

Differential Privacy

ϵ

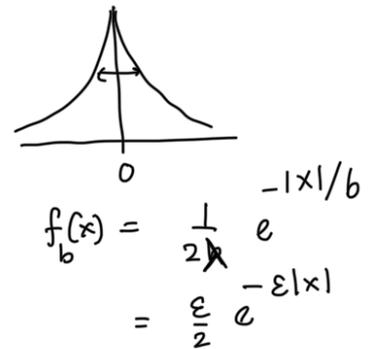
$$\begin{array}{ccc} \omega & & \omega' \\ \downarrow & & \\ x(\omega) & \sim & x(\omega') \end{array}$$

$$\begin{aligned} \Pr [x(\omega) \in S] &\leq e^\epsilon \Pr [x(\omega') \in S] \\ \Pr [x(\omega') \in S] &\leq e^\epsilon \Pr [x(\omega) \in S] \end{aligned}$$

Every cont. dist. has a density

$$\mathcal{X} = \{0, 1, 0, 1, \dots, 1\}$$

$$x(\omega) = \frac{\#1's}{T} + \underbrace{\text{Noise}}_{\sim \frac{1}{\epsilon}} f_{\lambda}$$



D.P. vs Utility

$$\begin{aligned} \Pr [x(\omega) \in S] &= \Pr [\text{Noise} + T \in [a, b]] \\ &= \Pr [\text{Noise} \in [a-T, b-T]] = \int_{a-T}^{b-T} \frac{1}{2\lambda} e^{-|x|/\lambda} dx \end{aligned}$$

$$\Pr [x(\omega) \in S] = \int_{a-T'}^{b-T'} \frac{1}{2\lambda} e^{-|x|/\lambda} dx$$

$$\begin{aligned} &\downarrow \\ &= \int_{a-T}^{b-T} \frac{1}{2\lambda} e^{-|x+\Delta|/\lambda} dx \end{aligned}$$

$\Delta = T - T'$
 $y = x - \Delta$
 $\downarrow a - T' - T + T' = a - T$

It is enough to show:

$$\lambda = \frac{1}{\epsilon} \quad \left| \frac{e^{-|x|/\lambda}}{e^{-|x+\Delta|/\lambda}} \right| \leq e^\epsilon$$

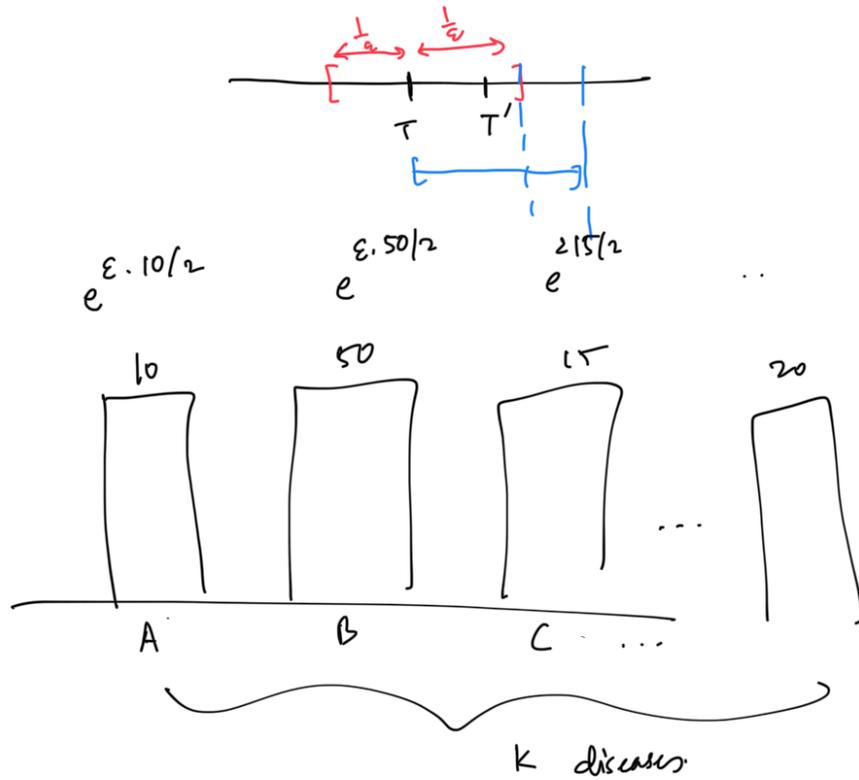
$$\leq e^{+\frac{\Delta}{\lambda}} \quad e^{-\frac{|x|}{\lambda} + \frac{|x+\Delta|}{\lambda}}$$

In general, Noise \sim Laplac $\left(\frac{\Delta}{\epsilon}\right)$, where $\Delta = \max_{\omega, \omega'} |x(\omega) - x(\omega')|$

$$\Pr \left[X > \frac{t}{e} \right] \leq e^{-t}$$

\downarrow
 $\text{Lap}(\frac{1}{e})$

Output: True Value + Noise $[-\frac{1}{e}, \frac{1}{e}]$



Thm: Suppose $X: \mathcal{D} \rightarrow \mathbb{R}^k$ $\mathcal{D}, \mathcal{D}'$: differ in only one "row"

$$\Delta := \sum_{i=1}^k |x_i(\mathcal{D}) - x_i(\mathcal{D}')| = \|X(\mathcal{D}) - X(\mathcal{D}')\|_1$$

$$X(\mathcal{D}) + \underbrace{(N_1, \dots, N_k)}_{\text{each of them from } \text{Lap}(\frac{\Delta}{2})}$$

Disease which is most prevalent?

Post processing cannot decrease D.P.





$$\Pr[\gamma(\mathcal{X}) = v] \sim \Pr[\gamma(\mathcal{X}') = v]$$

$$\Pr[\chi(\mathcal{X}) \in g^{-1}(v)] \quad \Pr[\chi(\mathcal{X}') \in g^{-1}(v)]$$

Auction:

n people

$$v_1, \dots, v_n \in [0, 1]$$

$$\text{Revenue} = p \cdot [\# \text{ people whose } v_i \geq p]$$

How much noise should we add?

— X —————>

[McSherry & Talwar '07] Exponential Mechanism.

* We have a set R of size r .
set of diceses, ...

Each outcome r has a utility $u(r)$

wanted to find $r \in R$ which has highest $u(r)$ $\leftarrow r^*$

Output $r \in R$ with probability proportional to $e^{u(r)/2}$

$$= \frac{e^{u(r)/2}}{\sum_{r' \in R} e^{u(r')/2}}$$

This mech. is ϵ -diff. private.

Erms? It is very likely that $u(r) \geq u(r^*) - \frac{\ln |R|}{\epsilon}$

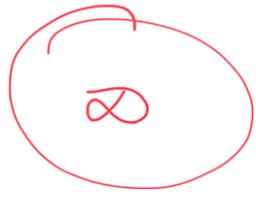
Composition Theorem



X
→

ϵ_1 -DP

$$e^{\epsilon_1} \cdot e^{\epsilon_2} = e^{\epsilon_1 + \epsilon_2}$$



ϵ_2 -D.P.