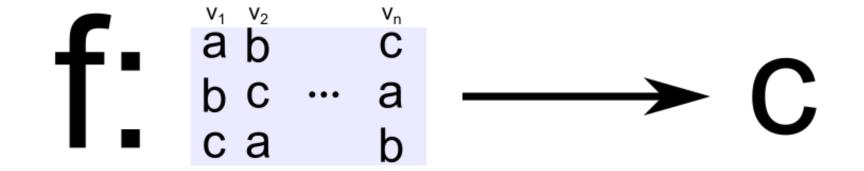
COL866: Special Topics in Algorithms

Lecture 13 Manipulation in Voting

VOTING RULE

A mapping from preference profiles to candidates.

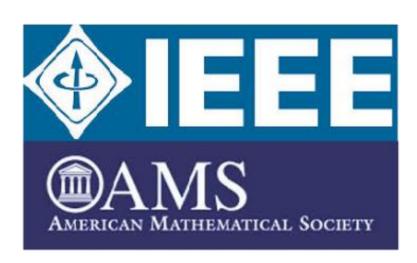


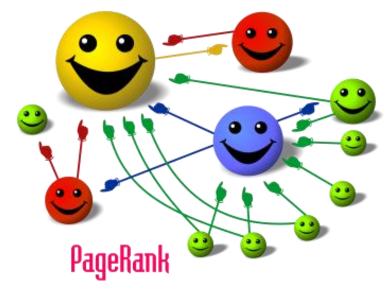
(also known as a social choice function)

Score-based	Runoff-based	Head-to-head election based
Plurality	Plurality with Runoff	Copeland
Borda Count	Single Transferable Vote	Schulze

Voting Rules are Everywhere!





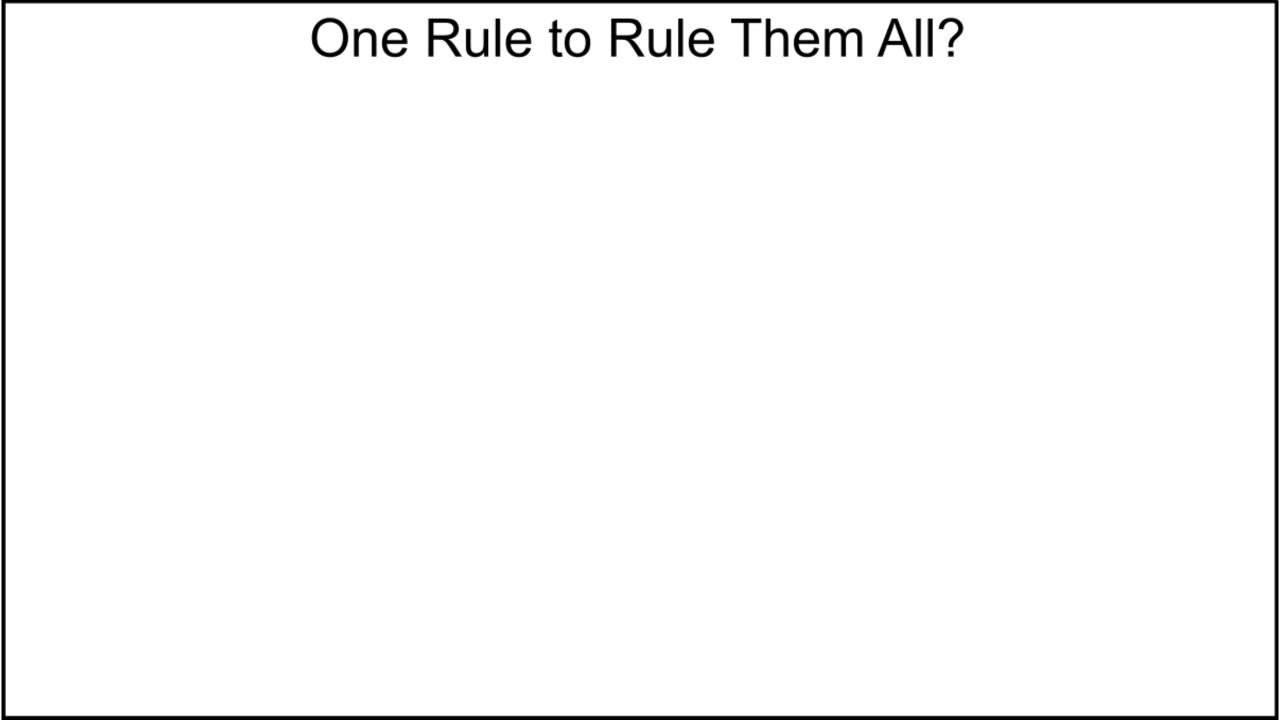








Kemeny-Young Copeland Winner **Ranked** Borda Majority-Judgement Approval Schul



One Rule to Rule Them All?

Axiomatic Approach

- Formulate a set of "reasonable" properties
- Check if there is a voting rule that satisfies them

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Unique voting rule with all the desirable properties



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Axiomatic Approach

- Formulate a set of "reasonable" properties
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Expectation

Unique voting rule with all the desirable properties



Reality

No such voting rule



ONTO

For any candidate "a", there exists a profile where "a" wins.

STRATEGYPROOF

No voter can improve by misreporting its preferences.

For any profile , any voter v_i , and any misreport , it must be that

$$f(\frac{v_1 - v_1 - v_2}{1 - v_1}) = f(\frac{v_1 - v_1}{1 - v_2})$$



An Obviously Bad Voting Rule

A voting rule is called a dictatorship if there exists a voter v_i such that for any preferences of the other voters, the voting outcome is the favorite candidate of voter v_i.

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A dictatorship is onto and strategyproof.





With three or more candidates, a voting rule is onto and strategyproof if and only if it is a dictatorship.





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Onto + strategyproof but two candidates:





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Majority





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≥3 candidates + strategyproof but not onto:





With three or more candidates, a voting rule is onto and strategyproof if and only if it is a dictatorship.

Onto + strategyproof but two candidates: Majority

≥3 candidates + strategyproof but not onto: Constant/Restricted Majority





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Onto + strategyproof but two candidates: Majority

≥3 candidates + strategyproof but not onto: Constant/Restricted Majority

≥3 candidates + onto but not strategyproof: Plurality/Borda/...

With three or more candidates, a voting rule is onto and strategyproof if and only if it is a dictatorship.

With three or more candidates, a voting rule is unanimous and monotone if and only if it is a dictatorship.

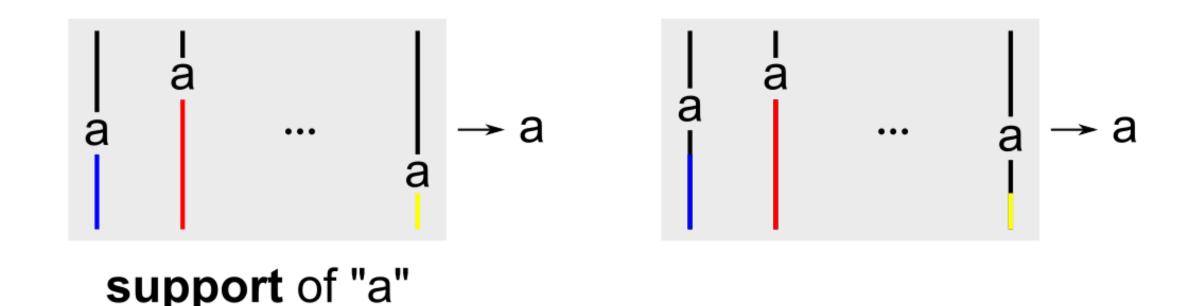
UNANIMOUS

If all voters agree on their top choice, then the voting rule picks it.

$$\mathbf{f}\left(\begin{smallmatrix}v_1&v_2&&v_n\\a&a&&a\\&\cdot&\cdot&&\cdot\\&\cdot&\cdot&&\cdot\\&\cdot&&\cdot\end{smallmatrix}\right)=\mathbf{a}$$

MONOTONE

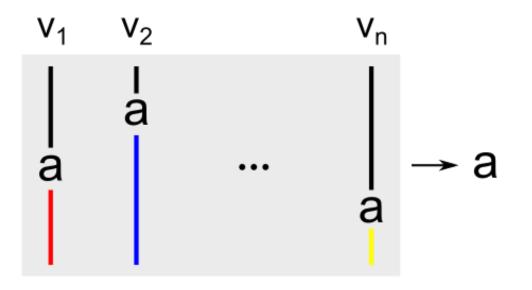
Suppose "a" is the current winner, and for each voter, any candidate that was ranked below "a" in its old vote is ranked below "a" in its new vote, then "a" continues to be the winner.

