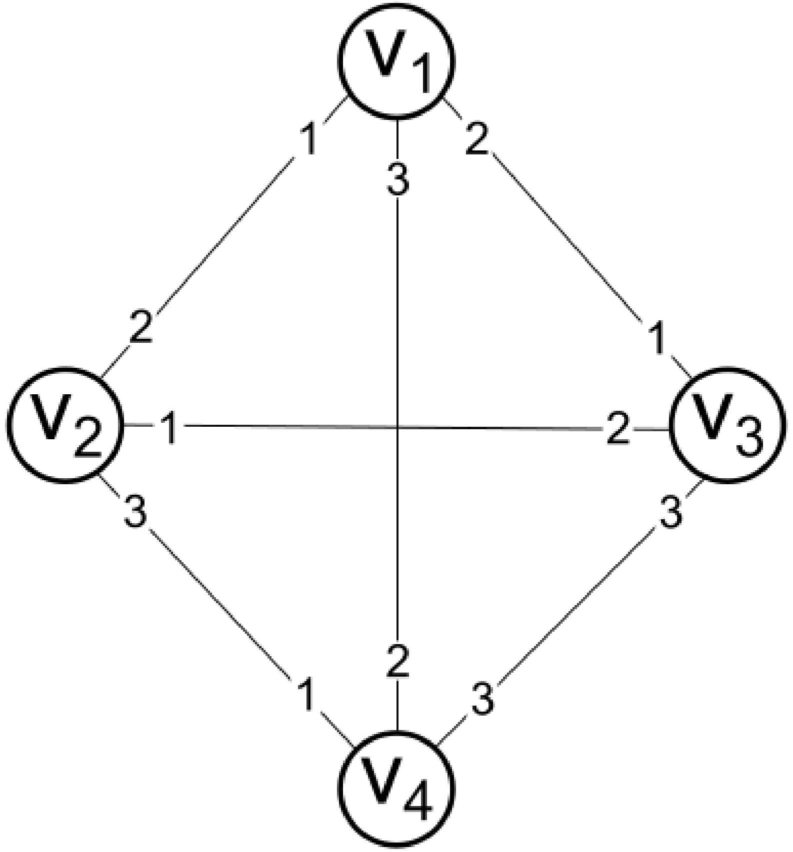

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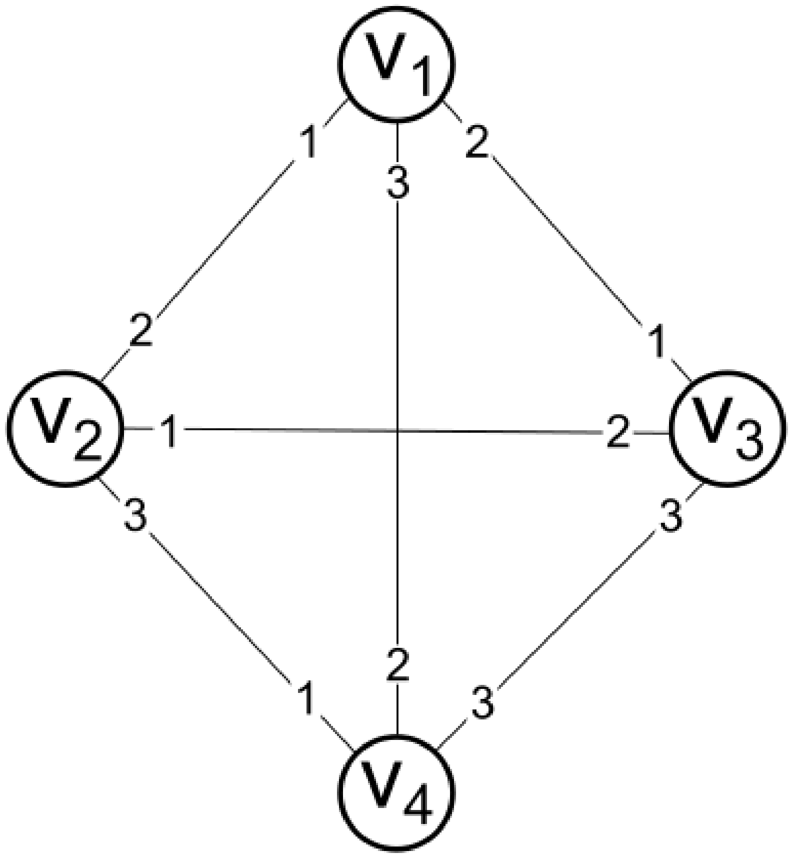