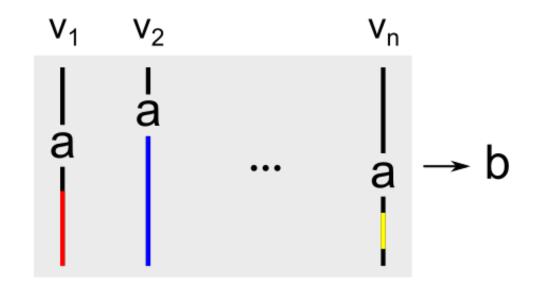


Strategyproofness = Monotonicity

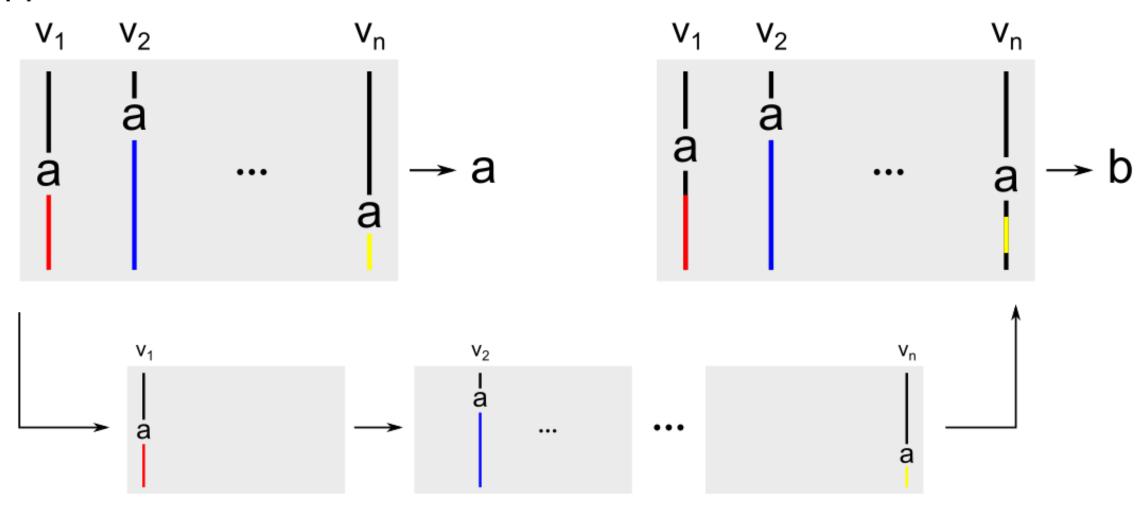
Suppose SP but not MON.

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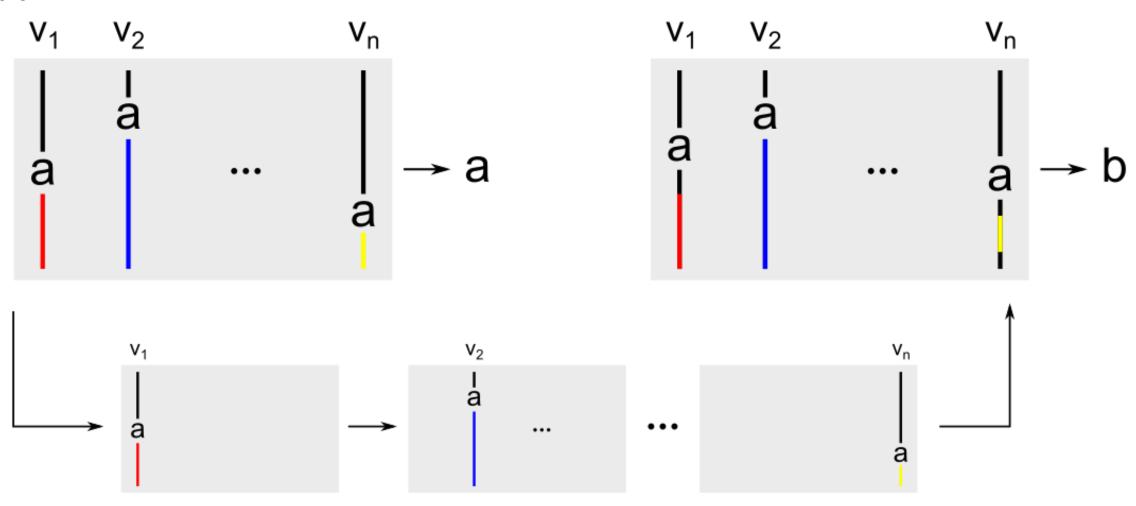




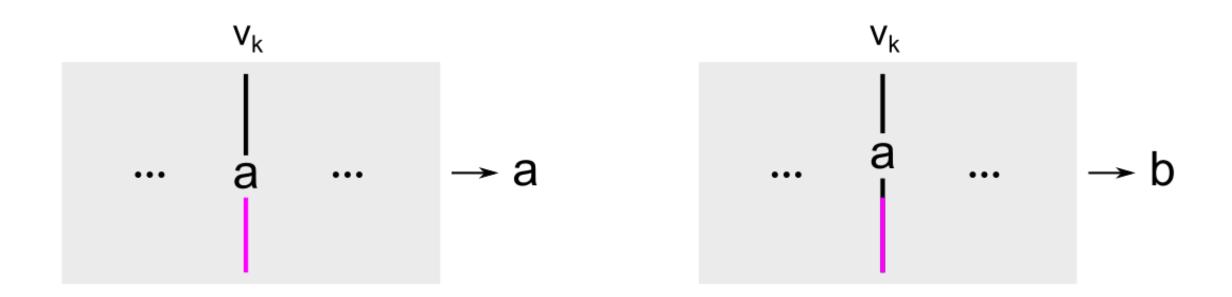
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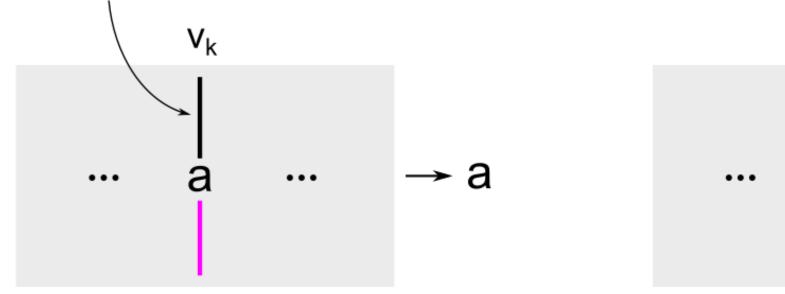
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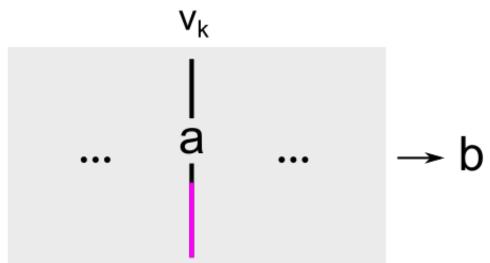


Suppose v_k is the first voter for which the outcome changes.

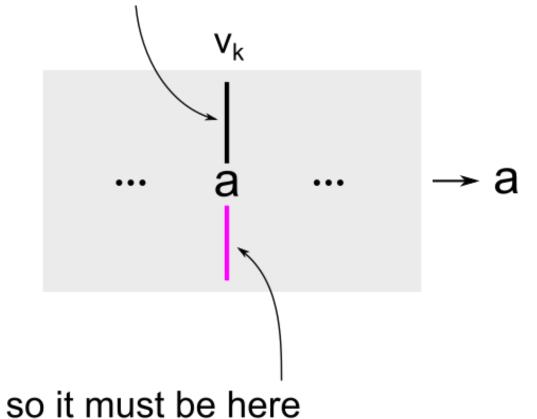


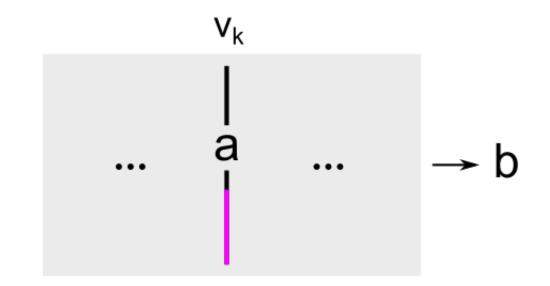
b can't be here (SP)



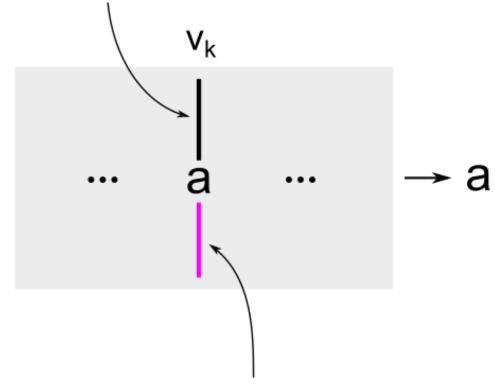


b can't be here (SP)

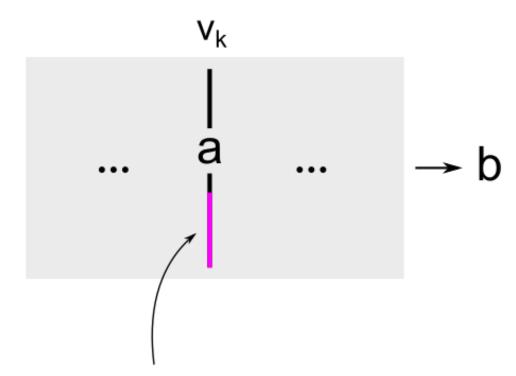




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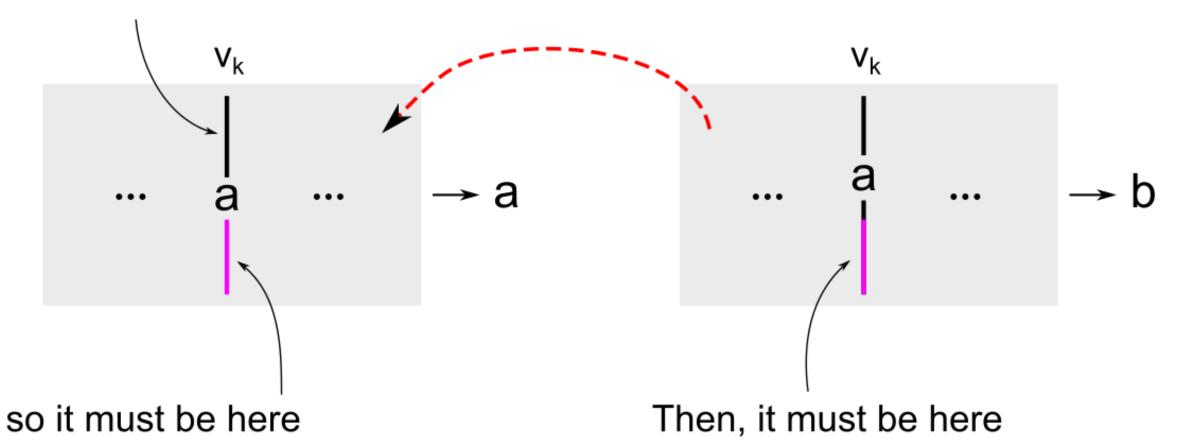


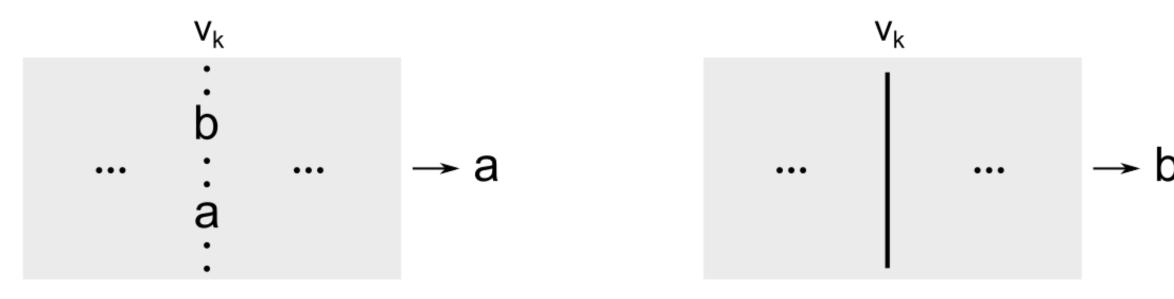
so it must be here

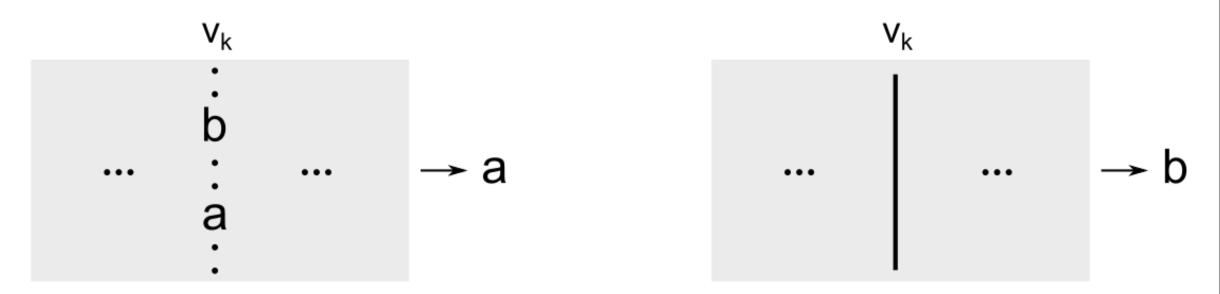


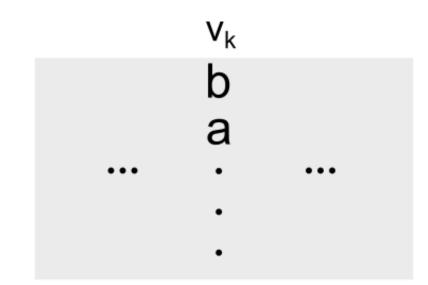
Then, it must be here

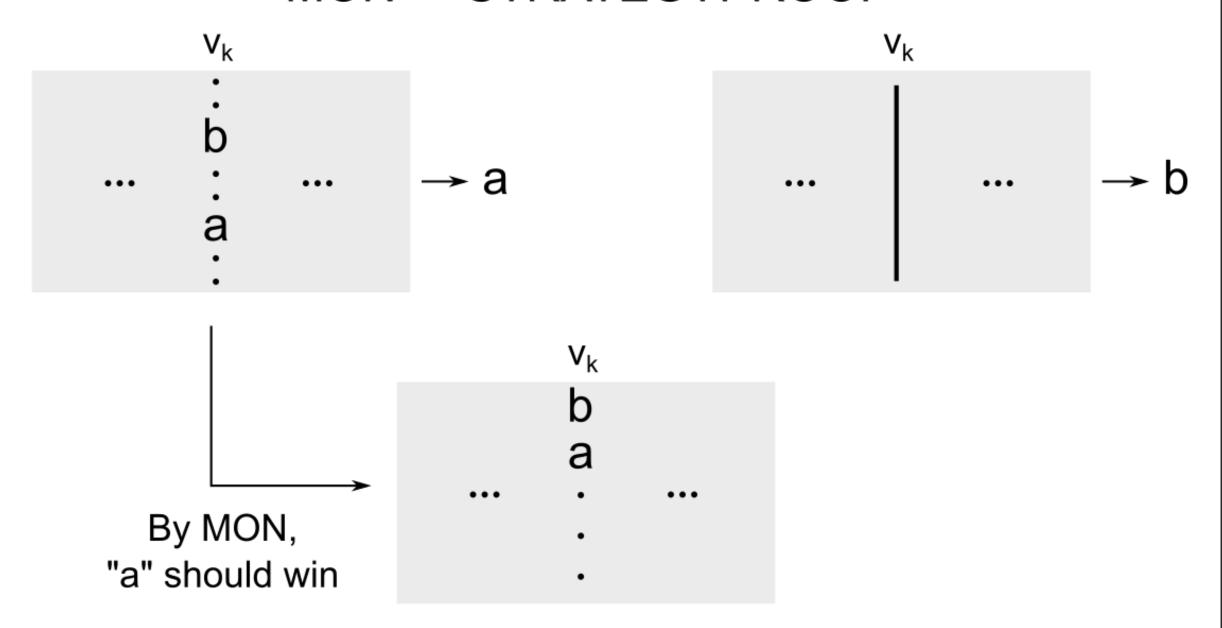
b can't be here (SP) But then there is a reverse manipulation