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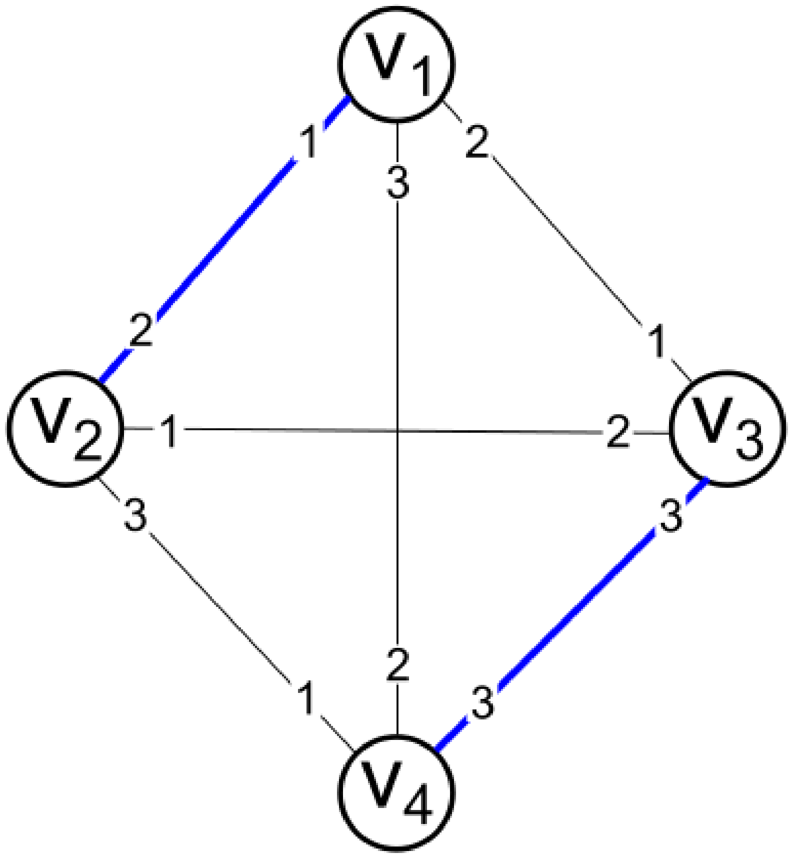

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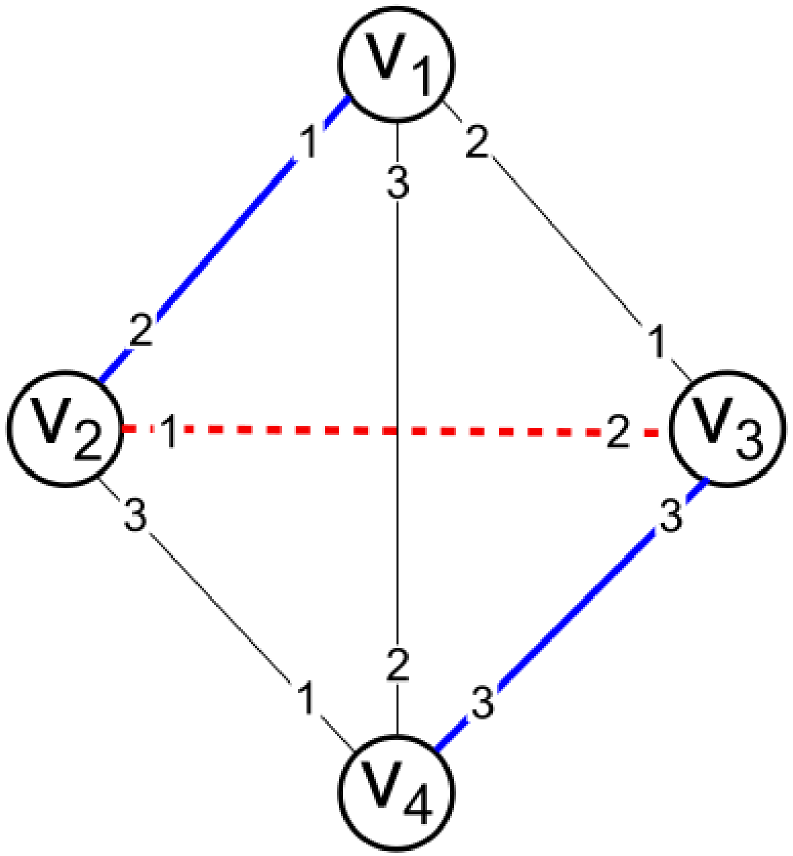