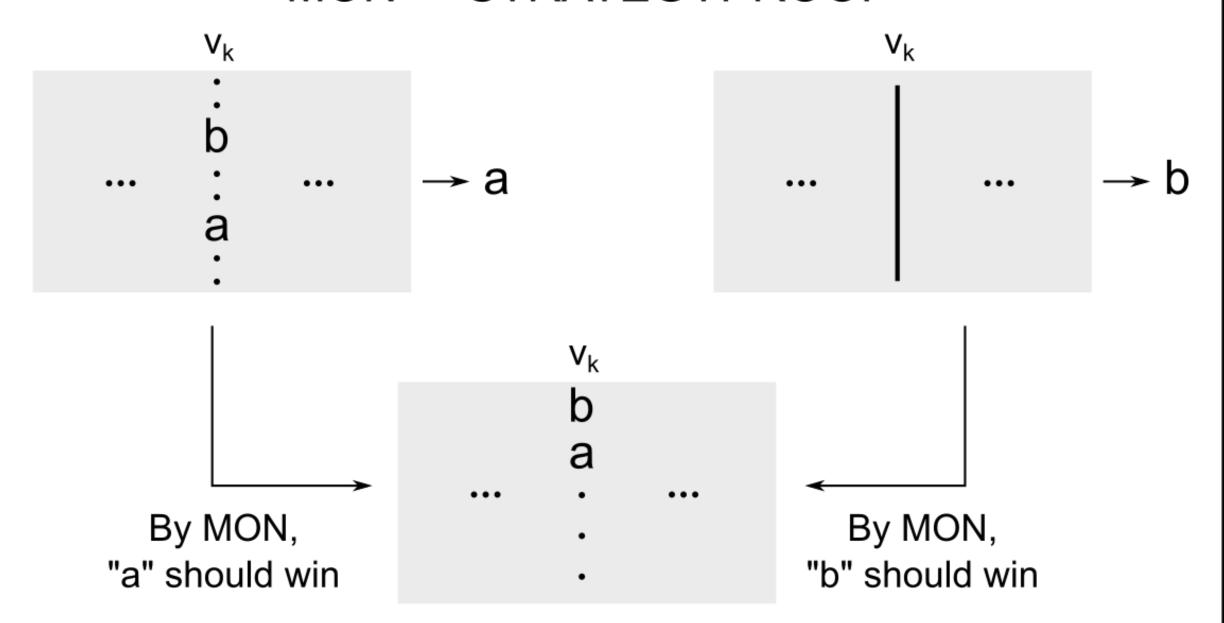








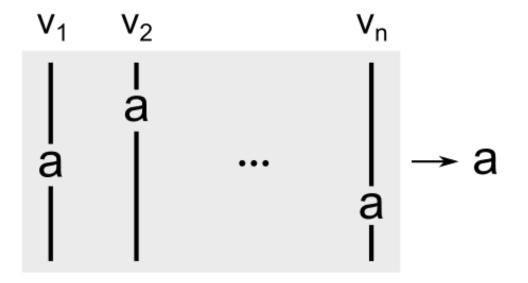




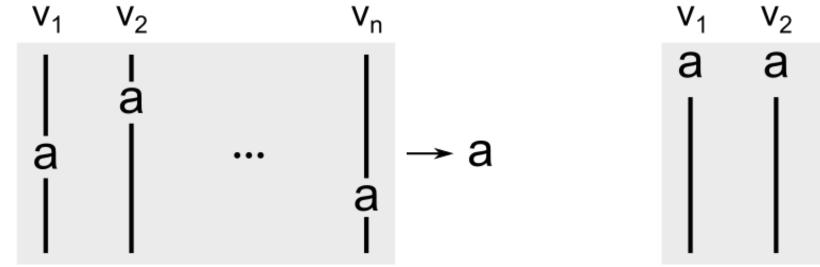
With three or more candidates, a voting rule is onto and strategyproof if and only if it is a dictatorship.

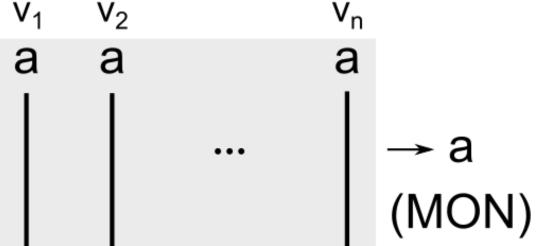
With three or more candidates, a voting rule is unanimous and monotone if and only if it is a dictatorship.

For any candidate "a", by ONTO, there is a profile where "a" wins

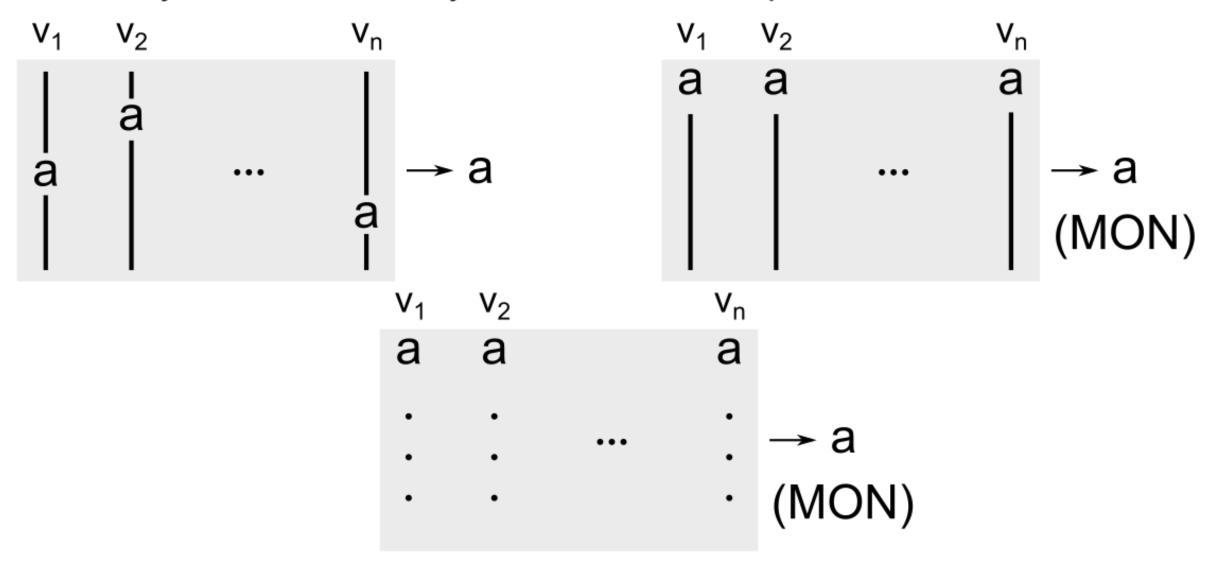


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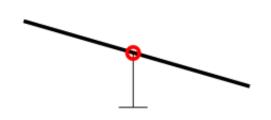




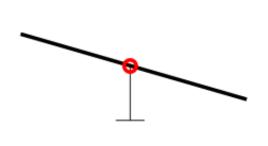
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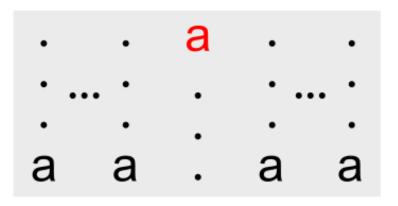


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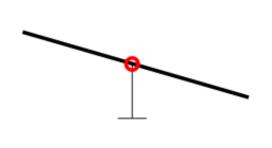
Identify a "pivotal voter" v_p for the pair $\{a,b\}$

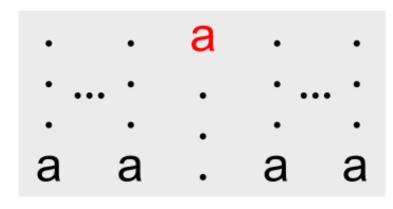


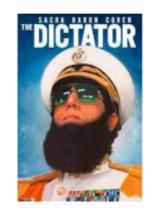


Identify a "pivotal voter" v_p for the pair $\{a,b\}$

Show that "a" wins if v_p likes it even if everyone else ranks it last



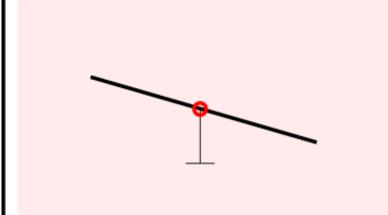


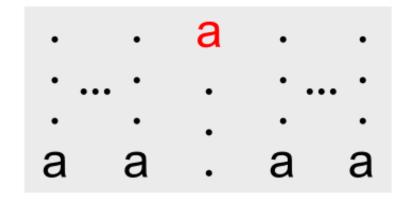


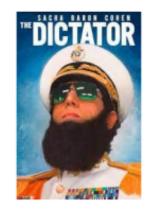
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v_p is the dictator for "a" (and every other candidate)





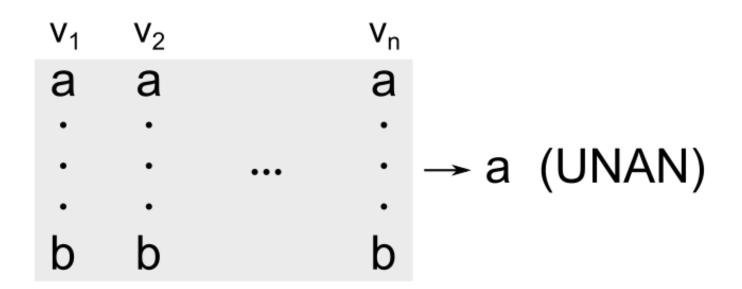


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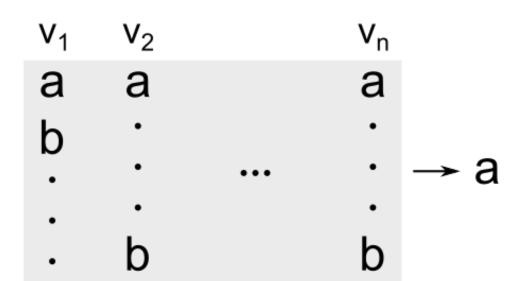
V_1	V_2		v_n
а	а		а
•	•		•
•	•	•••	•
•	•		•
b	b		b

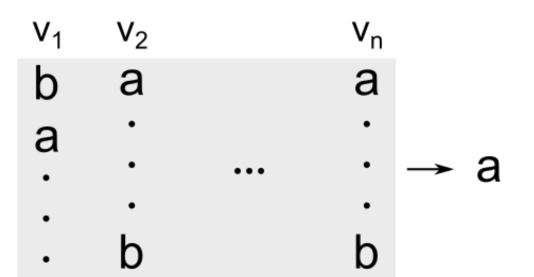


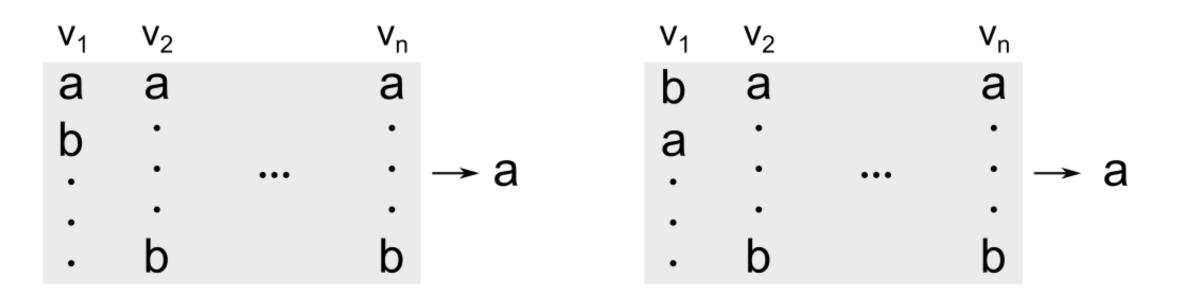
V_1	V_2		v_{n}
а	а		а
b	•		•
•	•	•••	•
	•		•
•	b		b

V_1	V_2		v_n		V_1	V_2		\mathbf{v}_{n}
а	а		а		b	а		a
b	•		•		а	•		•
•	•	•••	•	→ a	•	•	•••	•
•	•		•	(NAONI)		•		•
•	b		b	(MON)	•	b		b

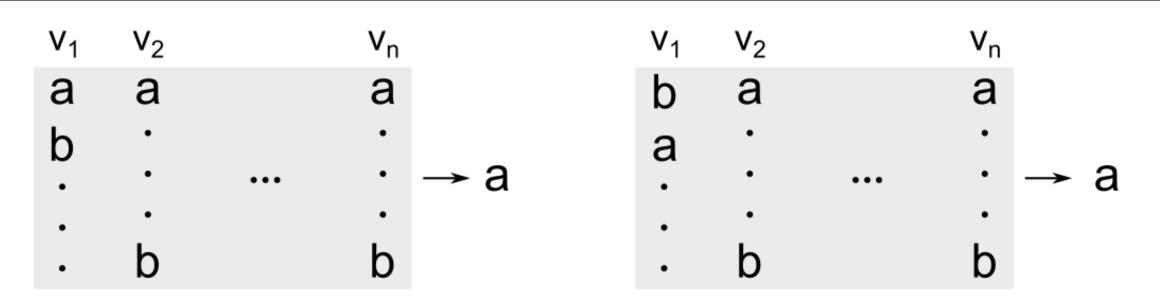
v_1	V_2		\mathbf{v}_{n}		v_1	V_2		\mathbf{v}_{n}	
a	а		а		b	а		а	
b	•		•		а	•		•	
•	•	•••	•	→ a	•	•	•••	•	→ {a,b}
•	•		•	(NAONI)		•		•	
•	b		b	(MON)	•	b		b	(MON)



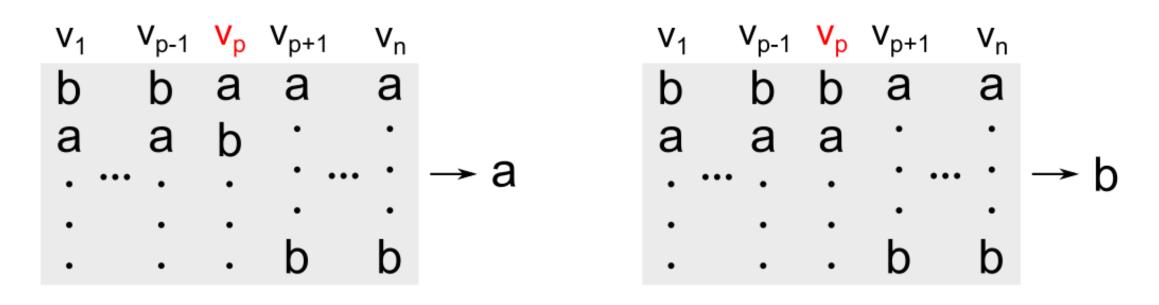


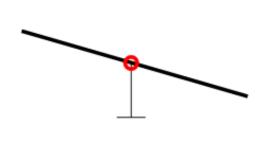


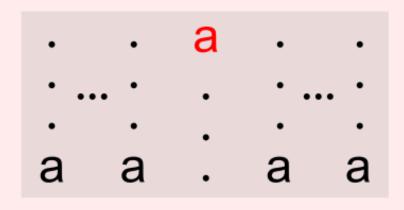
"b" must win eventually due to UNAN

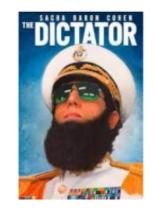


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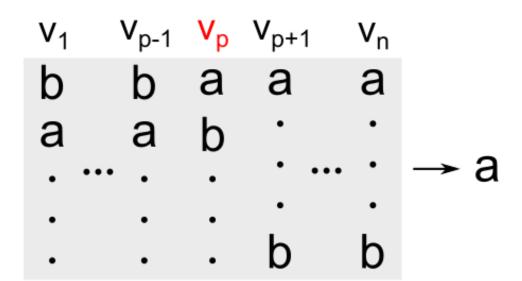


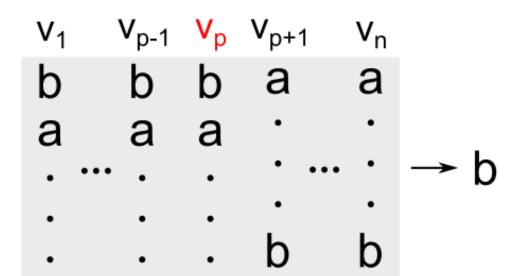


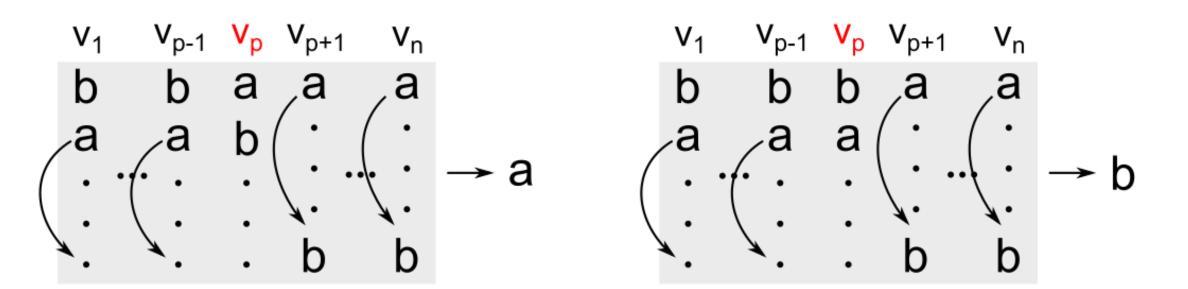
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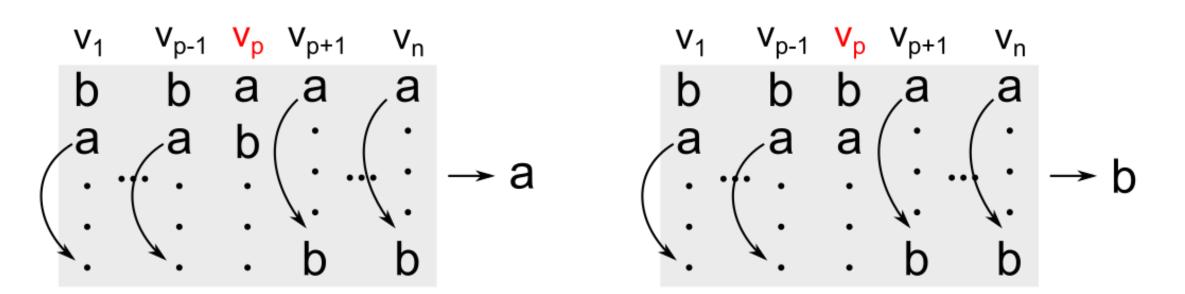
Show that "a" wins if v_p likes it even if everyone else ranks it last

v_p is the dictator for "a" (and every other candidate)