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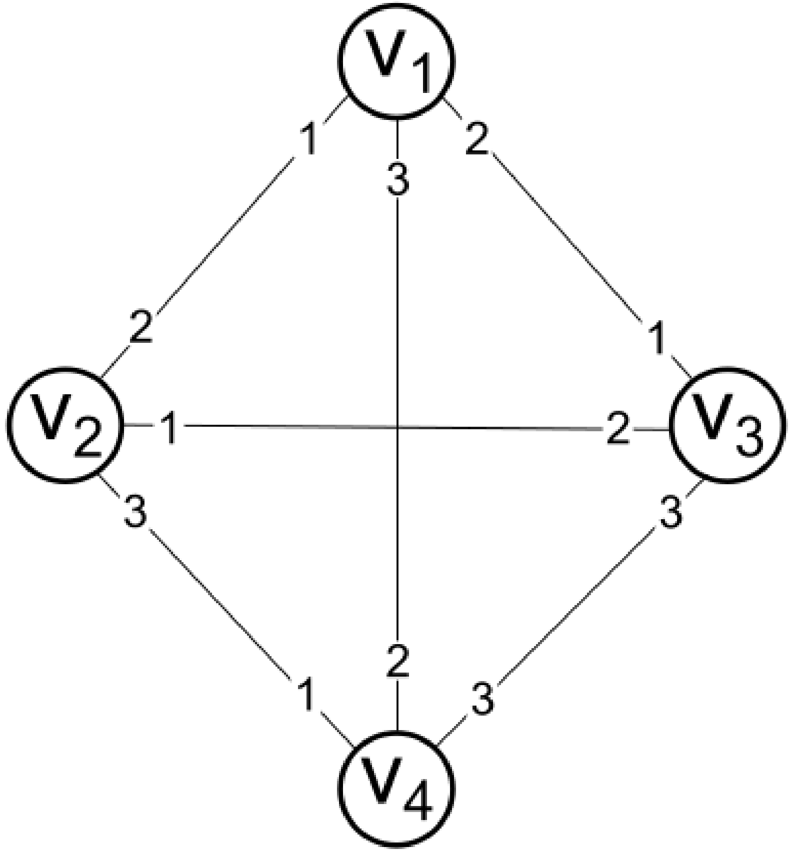

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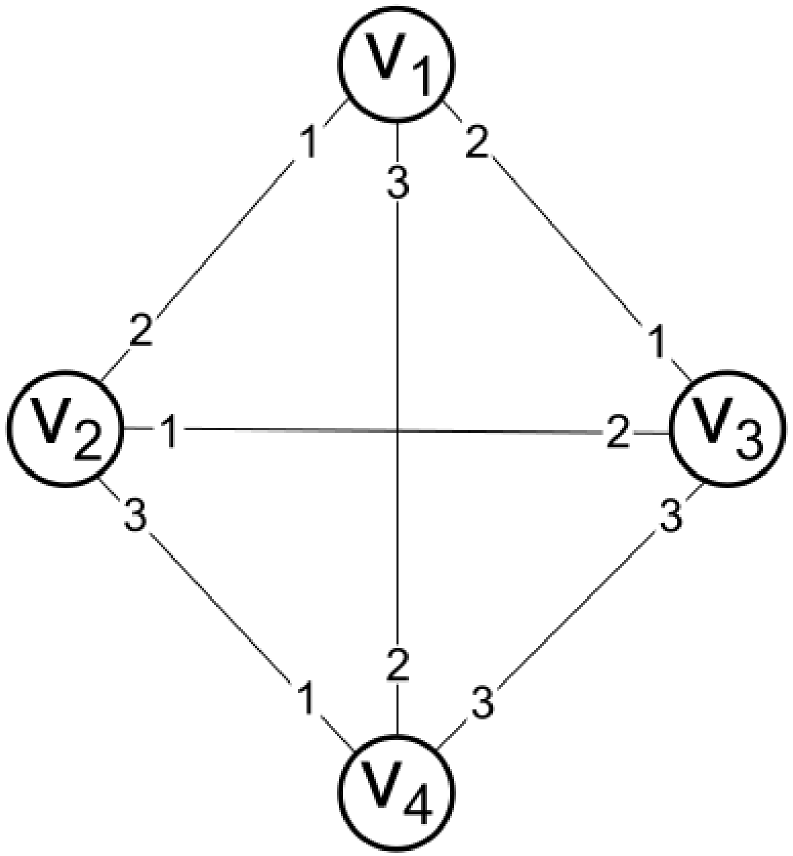