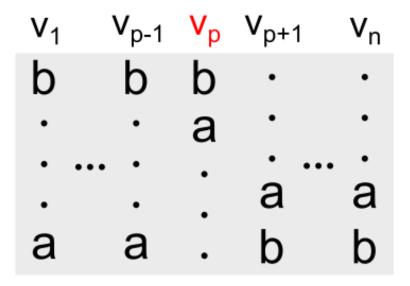


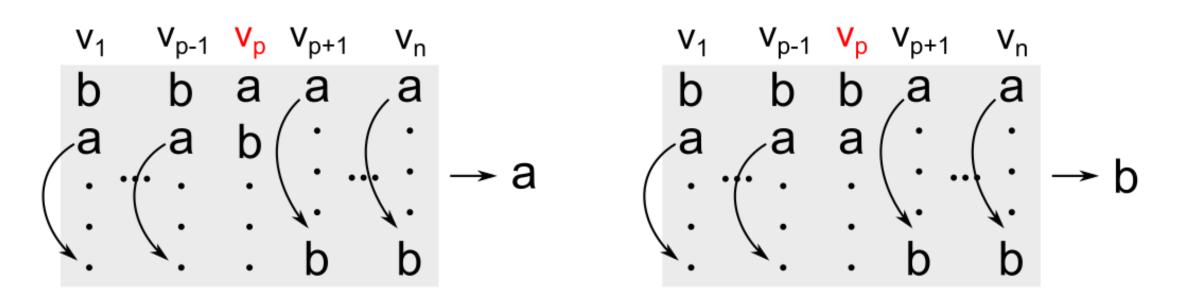


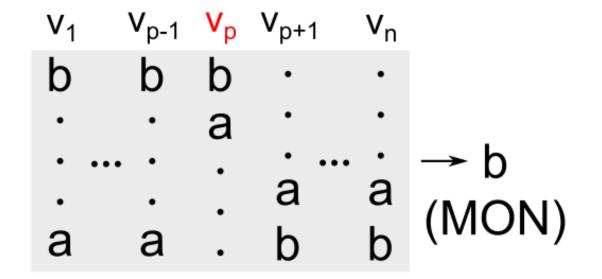


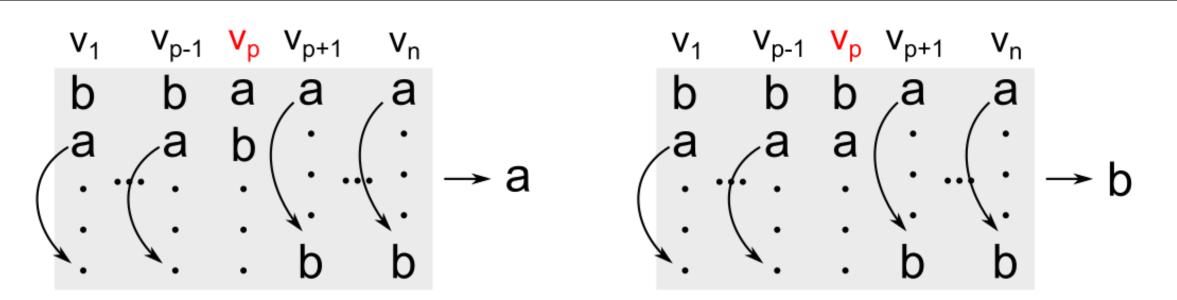


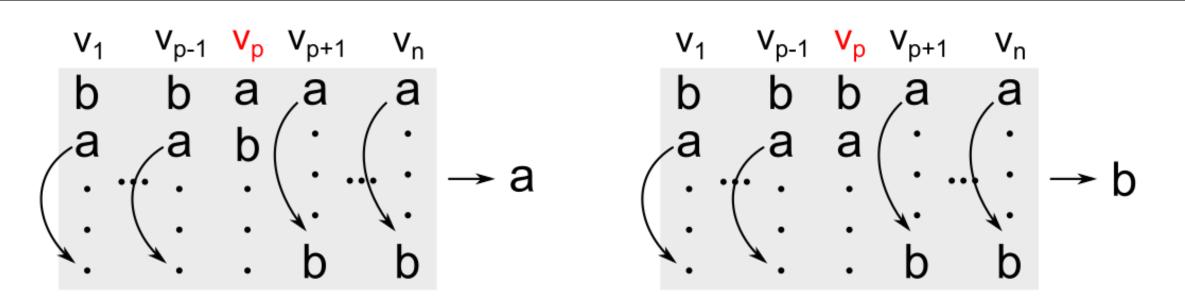
V_1	V_{p-1}	\mathbf{V}_{p}	V_{p+1}	v_n
b	b	а	•	•
•	•	b	•	•
• •	••	•	•	• •
•	•	•	а	а
a	а	•	b	b

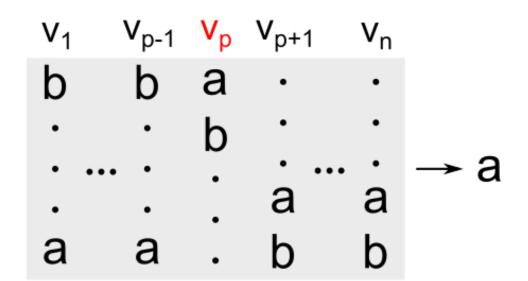






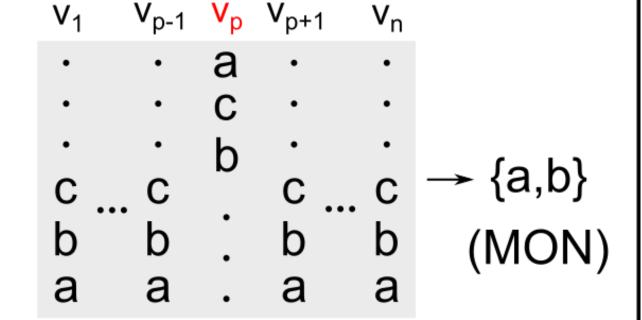


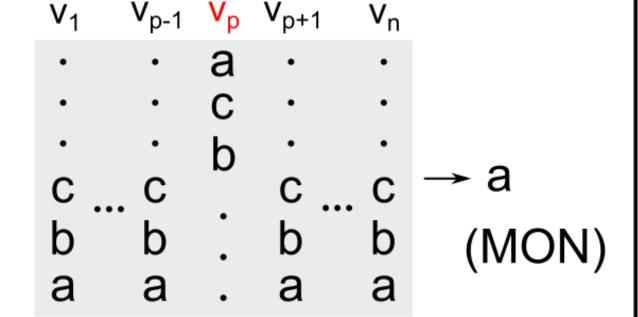


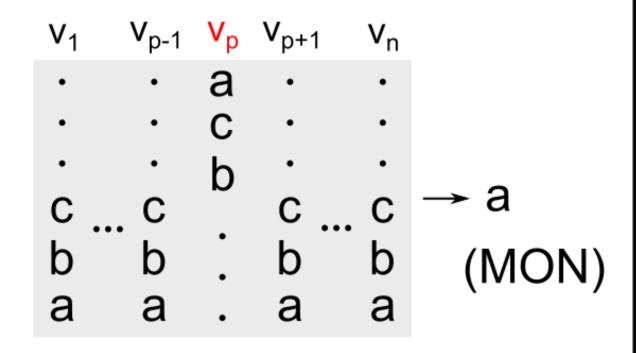


V_1	V_{p-1}	V_p	V_{p+1}	V_n
•	•	a	•	•
•	•	С	•	•
•	•	b	•	•
C	С		С.	. C
b	b	•	b	b
а	а		а	а

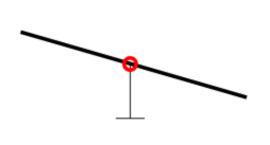
$$v_1$$
 v_{p-1} v_p v_{p+1} v_n
 \vdots a \vdots \vdots a \vdots

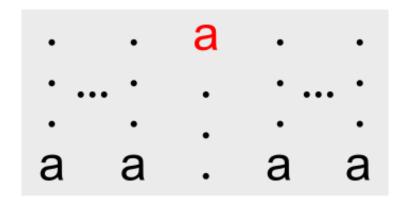






UNAN + MON ⇒ DICTATOR



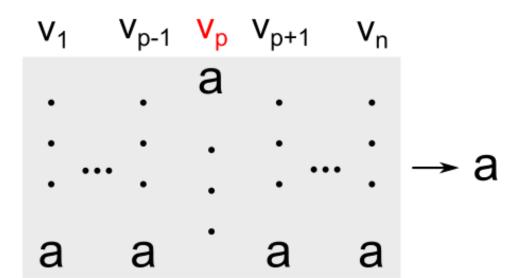


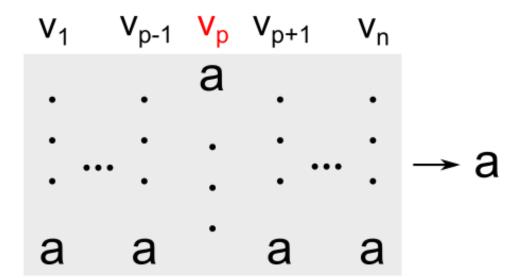


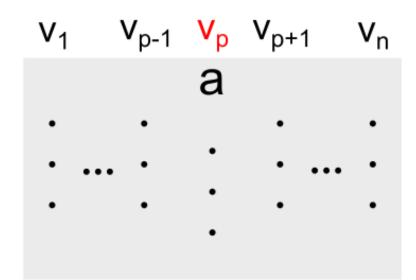
Identify a "pivotal voter" v_p for the pair $\{a,b\}$

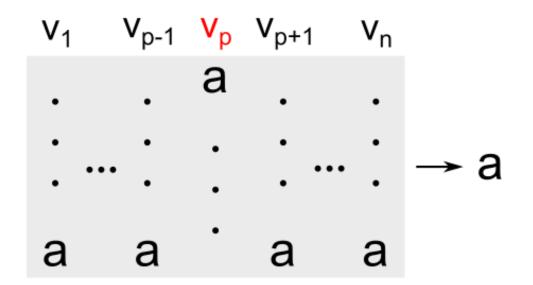
Show that "a" wins if v_p likes it even if everyone else ranks it last

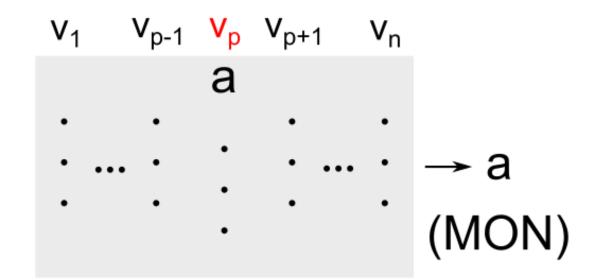
v_p is the dictator for "a" (and every other candidate)

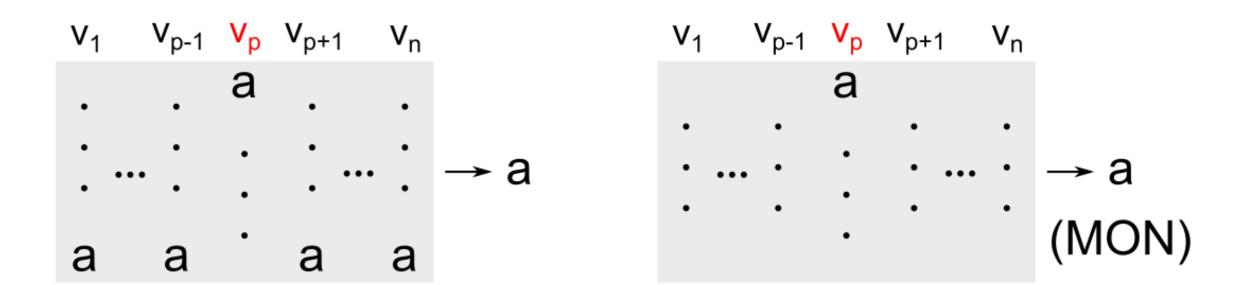




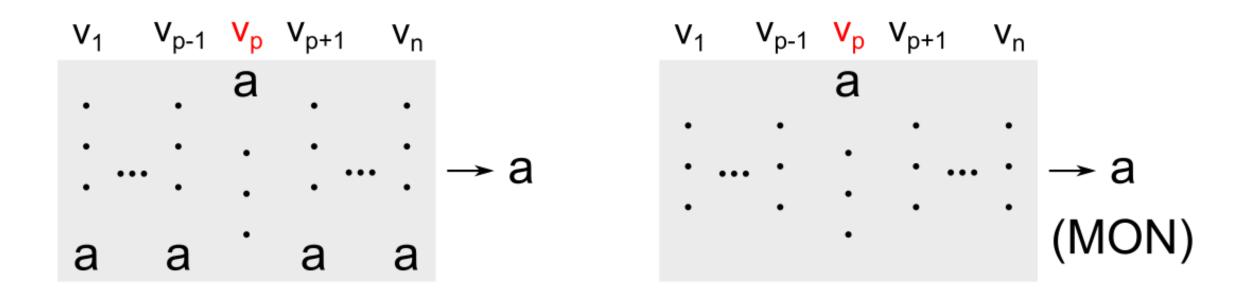




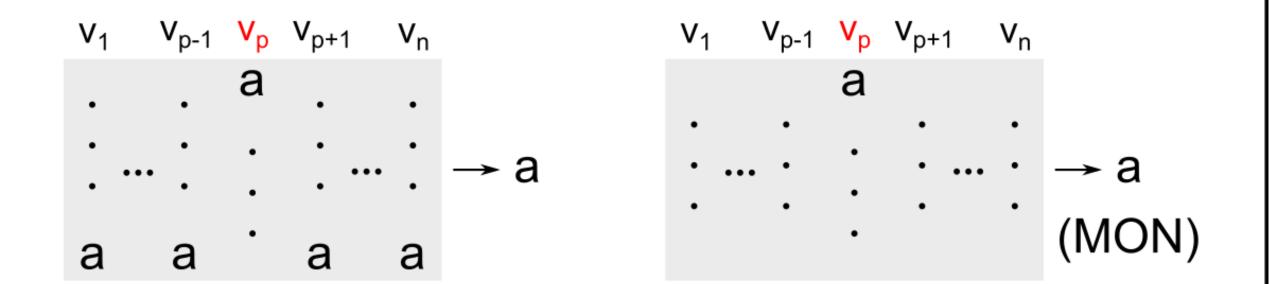




v_p is the dictator for "a"

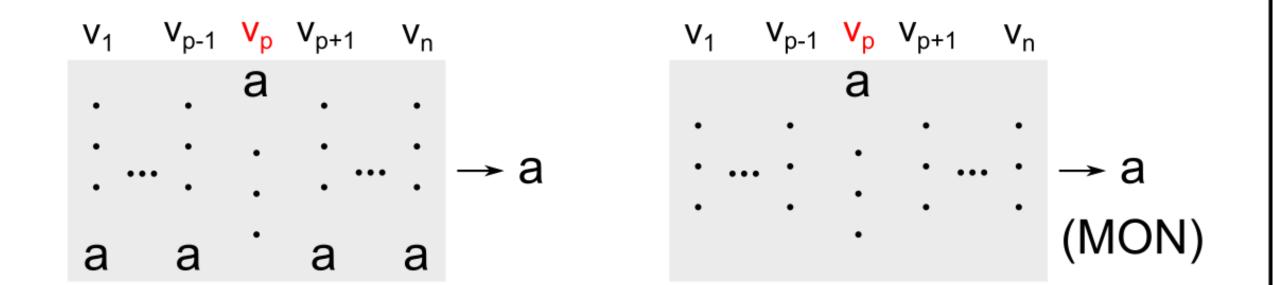


v_p is the dictator for "a"



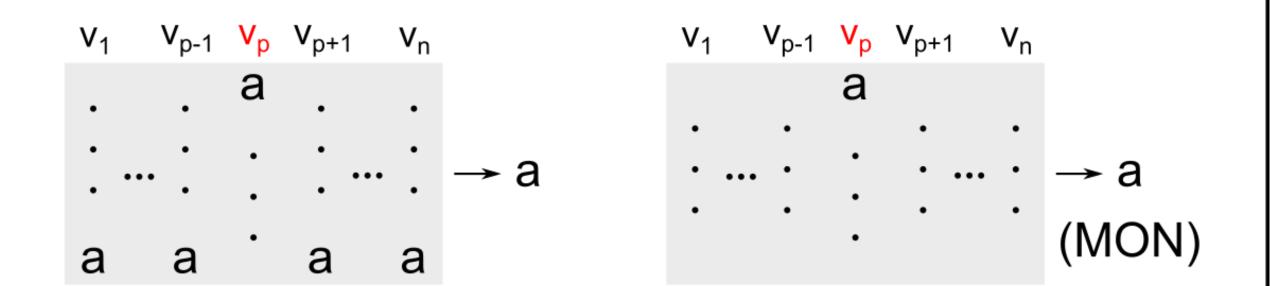
v_p is the dictator for "a"

But there can't be distinct dictators for different candidates.



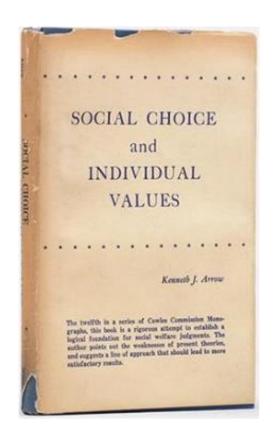
ν_p is the dictator for "a"

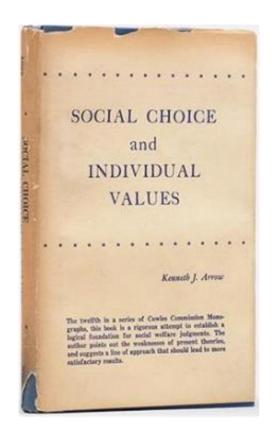
But there can't be distinct dictators for different candidates. So, v_{p} is the dictator for all candidates.