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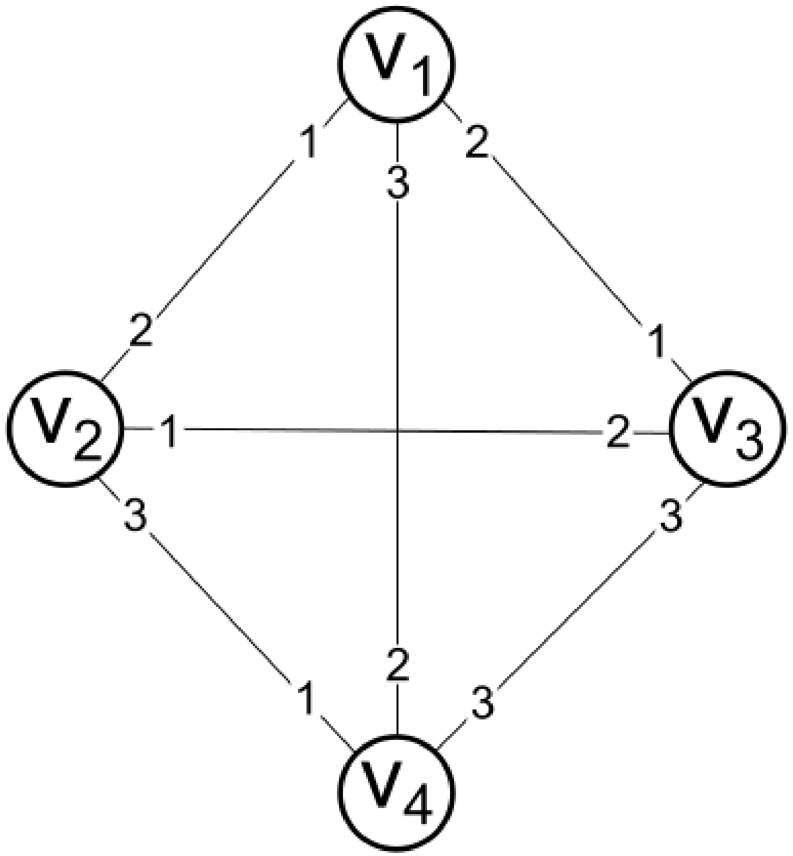

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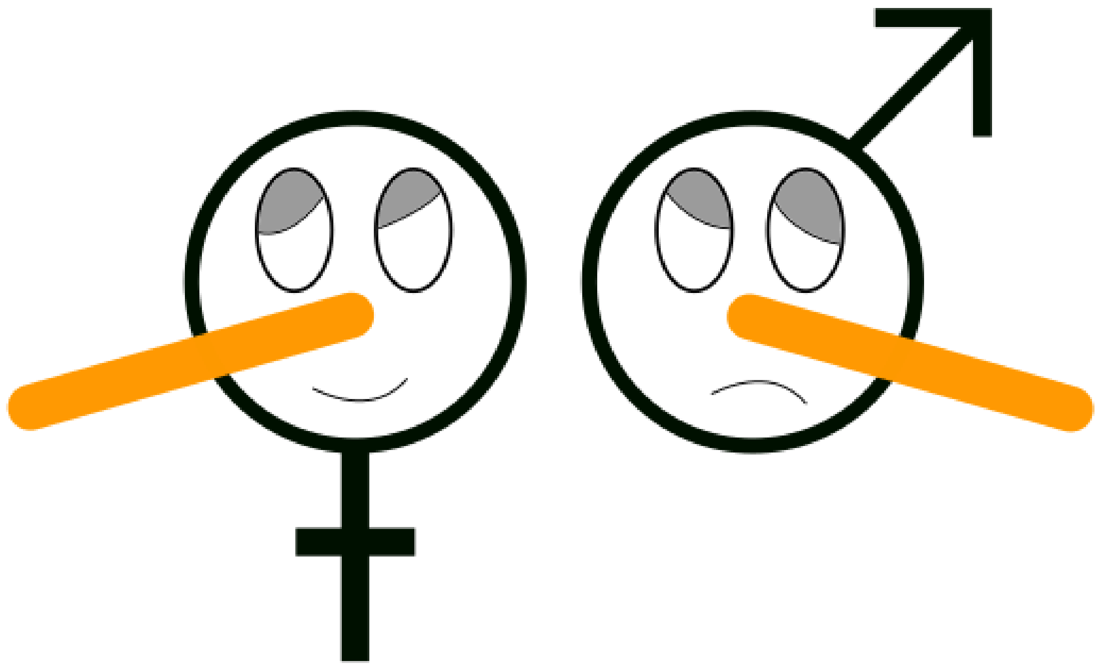
Whoever is matched with v_4 will block with one of the other two agents.

Next Time

Incentives in the Stable Matching Problem



References

- Structure of the Set of Stable Matchings

Alvin Roth and Marilda Sotomayor
"Two-Sided Matching: A Study in Game-Theoretic Modeling and Analysis"
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