ν_p is the dictator for "a"

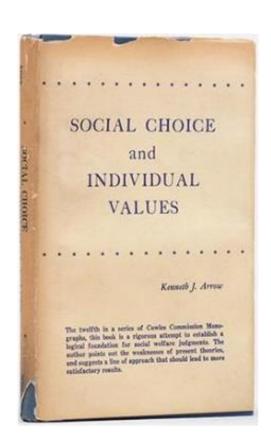
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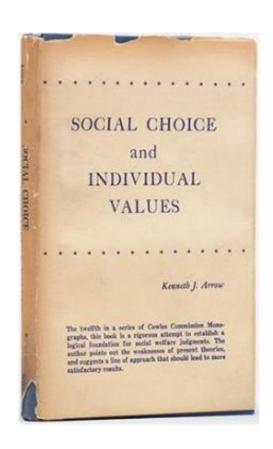
1951 PhD thesis of Kenneth Arrow





1951 PhD thesis of Kenneth Arrow



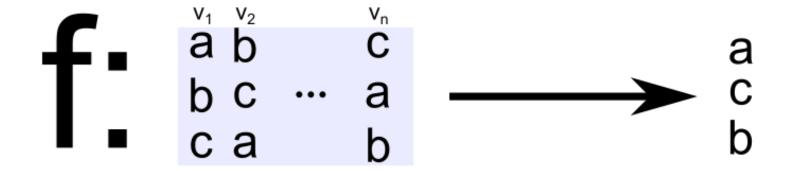




1951 PhD thesis of Kenneth Arrow

Voting Rules that Output Rankings

A mapping from preference profiles to rankings over candidates.



(also known as a social welfare function or SWF)

UNANIMOUS

If all voters prefer "a" over "b", then so does the SWF.

$$\mathbf{f}\left(\begin{array}{cccc} v_1 & v_2 & & v_n \\ \vdots & a & & \vdots \\ a & b & \dots & a \\ b & \vdots & \ddots & b \\ \vdots & \vdots & \ddots & \vdots \end{array}\right) \mathbf{f}$$

INDEPENDENCE OF IRRELEVANT ALTERNATIVES

If the relative ranking of "a" and "b" in each vote is unchanged, then their relative ranking in the SWF outcome is also unchanged.

$$\mathbf{f} \left(\begin{smallmatrix} v_1 & v_2 & & v_n \\ & b & & \\ a & a & \dots & a \\ b & & & \\ b & & & \\ \end{smallmatrix} \right) \mathbf{h}$$

DICTATORSHIP

An SWF that mimics the preferences of a fixed voter on all inputs.

DICTATORSHIP

An SWF that mimics the preferences of a fixed voter on all inputs.

A dictatorship is unanimous and IIA.



[Arrow'51]

With three or more candidates, an SWF is unanimous and IIA if and only if it is a dictatorship.



[Arrow'51]

With three or more candidates, an SWF is unanimous and IIA if and only if it is a dictatorship.

Proof almost identical to the one we saw for Gibb-Satt thm.

Sort: ♦	+	\$	\$	\$	\$	\$	\$	\$	\$	\$	\$	\$	\$	+	\$	\$	\$	\$	\$	+	+	\$	
Criterion	Majority	Maj.	Mutual	Condorcet	Cond.	Smith/	LIIA	IIA	Cloneproof	Monotone	Consistency	Participation	Reversal	Polyt	ime/	Summable	Later-no-		No favorite	Ballot	Ra	Ranks	
Method	inajority	loser	maj.	Johnson	loser	ISDA			Gioriopicoi	monotono	Control	T di doipadon	symmetry	resolv	able	Gammasio	Harm	Help	betrayal	type	=	>2	
Approval	Rated ^[a]	No	No	No ^{[b][c]}	No	No ^[b]	Yes	Yes ^[d]	Yes ^[e]	Yes	Yes	Yes	Yes	O(N)	Yes	O(N)	No	Yes ^[f]	Yes	Approvals	Yes	No	
Borda count	No	Yes	No	No ^[b]	Yes	No	No	No	Teams	Yes	Yes	Yes	Yes	O(N)	Yes	O(N)	No	Yes	No	Ranking	Yes	Yes	
Bucklin	Yes	Yes	Yes	No	No	No	No	No	No	Yes	No	No	No	O(N)	Yes	O(N)	No	Yes	If equal preferences	Ranking	Yes	Yes	
Copeland	Yes	Yes	Yes	Yes	Yes	Yes	No	No ^[b]	Teams, crowds	Yes	No ^[b]	No ^[b]	Yes	O(N ²)	No	O(N ²)	No ^[b]	No	No ^[b]	Ranking	Yes	Yes	
IRV (AV)	Yes	Yes	Yes	No ^[b]	Yes	No ^[b]	No	No	Yes	No	No	No	No	O(N ²)	Yes ^[g]	O(N!) ^[h]	Yes	Yes	No	Ranking	No	Yes	
Kemeny-Young	Yes	Yes	Yes	Yes	Yes	Yes	Yes	No ^[b]	Spoilers	Yes	No ^[b] [i]	No ^[b]	Yes	O(N!)	Yes	O(N ²) ^[j]	No ^[b]	No	No ^[b]	Ranking	Yes	Yes	
Highest median/Majority judgment ^[k]	Rated ^[l]	Yes ^[m]	No ^[n]	No ^{[b][c]}	No	No ^[b]	Yes	Yes ^[d]	Yes	Yes	No ^[o]	No ^[p]	Depends ^[q]	O(N)	Yes	O(N) ^[r]	No ^[s]	Yes	Yes	Scores ^[t]	Yes	Yes	
Minimax	Yes	No	No	Yes ^[u]	No	No	No	No ^[b]	Spoilers	Yes	No ^[b]	No ^[b]	No	O(N ²)	Yes	O(N ²)	No ^{[b][u]}	No	No ^[b]	Ranking	Yes	Yes	
Plurality/FPTP	Yes	No	No	No ^[b]	No	No ^[b]	No	No	Spoilers	Yes	Yes	Yes	No	O(N)	Yes	O(N)	N/A ^[v]	N/A ^[v]	No	Single mark	N/A	No	
Score voting	No	No	No	No ^{[b][c]}	No	No ^[b]	Yes	Yes ^[d]	Yes	Yes	Yes	Yes	Yes	O(N)	Yes	O(N)	No	Yes	Yes	Scores	Yes	Yes	
Ranked pairs	Yes	Yes	Yes	Yes	Yes	Yes	Yes	No ^[b]	Yes	Yes	No ^[b]	No ^{[p][b]}	Yes	O(N ³)	Yes	O(N ²)	No ^[b]	No	No ^{[p][b]}	Ranking	Yes	Yes	
Runoff voting	Yes	Yes	No	No ^[b]	Yes	No ^[b]	No	No	Spoilers	No	No	No	No	O(N) ^[w]	Yes	O(N) ^[w]	Yes	Yes ^[x]	No	Single mark	N/A	No ^[y]	
Schulze	Yes	Yes	Yes	Yes	Yes	Yes	No	No ^[b]	Yes	Yes	No ^[b]	No ^{[p][b]}	Yes	O(N ³)	Yes	O(N ²)	No ^[b]	No	No ^{[p][b]}	Ranking	Yes	Yes	
STAR voting	No ^[z]	Yes	No ^[aa]	No ^{[b][c]}	Yes	No ^[b]	No	No	No	Yes	No	No	Depends ^[ab]	O(N)	Yes	O(N ²)	No	No	No ^[ac]	Scores	Yes	Yes	
Sortition, arbitrary winner ^[ad]	No	No	No	No ^[b]	No	No ^[b]	Yes	Yes	No	Yes	Yes	Yes	Yes	O(1)	No	O(1)	Yes	Yes	Yes	None	N/A	N/A	
Random ballot ^[ae]	No	No	No	No ^[b]	No	No ^[b]	Yes	Yes	Yes	Yes	Yes	Yes	Yes	O(N)	No	O(N)	Yes	Yes	Yes	Single mark	N/A	No	

"Most systems are not going to work badly all of the time."

All I proved is that all can work badly at times."



Next Time

Algorithms for Manipulating Voting Rules
Oct 8 (Sat)

References

A shared proof of Arrow's and Gibbard-Satterthwaite theorems:
 "Arrow's Theorem and the Gibbard-Satterthwaite Theorem: A Unified Approach" by Philip J. Reny.
 https://www.sciencedirect.com/science/article/pii/S0165176500003323

- Another nice proof of Arrow's theorem: https://www.youtube.com/watch?v=QLi_5LCwJ20
- The "big table" of voting rules is from the Wikipedia article "Comparison of Electoral Systems".
- A tribute to Arrow (see the end of the article for a nice anecdote):
 <u>https://www.nytimes.com/2017/02/21/business/economy/kenneth-arrow-dead-nobel-laureate-in-economics.html</u>